

CSCI 334:
Principles of Programming Languages

Lecture 23: Why PL Matters

Instructor: Dan Barowy
Williams

Topics

Wildly different models of programming.
Which language do I use?

Your to-dos

1. **Mostly working** project, **due Sunday 12/10**
2. Want to talk about your project?
Last chance.
Office hours tomorrow 10-11am, 1:30-2:30pm.

Final project timeline

- ~~1. Project proposal (Lab 8), **due Sun 11/12**~~
- ~~2. Minimally working version (Lab 9), **due Sun 11/19**~~
- ~~3. Language specification doc (Lab 10), **due Sun 12/3**~~
4. Mostly working version (Lab 11), **due Sun 12/10**
5. Project + video presentation (Lab 12), **due Sun 12/17**

Announcements



CS Holiday Party

Friday, Dec 8 @ 2:35pm
CS Common Room

Join the CS faculty and your peers for an end-of-semester celebration. We will have hot cocoa and treats for you to enjoy. Last gathering of the year!

Ingalls Test for Extensibility

i.e., the “rectangle test”

- The test is about **the ability to extend software *after* it has already been designed and written.**

Java, Python, etc. pass the rectangle test

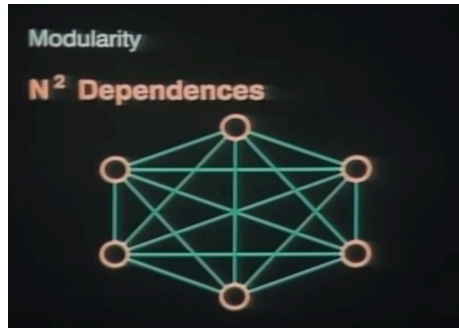
The right choice depends on the problem

- OO offers a **different kind of extensibility** than functional (or function-oriented) languages.
- Suppose you’re **modeling a hospital**.

Operation	Doctor	Nurse	Orderly
Print	Print Doctor	Print Nurse	Print Orderly
Pay	Pay Doctor	Pay Nurse	Pay Orderly

- FP makes it **easy to add operations** (**rows** above).
- OOP makes it **easy to add data** (**columns** above).

Why I like OO



OO is fundamentally based on **the idea that people matter** in the design of a programming language.

How do we **minimize human effort** while designing large pieces of software?

Prolog

Prolog

- The goal of AI is to enable a computer to answer declarative queries.
- I.e., it already knows how to answer you.
- Prolog was an attempt to solve this problem.
- Since this was early work, the input language was somewhat primitive: predicate logic.
- As you will see, formulating queries in pure logic is not the easiest thing to do.
- However, for certain classes of logic, there are known efficient, deterministic algorithms for solving every possible query.

Logic Programming



- Logic programming began as a collaboration between AI researchers (e.g., John McCarthy) and logicians (e.g., John Alan Robinson) to solve problems in artificial intelligence.
- Cordell Green built the first "question and answer" system using Robinson's "unification algorithm," demonstrating that it was practical to prove theorems automatically.

Prolog

- Alain Colmerauer and Phillippe Roussel at Aix-Marseille University invented Prolog in 1972.
- They were inspired by a particular formulation of logic, called "Horn clauses," popularized by the logician Robert Kowalski.
- Horn clauses have a "procedural interpretation," meaning that they suggest a simple procedure for solving them, called "resolution."
- John Alan Robinson's unification algorithm is an efficient algorithm for doing resolution, and this is essentially the algorithm used by Prolog.



Horn Clause

- Horn clauses are composed of two simple pieces:
 - facts
 - rules (clauses)
- Rules are composed of facts
- Complex facts may also be composed using conjunction.
- We will explore these concepts using Prolog syntax.
- Note that Horn clauses can be "satisfied" in polynomial time.
- In fact, Horn logic is the most expressive form of logic known to be satisfiable in polynomial time.

Facts (Prolog syntax)

- Here are some facts:

```
raining.  
cloudy.  
thursday.
```

- Facts are assumed to be true.
- Facts of this form are sometimes called "atoms", since they are indivisible.
- The meaning of these facts is up to the programmer.
- Facts can also be compound:

```
raining,cloudy.  
cloudy,thursday.
```

- ", " denotes "logical and".
- Note that, in Prolog, facts are always lowercase and must begin with a letter.

