

CSCI 136:
Data Structures
and
Advanced Programming

Lecture 10
Recursion, part 2

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Topics

- Quiz 2—best case
- Recursion costs
- Mathematical Induction

Your to-dos

1. Lab 3, **due Tuesday 3/1 by 10pm**
2. Read **before Wed**: Bailey, Ch 9.4–9.5.

Announcements

- Lab 1: feedback today
- Lab 1: if feedback has mistakes...



What is the “best case”?

```
public static <E> void rev(Vector<E> orig, Vector<E> flip) {  
    for (int i = orig.size() - 1; i >= 0; i--) {  
        flip.add(orig.get(i));  
    }  
}
```

2. What is the *best case* runtime for rev in Big-O notation? You can assume that n is the length of orig.

Note: write an upper bound using O as in “ $O(g(n))$ ”

Your answer: _____

Does the **best case** depend on n ? (where n is the size of the input)

```
public static <E> void rev(Vector<E> orig, Vector<E> flip) {  
    for (int i = orig.size() - 1; i >= 0; i--) {  
        flip.add(orig.get(i));  
    }  
}
```

Yes. Don't be fooled by the fact that n can be 0.

```
public static int findFirstSpace(String s) {  
    for(int i = 0; i < s.length(); i++) {  
        if(s.charAt(i) == ' ') {  
            System.out.println("Found space at i = " + i);  
            return i;  
        }  
    }  
    return -1;  
}
```

No.

The algorithm can finish early **even if n is large.**

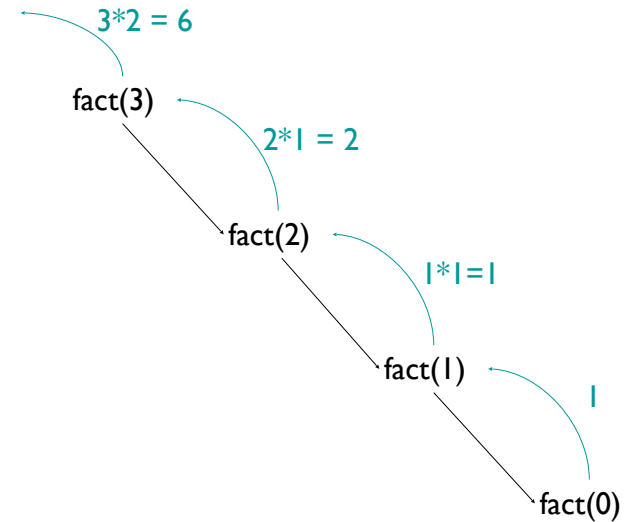
Practice Quiz

Recall: Factorial

- $n! = n \times (n-1) \times (n-2) \times \dots \times 1$
- Work with a partner and see if you can come up with a recursive solution.

How much does a **recursive** solution **cost**?

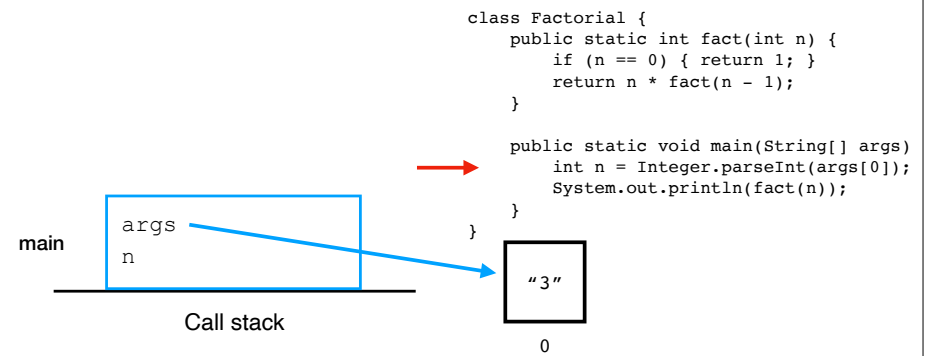
Graphically...



