## **Midterm 2 Study Guide Solutions** Handout 27

CSCI 334: Spring 2024

## $\overline{\phantom{a}}$  Solutions  $\overline{\phantom{a}\phantom{a}\phantom{a}}$

- **Q1.** (10 points) . Terminology You should do this exercise on your own.
- **Q2.** (10 points) . Decision Problem

We prove that bothHalt is not computable by contradiction. Assume bothHalt exists. Then we can write the following function.

let halt  $p$  i = bothHalt p p i

If p halts, then bothHalt will return true because "both" p programs halt. Likewise, if p does not halt, then bothHalt will return false because "both" p programs do not halt. Since we are able to construct a halt function using the ordinary rules of logic  $(F#)$  with the only assumption being the existence of bothHalt, and we know that halt is not computable, then bothHalt must not be computable.

An interesting alternative raised in class is to define a different function, like so

```
let halt p i =let f x = xbothHalt p f i
```
This also works because when p halts, f also halts (because it always halts), so bothHalt returns true. If p does not halt, bothHalts returns false because it is sufficient for either program to not halt for bothHalt to return false. Good thinking!

**Q3.** (10 points) . Parsing and Evaluation

Here is one possible solution.

```
open Combinator
type Expr =
| Increment
| Decrement
| MoveLeft
| MoveRight
| Print
type State = { tape: int list; pos: int }
let op =
     (
          (pchar '+' | \rangle fun _ -> Increment) <|>
          (\text{pchar } ' -' /\text{)} fun _ -> Decrement) <|>
          (\text{pchar } \langle \langle \cdot \rangle | \cdot) fun -> MoveLeft) \langle \cdot \rangle
```

```
(\text{pchar '>'} |>> fun _ -> MoveRight) <|>
         (pchar 'p' | \rangle fun _ -> Print)
    ) <!> "op"
let expr = pmany0 op <!> "expr"
let grammar = pleft expr peof <!> "grammar"
let rec getAtIndex xs i =
    match xs with
    | [] -> failwith "Cannot find index in list."
    | y:: when i = 0 -> y
    | _::ys -> getAtIndex ys (i - 1)
let rec updateAtIndex xs i v =match xs with
    | [] when i <> -1 -> failwith "Cannot find index in list."
    | | | \rightarrow || \angle::ys when i = 0 -> v::ys
    | y::ys -> y::updateAtIndex ys (i - 1) v
let parse input =
    let i = prepare input
    match grammar i with
    | Success(ast,_) -> Some ast
    | Failure(_,_) -> None
let ephsPrint v =
    let s =match v with
        | 0 -> "a"
        | 1 - > "b"| 2 \rightarrow "c"| 3 - > "d"| 4 - > "e"| 5 - > "f"| 6 -> "g"
        | 7 -> "h"
        | 8 -> "i"
        | 9 -> "j"
        | 10 \rightarrow "k"| 11 -> "l"
        | 12 - > "m"| 13 - > "n"| 14 - \rangle "o"
        | 15 -> "p"
        | 16 -> "q"
        | 17 \rightarrow "r"| 18 \rightarrow "s"
        | 19 -> "t"
        | 20 -> "u"
        | 21 \rightarrow "v"| 22 - > "w"| 23 \rightarrow "x"
```

```
| 24 \rightarrow "y"| 25 \rightarrow "z"
        | -> "poo"
    printf "%s" s
let evalOp (e: Expr)(s: State) : State =
    match e with
    | Increment ->
        let cur = getAtIndex s.tape s.pos
        let upd = cur + 1let tape' = updateAtIndex s.tape s.pos upd
        { s with tape = tape' }
    | Decrement ->
        let cur = getAtIndex s.tape s.pos
        let upd = cur - 1let tape' = updateAtIndex s.tape s.pos upd
        \{ s with tape = tape' }
    | MoveLeft ->
        if s.pos > 0 then
             \{ \text{ tape} = \text{s.tape}; \text{pos} = \text{s.pos} - 1 \}else
             printfn "Tape head out-of-bounds (on left)."
             exit 1
    | MoveRight ->
         if s.pos < 10 then
             \{ \text{ tape} = \text{s.tape}; \text{pos} = \text{s.pos} + 1 \}else
             printfn "Tape head out-of-bounds (on right)."
             exit 1
    | Print ->
        let cur = getAtIndex s.tape s.pos
        ephsPrint cur
        s
let rec eval (es: Expr list)(s: State) : State =
    match es with
    | | \rightarrow s
    | op::ops ->
        let s' = evalOp op seval ops s'
[<EntryPoint>]
let main args =
    let file = args[0]let input = System.IO.File.ReadAllText file
    let ast_maybe = parse input
    match ast_maybe with
    | Some ast ->
        eval ast { tape = [0;0;0;0;0;0;0;0;0;0;0]; pos = 0 } |> ignore
        0
    | None ->
        printfn "Invalid program"
        1
```
## **Q4.** (10 points) . Partial and Total Functions

- (a) Partial function:  $\{(n, \text{fibonacci } n) \mid n \in \mathbb{Z} \land x \geq 0\}$
- (b) Looking at our gcd function, the only operation that may be problematic is %. The answer I gave in class assumed mod was undefined for negative numbers, but actually, I had it backward: mod is defined for all divisors and quotients in Z. Some programming languages, like C, assume that the quotient and divisor are not negative, and will give the wrong answer otherwise. If we stick with the mathy definition (instead of C), then gcd appears to be undefined nowhere (because the case where  $\mathbf{b} = 0$  is explicitly handled and defined), so the function is total:  $\{\langle a, b, \mathbf{g} \mathbf{c} \mathbf{d} \, a \, b \rangle \mid a, b \in \mathbb{Z}\}\$
- (c) Total function:  $\{\langle x, |x| \rangle \mid x \in \mathbb{Z}\}\$

One can use function graphs to enforce preconditions when implementing the above functions. These are also a concise form of documentation that you might consider putting into a Javadoc or Python docstring. Other programmers (or "future you") will thank you.