Rules (Prolog syntax)

- Here are some rules:

```
sleep_deprived :- thursday.  
unhappy :- raining,cloudy.
```

- The interpretation of a rule $x :- Y$ is:
if Y is true, then x is true
- In other words, Y is the antecedent and x is the consequent.
- So, we might interpret the above as:
"students are sleep deprived if it is Thursday"
"I am unhappy if it is raining and cloudy."

Variables (Prolog syntax)

- Note that I just used a generalization of rules without definition:

```
X :- Y
```

- Prolog explicitly allows generalizations of facts like this.
- We call these generalizations "variables", because their precise values (i.e., facts) may not be known to us.
- In the "execution" of a Prolog program, we seek to "instantiate" variables with facts.
- In Prolog, variables are always written starting with an uppercase letter.
- We will come back to variables shortly.

Complex facts (Prolog syntax)

- Prolog allows one additional form:

```
musician(mia).  
musician(john).  
friends_with(mia, john).
```

- Statements of this form are called "complex facts."
- Again, the interpretation is up to you.
- E.g.,
"Mia is a musician."
"John is a musician."
"Mia is friends with John."
- Note that we do not automatically assume that
"John is friends with Mia!"

Queries

- Taken together, facts and rules form a "knowledge base."

```
raining.  
cloudy.  
thursday.  
sleep_deprived :- thursday.  
unhappy :- raining, cloudy.
```

- A query asks the knowledge base a question. E.g.,

```
?- sleep_deprived.  
true  
?- unhappy.  
true
```

Resolution

- "Resolution" is the name of the procedure that Prolog uses to "satisfy" a query.

```
raining.  
cloudy.  
thursday.  
sleep_deprived :- thursday.  
unhappy :- raining, cloudy.
```

- Essentially, we seek to reduce a query expression to the expression true by substitution.
- Remember that facts are assumed to be true.

Resolution

1. raining.
2. cloudy.
3. thursday.
4. sleep_deprived :- thursday.
5. unhappy :- raining,cloudy.

?- sleep_deprived.

- For a given query, we first seek either a fact that immediately makes the query true, or we seek a rule whose consequent is the query.

- When a rule is reduced to the form $X \text{ :- true}$, then X is true.

6. sleep_deprived :- thursday (by KB4)
7. sleep_deprived :- true (by KB3)
8. true (by deduction)

Resolution

- Given the following knowledge base,

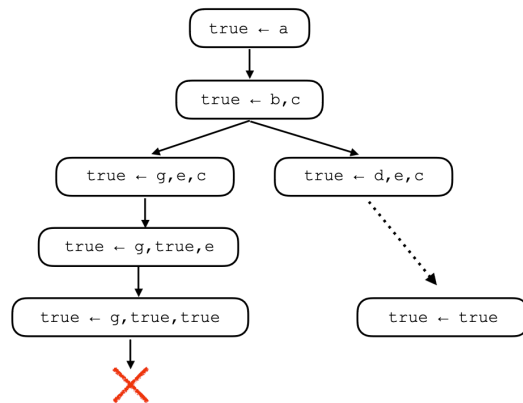
1. a :- b,c.
2. b :- d,e.
3. b :- g,e.
4. c :- e.
5. d.
6. e.
7. f :- a,g.

- let's try to satisfy the following query using resolution:

?- a.

Proof Search

- Nonetheless, Prolog is not generally sensitive to the order of the facts in a database. How does this work?
- The answer is that resolution is actually a form of backtracking search.



Resolution with Variables

- Resolution with variables can be very computationally expensive.
- Unification allows resolution with variables to be completed in polynomial time.
- The basic insight is to "instantiate" variables "on demand" instead of enumerating all possible variable instantiations into facts.
- Hindley-Milner is essentially just unification.

1. musician(mia).
2. musician(john).
3. friends_with(X,Y) :- musician(X),musician(Y).

- Let's resolve the following query:

?- friends_with(mia, john).