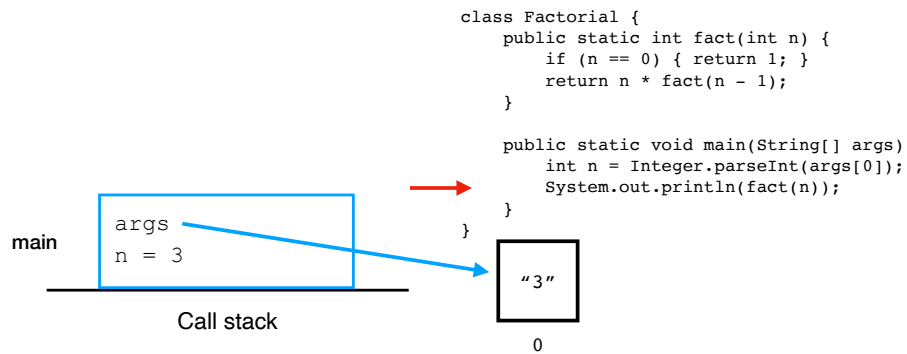
```
class Factorial {  
    public static int fact(int n) {  
        if (n == 0) { return 1; }  
        return n * fact(n - 1);  
    }  
    → public static void main(String[] args)  
        int n = Integer.parseInt(args[0]);  
        System.out.println(fact(n));  
    }  
}
```

Call stack

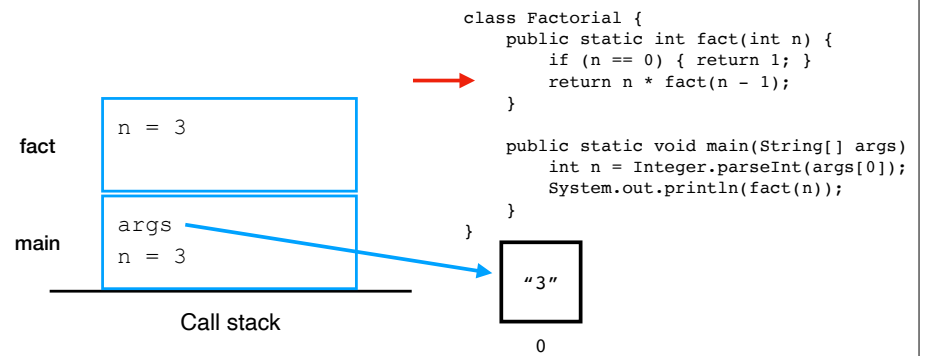
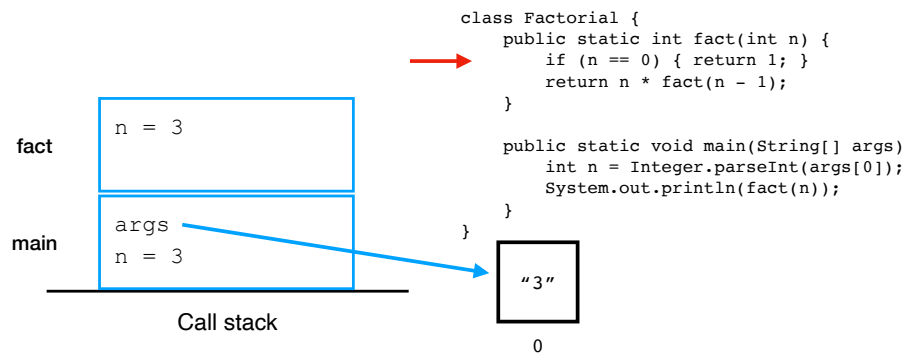
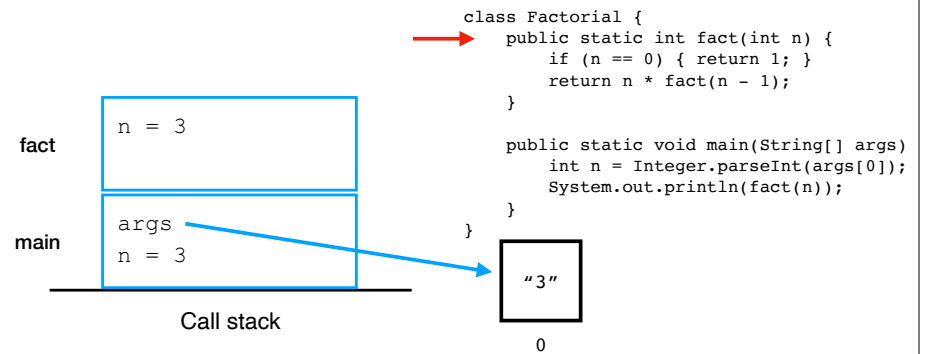
Call program with input "3".

