## CS 361: Theory of Computation

Assignment 3 (due 09/24/2025)

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## LATEX Source for Solutions: https://www.overleaf.com/read/wrpptzxwzfkx#835158

**Problem 1.** In this problem, we are going to prove the only-if direction of the Myhill-Nerode theorem. Let L be a language over the alphabet  $\Sigma$ . We will show that if the relation  $\equiv_L$  over  $\Sigma^*$  has finite number of equivalence classes then L is regular. Let k be the number of equivalence classes defined by  $\equiv_L$  over  $\Sigma$ . Let the set C contain exactly one representative string from each of the k equivalence classes, that is, |C| = k and every pair of strings in C are mutually distinguishable wrt L. Let  $e: \Sigma^* \to C$  be an onto function that maps any string in  $\Sigma^*$  to its representative element from C.

We now construct a DFA  $M = (Q, \Sigma, \delta, q_0, F)$  for L with exactly k states. Let Q = C, that is, create k states, labeled with strings from C.

(a) Identify the start and accept states of M in terms of the function e and language L.

Solution.  $q_0 = \boxed{\text{STATE HERE}}$ 

 $F = \boxed{\text{STATE HERE}}$ 

(b) State the transition function  $\delta$  in terms of the function e.

Solution.

(c) Consider a string  $w = w_1 w_2 \cdots w_n$ , where each  $w_i \in \Sigma$ . Let  $w[0] = \varepsilon$  and  $w[i] = w_1 w_2 \cdots w_i$ . What state is M in after reading i symbols from w, where  $i = 0, \ldots, n$  (in terms of the function e)?

Solution. STATE HERE

(d) Using part (c) above, give a one-sentence justification of why M recognizes L.

Solution.

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**Problem 2.** Use the Myhill Nerode theorem to answer the following questions about the language  $L_k$  defined for a fixed number  $k \geq 1$  and  $\Sigma = \{a\}$ .

$$L_k = \{ a^{\ell} \mid \ell \text{ is a multiple of } k \}$$

(a) State all the equivalence classes of the relation  $\equiv_L$  over  $\Sigma^*$ . Prove that they form equivalence classes—that is—any pair of strings in the same class are indistinguishable wrt L, any pair of strings in different classes are mutually distinguishable wrt L and finally that the classes partition  $\Sigma^*$ .

Solution.

(b) Using (a) argue that  $L_k$  is regular. How many states does a minimal DFA for L require? Solution.

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**Problem 3.** Prove that the following languages are not regular using the **Myhill Nerode** theorem.

(a)  $L = \{0^n 1^{m+n} 0^n \mid m, n \ge 1\}.$ 

Solution.

(b)  $L = \{w \in \{0,1\}^* \mid w = w^R \text{ and } |w| \text{ is divisible by 3}\}$ . (That is, the set of binary strings that are palindromes and have length a multiple of 3.) *Hint. The divisibility by 3 is a bit of a distraction, think of ways of avoiding it.* 

Solution.

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**Problem 4.** For each of the following languages over the alphabet  $\Sigma = \{0, 1\}$ , identify whether it is regular or not regular. If it is regular, provide a DFA, NFA or regular expression for it. If it is not regular, provide a proof using the **pumping lemma** or **closure properties**.

- (a)  $L = \{w \mid \text{ substrings 01 and 10 appear the same number of times in } w\}.$ Solution.
- (b)  $L = \{ w \mid \text{ substrings 00 and 11 appear the same number of times in } w \}.$  Solution.
- (c)  $L = \{xy \mid \text{where } x = 1^k \text{ and } y \text{ contains at least } k \text{ 1s, for some } k \geq 1\}.$ Solution.
- (d)  $L = \{xy \mid \text{ where } x = 1^k \text{ and } y \text{ contains at most } k \text{ 1s, for some } k \geq 1\}.$ Solution.