Resolution with Variables

- When asking a query that utilizes variables, Prolog will both search for a satisfying assignment and it will return that assignment.
- There may be more than one possible assignment.
- If so, use the ";" command to ask for another solution.
- Let's resolve the following query:

```
?- friends_with(mia,Who).
```

- We may even ask:

```
?- friends_with(Who1,Who2).
```

Exercise

- Construct the a knowledge base containing the following facts:
 - "Giants eat people."
 - "Giants eat bunnies."
 - "Bunnies eat grass."
 - "People eat bunnies."
 - "People eat people."
 - "Those who are eaten by others hate those others."
 - "Monsters love those who hate themselves."
- Then supply a query that can answer:
 - "Who do monsters love?"

Turing Tarpit

A **Turing tarpit** is a programming language flexible enough to do anything (i.e., it is **Turing equivalent**) while also being **difficult to learn and use** for everyday tasks.

"Beware of the Turing tar-pit in which everything is possible but nothing of interest is easy." —Alan Perlis

Examples:

- Turing machines
- The Lambda Calculus
- Maybe Prolog?
- C?

FlashMeta

Programming **of** the People, **by** the People, and **for** the People

Daniel Barowy
UMassAmherst

Williams College, January 9, 2017

```
1 module Parser
2 where
3
4 import Data.Char
5 import Control.Monad
6
7 newtype Parser a = P (String -> [(a, String)])
8
9 instance Monad Parser where
10   — identity
11   return v = P (\inp -> [(v,inp)])
12   — bind
13   p >>= f = P (\inp -> case parse p inp of
14     [] -> []
15     _ -> parse (f v) out)
```

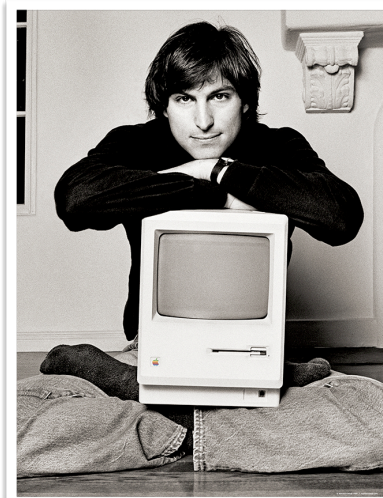


Hang on, let me fire up VSCode.



In the future: There are **no** programmers.

A Bicycle for the Mind



This work is not done yet.

Some things to think about

Would it be better to

- Have a programming language where **bugs are impossible**, but **programming is difficult**, or
- have a programming language where **bugs are possible** but their **consequences are minimized**?

Coq

```
296 Defined.
297
298 (* 3.2 *)
299 theorem mod_of_multiple_term:
300   forall n m p: nat,
301   (n * p + m) mod p = m mod p.
302 Proof.
303 intros n m p.
304 destruct (eq_nat_dec p 0) as [H|H1]: [rewrit
305 pose (H2 := div_mod_rep (n * p + m) p H1).
306 rewrite (div_mod_rep m p H1) in H2 at 1.
307 rewrite plus_assoc in H2.
308 rewrite <- Nat.mul_add_distr_r in H2.
309 symmetry.
310 refine (mod_unique _ _ H2); apply m
311 Qed.
312
313 (* 2.3 *)
314 theorem mod_plus:
315   forall n m k: nat,
316   (n + m) mod k = (n mod k + m mod k) mod k.
317 Proof.
318 intros n m [k]; [tauto].
319 assert (H1: S k <= 0); [discriminate].
320 rewrite (div_mod_rep (n + m) k H1) at 1.
```

TypeScript

```
1 interface Person {
2   firstName: string;
3   lastName: string;
4 }
5
6 function greeter(user: Person): string {
7   return "Hello " + user.firstName + " " + user.lastName;
8 }
9
10 let john: Person = {
11   firstName: "John",
12   lastName: "Doe"
13 }
14 console.log(greeter(john));
15
16
```

1: bash

Hello John Doe

Would it be better to

- Have **one language to rule them all**, or
- have many **different, small special-purpose languages**?