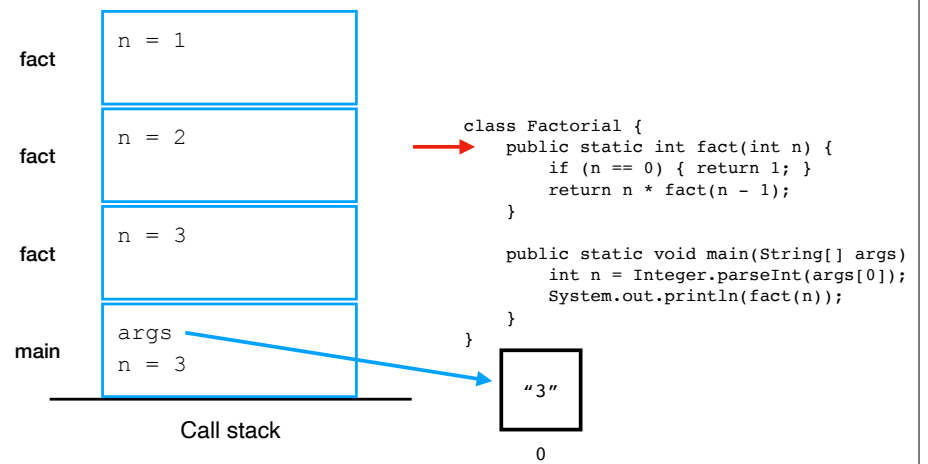
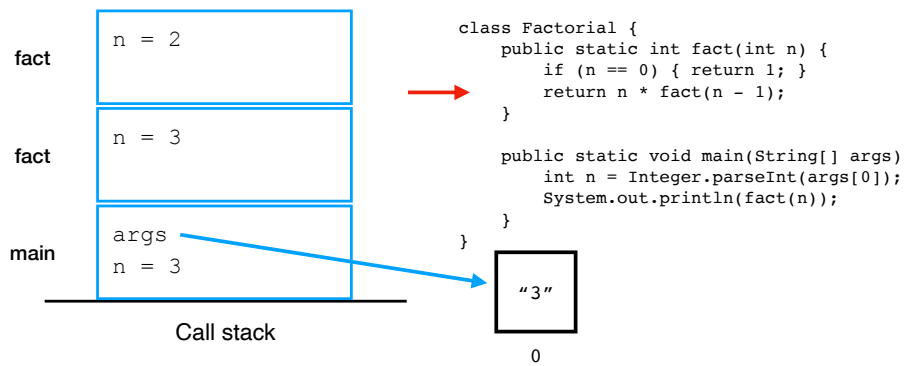
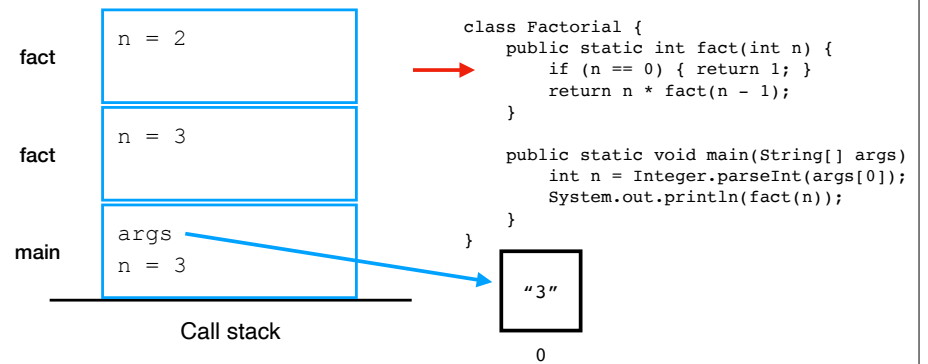
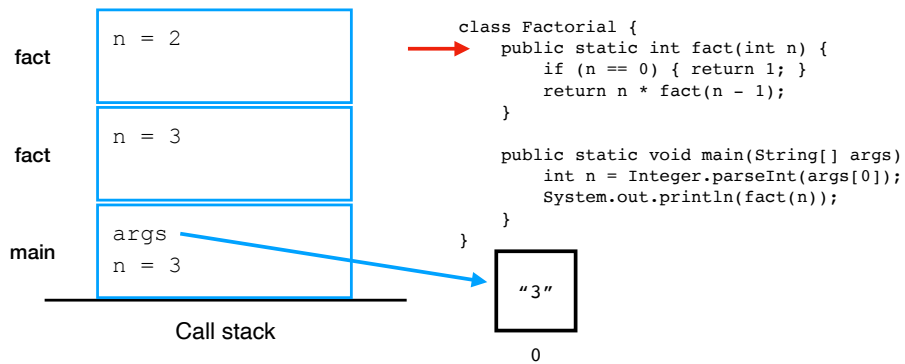


