Reduction Proofs: Solution

— Halt-no-Input —

Prove that the Halt-no-Input problem is undecidable.

Halt-no-Input problem: given a program P that requires no input, does P halt?

We want to prove that the Halt-no-Input problem is undecidable. Therefore, we assume the existence of a function that solves this problem. Suppose Haltni is a function that computes a solution to the Halt-no-Input problem:

```
let HaltNI (P: unit -> 'a) : bool = ...
```

Although the problem does not state precisely what type P is, we are allowed to make reasonable assumptions. For example, programs are in some sense functions, which is why P has the type unit \rightarrow 'a above.

Let's now construct a reduction.

```
let Halt (P: string -> 'a) (i: string) : bool =
let P'() = P i
HaltNI P'
```

The above reduction reduces Halt-no-Input to the Halting Problem. It works by showing that we can solve the Halting Problem by changing the form of the problem slightly, then by calling Haltni. However, we know that we cannot solve the Halting Problem, because it is undecidable. Since all of the above code is perfectly legal F#, then the only flaw in our logic must have been our assumption that Haltni is decidable. Therefore, Haltni is undecidable.

Alternatively, we can think of programs as strings if it is more convenient, and we will use string manipulation to construct our reduction. Therefore, we make a different set of reasonable assumptions about HaltNI:

```
let HaltNI (P: string) : bool = ...
```

Suppose HaltNI expects that P be a function of the form let MyFunc () = ... Then we can construct a different reduction:

```
let Halt (P: string) (i: string) : bool =
let P' = "let MyFunc2 () = \n\t" + P + "\n\tMyFunc \"" + i + "\""
HaltNI P'
```

where Halt takes a P of the form let MyFunc (i: string) = ... For example, if we call Halt with P defined as let MyFunc (i: string) = "hello" + i and i defined as "world", then the reducer constructs the string P',

```
let MyFunc2 () =
let MyFunc (i: string) = "hello " + i
MyFunc "world"
```

in order to convert the problem into the form expected by Haltni. Since string manipulation is clearly legal in F#, we must again conclude that Haltni is undecidable.

Observe that in both proofs, the reduction does not solve the Halt-no-Input problem directly. Instead, the reduction algorithm changes the form of the problem, delegating computation of the solution to the function that we assumed exists (HaltNI).