Clojure

```
(ns poetry.core
  (:require [clj-http.client :as http]
            [clojure.string :as str]))

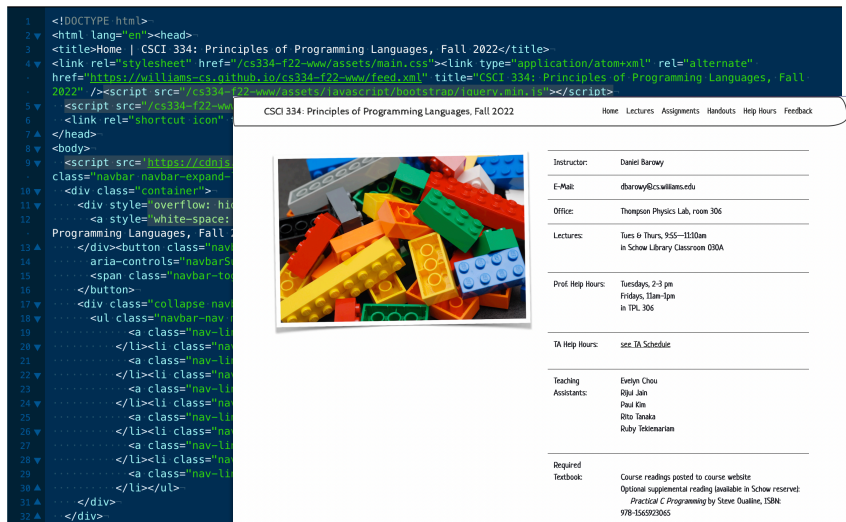
(def haiku-url
  "http://search.twitter.com/search.json?q=%23haiku")

(defn raw-haikus []
  (-> (http/get haiku-url {:as :json})
      :body
      :results
      (map :text)))

(defn trim-lines [s]
  (-> (str/split-lines s)
      (map str/trim)
      (remove str/blank?)
      (str/join "\n")))

(defn sanitize-haiku [haiku]
  (-> haiku
      (str/replace #"RT" "")
      (str/replace #"#\W+" "")
      (str/replace #"@\W+?" ""))
```

HTML



The image shows a screenshot of an HTML document. On the left, the raw HTML code is visible, including tags for <doctype>, <html>, <head>, <title>, <link>, <script>, <body>, <div>, , and </div>. On the right, the rendered page is shown, featuring a navigation menu with links for Home, Lectures, Assignments, Readouts, Help hours, and Feedback. Below the menu is a photo of colorful building blocks (LEGO bricks) in various colors like red, yellow, green, and blue.

And if “many languages” is the answer,
would it be better to

- Have many **standalone languages**, or
- have many languages that **can be embedded in a host language**?

SQL

```
mysql> describe book_stock;
+-----+-----+-----+-----+-----+-----+
| Field | Type | Null | Key | Default | Extra |
+-----+-----+-----+-----+-----+-----+
| book_id | int(6) | YES | | NULL | |
| book_qty | int(6) | YES | | NULL | |
| booktype | varchar(20) | YES | | NULL | |
+-----+-----+-----+-----+-----+-----+
3 rows in set (0.04 sec)

mysql> insert into book_stock values(1001,20,'education');
Query OK, 1 row affected (0.00 sec)

mysql> select * from book_stock;
+-----+-----+-----+
| book_id | book_qty | booktype |
+-----+-----+-----+
| 1001 | 20 | education |
+-----+-----+-----+
1 row in set (0.07 sec)

mysql>
```

LINQ

```
NorthwindDataContext db = new NorthwindDataContext();
var products = from p in db.Products
                where p.Category.CategoryName == "Beverages"
                select p;
```

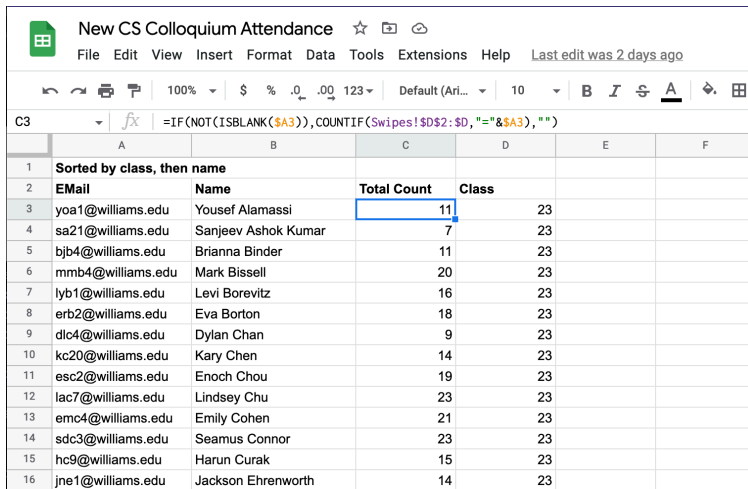
Would it be better to

- Leave programming to the **experts**, or
- let **anybody** do it?