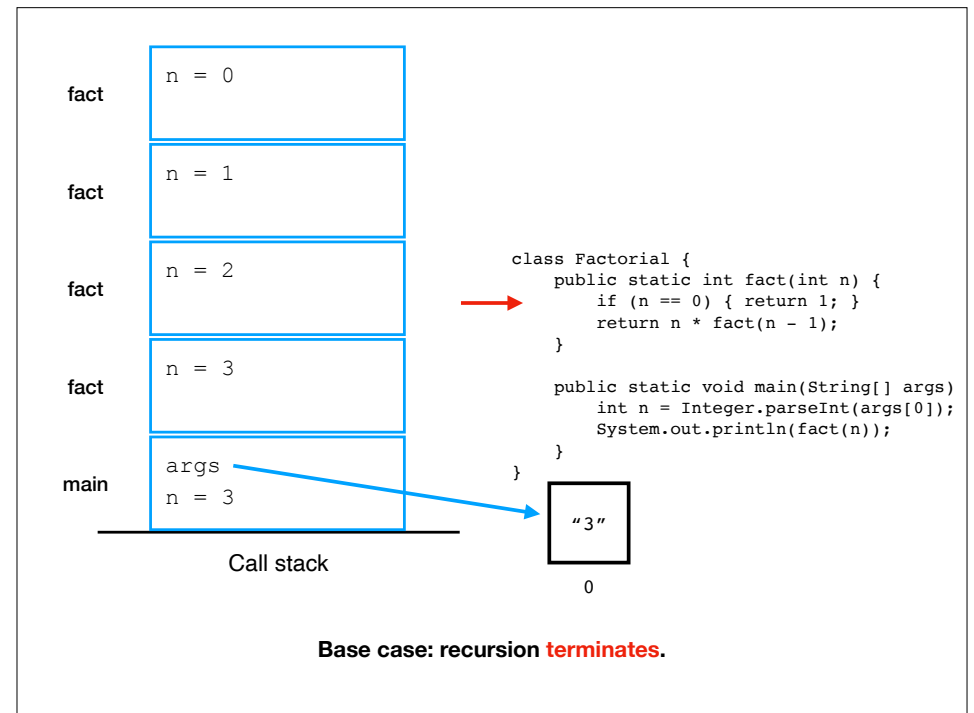
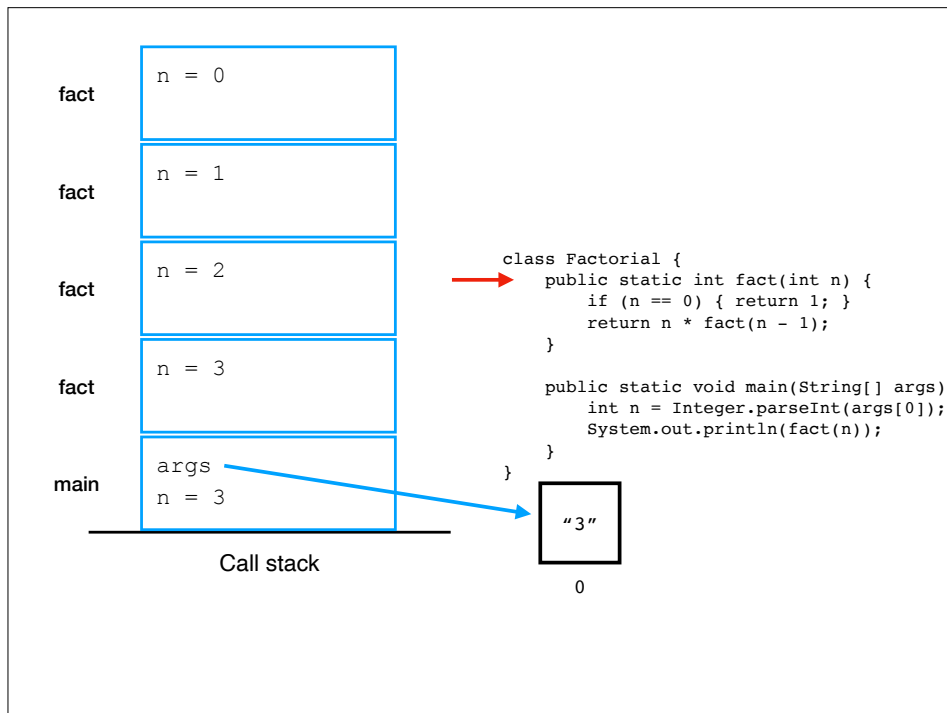
