

# Heaps and Priority Queues

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# Admin

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- Final review Friday (no quiz!)
- If you have an “exam hardship” let me know as soon as possible
- Lab 6 graded; back any minute now. Lab 7 waiting on a backend issue; probably later today
  - Labs 8 and 9, and last quiz, should be soon...
- Please bring your computers (or something) on Wednesday for course evals at end of class
- Any questions?

# Adjacency List vs Adjacency Matrix

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- Adjacency List is (often) much faster for listing neighbors of a vertex:
  - Adjacency Matrix gives time proportional to the total number of vertices, Adjacency List gives time proportional to the degree.
- Adjacency Matrix is much faster for looking up if there is an edge between two vertices
- Adjacency List (on a graph with  $n$  vertices and  $m$  edges) is much more space efficient if  $m < n^2$

# Shortest Path in Graph

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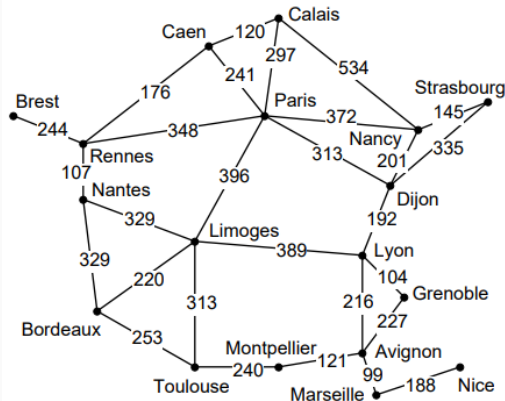
## Shortest path in a graph

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- Breadth-first search finds the path with the smallest length in a graph
- Let's look at a visualization of this

## Shortest path: what do we really want?

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- Not all edges are the same!
- It takes a different amount of time to travel down different roads, or take different flights
- What if we have numerical labels, and the length of the path is the sum of the labels?

## Trying out BFS

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- Does BFS work if we want to take path labels into account?
- Can we come up with an example where it doesn't?

# What do we want?

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- Want to explore the paths *in order of length*
- So: something like BFS, but want to explore the shortest path next
- Called Dijkstra's algorithm. We'll discuss in detail on Wednesday
- Let's look at how it runs
- For now: what do we need to implement this?



## New operation on a data structure

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- Need to know: what is the **shortest remaining path to explore**
- In BFS, we kept all vertices we wanted to explore in a queue
- Now, we don't want first in first out. We want the vertex with **smallest path length** out.
- We want to keep a collection of vertices in a data structure, with the ability to remove the smallest
- This is called a *priority queue*

## **Priority Queues (Slide Deck Change)**

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# Benefits of heaps?

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- Balanced binary search tree: time for removeMin? add?
  - Both  $O(\log n)$
- Heap?
  - Both  $O(\log n)$
  - But *much* better constants, much simpler
- Creating a balanced binary tree of size  $n$  using an unsorted Vector?
  - $O(n \log n)$
- Creating a heap of size  $n$  using an unsorted Vector?
  - $O(n)$

# Heap vs priority queue

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- Priority queue is the interface
- Heap is the specific implementation
- Like Map vs Hashtable. There are other ways to implement a Map; similarly, there are other ways to implement a priority queue

# Summary

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- In short: heaps are much simpler and have much better constants
- Extremely common in practice!
- HeapSort is one of the most common sorting methods, especially if you want  $O(n \log n)$  guaranteed worst-case running time
- We saw min heaps. Can get a “max heap” by flipping the requirement: the root element must be largest in the heap. Then can get good `removeMax` performance