C++

```
494 // Look for a matching path substitution and return the path this command should use
495 string Command::substitutePath(string p) noexcept {
496     auto iter = _current_run._substitutions.find(p);
497     if (iter == _current_run._substitutions.end()) return p;
498
499     LOG(exec) << this << ": Replacing path " << p << " with " << iter->second;
500
501     return iter->second;
502 }
503
504 // Inform this command that it used a temporary file
505 void Command::addTempfile(shared_ptr<Artifact> tempfile) noexcept {
506     // Add the tempfile and mark it as not accessed (the value in the map)
507     _current_run._tempfiles.emplace(tempfile, false);
508 }
509
510 // Get a reference from this command's reference table
511 const shared_ptr<Ref>& Command::getRef(Ref::ID id) noexcept {
512     ASSERT(id >= 0 && id < _current_run._refs.size())
513     << "Invalid reference ID " << id << " in " << this;
514     ASSERT(!_current_run._refs[id] << "Access to null reference ID " << id << " in " << this;
515     return _current_run._refs[id];
516 }
517
518 // Store a reference at a known index of this command's local reference table
519 void Command::setRef(Ref::ID id, shared_ptr<Ref> ref) noexcept {
520     ASSERT(ref) << "Attempted to store null ref at ID " << id << " in " << this;
521
522     // Are we adding this ref onto the end of the refs list? If so, grow as needed
523     if (id >= _current_run._refs.size()) _current_run._refs.resize(id + 1);
```

Spreadsheets



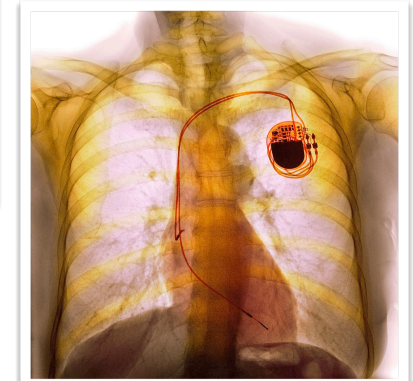
New CS Colloquium Attendance

File Edit View Insert Format Data Tools Extensions Help Last edit was 2 days ago

C3 =IF(NOT(ISBLANK(\$A3)),COUNTIF(Swipes!\$D\$2:\$D,"="&\$A3),"")

	A	B	C	D	E	F
1	Sorted by class, then name					
2	Email	Name	Total Count	Class		
3	yoa1@williams.edu	Yusef Alamassi	11	23		
4	sa21@williams.edu	Sanjeev Ashok Kumar	7	23		
5	bjb4@williams.edu	Brianna Binder	11	23		
6	mmb4@williams.edu	Mark Bissell	20	23		
7	lyb1@williams.edu	Levi Borevitz	16	23		
8	erb2@williams.edu	Eva Borton	18	23		
9	dlc4@williams.edu	Dylan Chan	9	23		
10	kc20@williams.edu	Kary Chen	14	23		
11	esc2@williams.edu	Enoch Chou	19	23		
12	lac7@williams.edu	Lindsey Chu	23	23		
13	emc4@williams.edu	Emily Cohen	21	23		
14	sdc3@williams.edu	Seamus Connor	23	23		
15	hc9@williams.edu	Harun Curak	15	23		
16	jne1@williams.edu	Jackson Ehrenworth	14	23		

PL matters because computers
are everywhere



It's up to us how we want
to use our machines



"A tasteful watercolor painting of a person pondering what to do with a computer."
(DALL·E, Dec 2022)

Which language do I use?

Next steps

(aka, some things to do over the summer)

- **Teach yourself** another programming language.
- Dig in to **a problem that bugs you**.
(me: I've always wanted to write a computer algebra solver)
- **Keep playing** with your project! It's yours! (and you can **show it off** to interviewers)
- Most of all, **do something that excites you**.