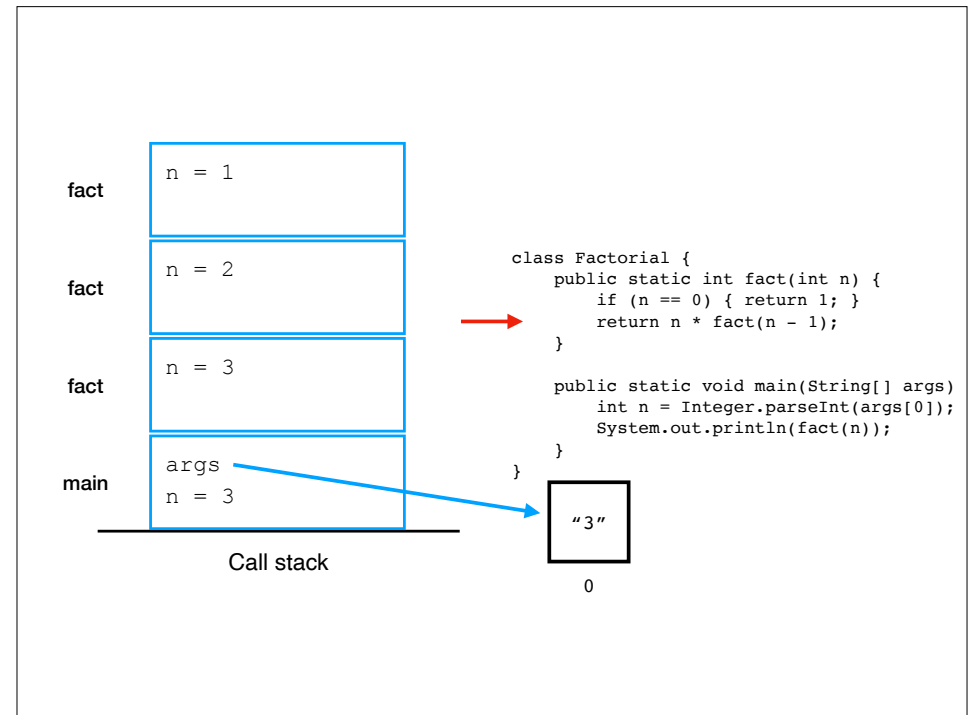
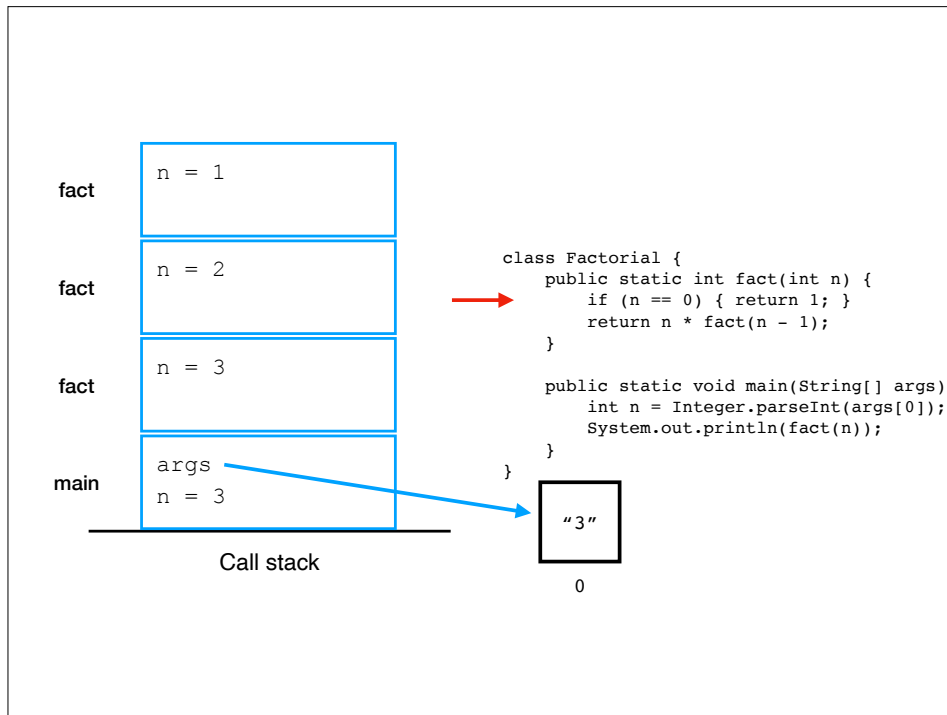
Call program with input "3".

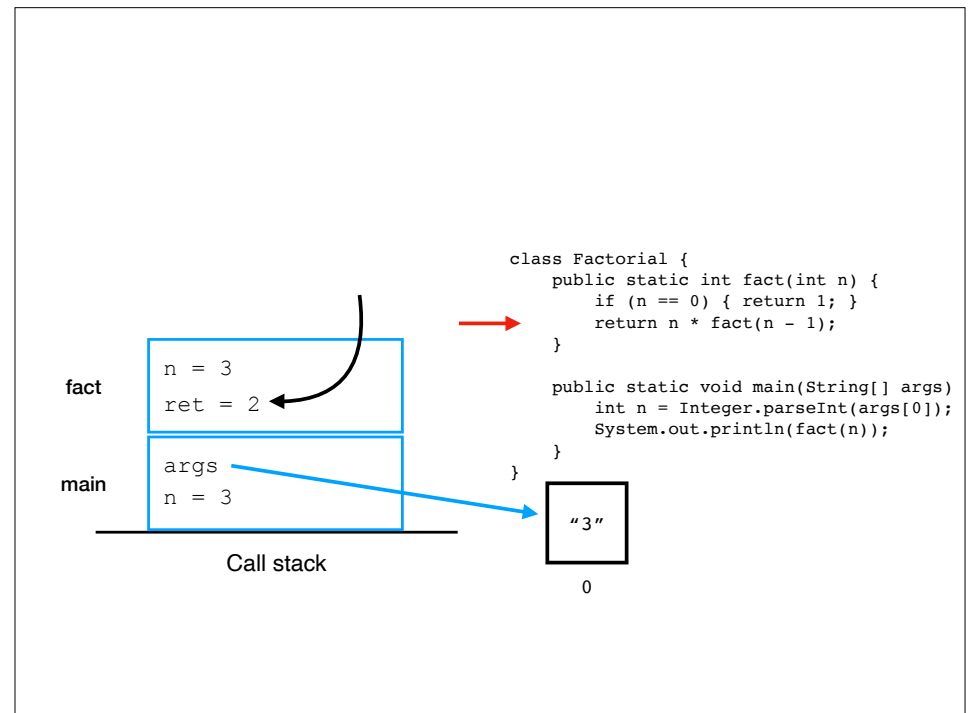
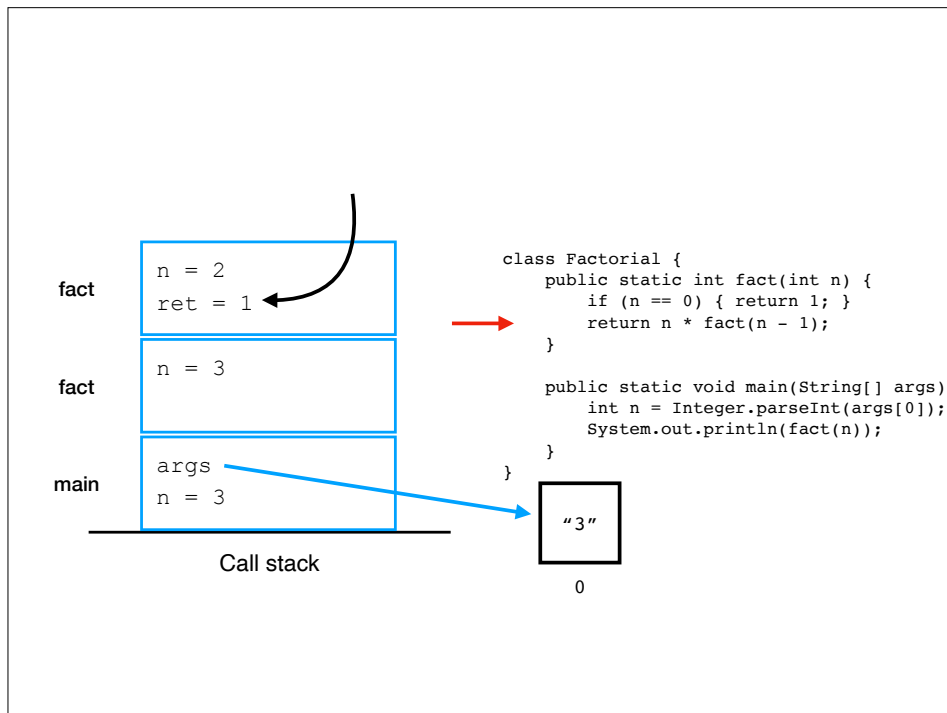
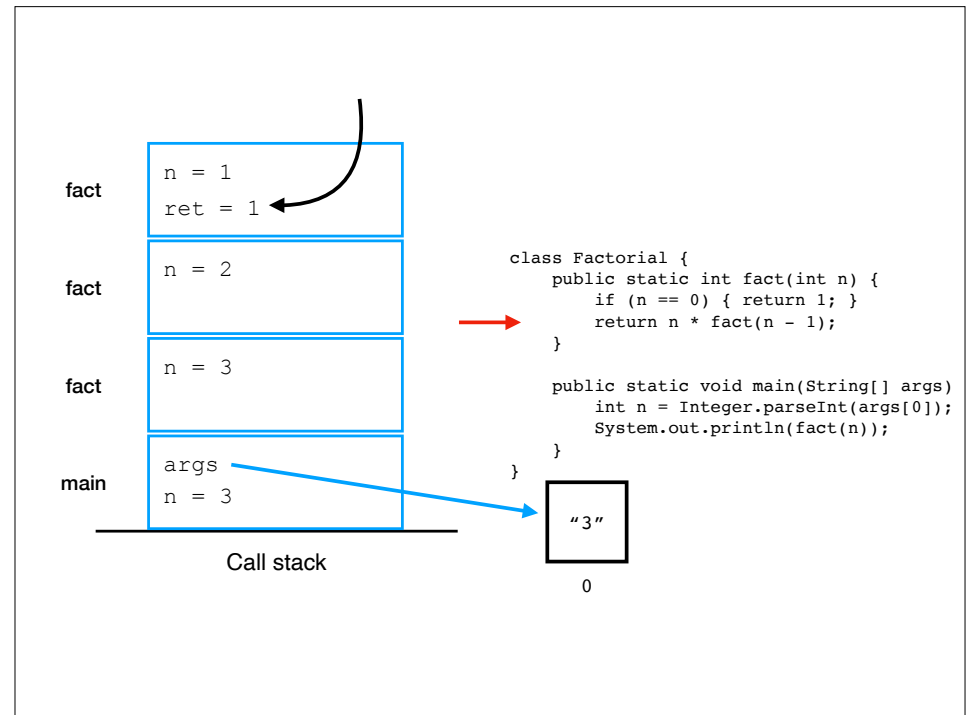
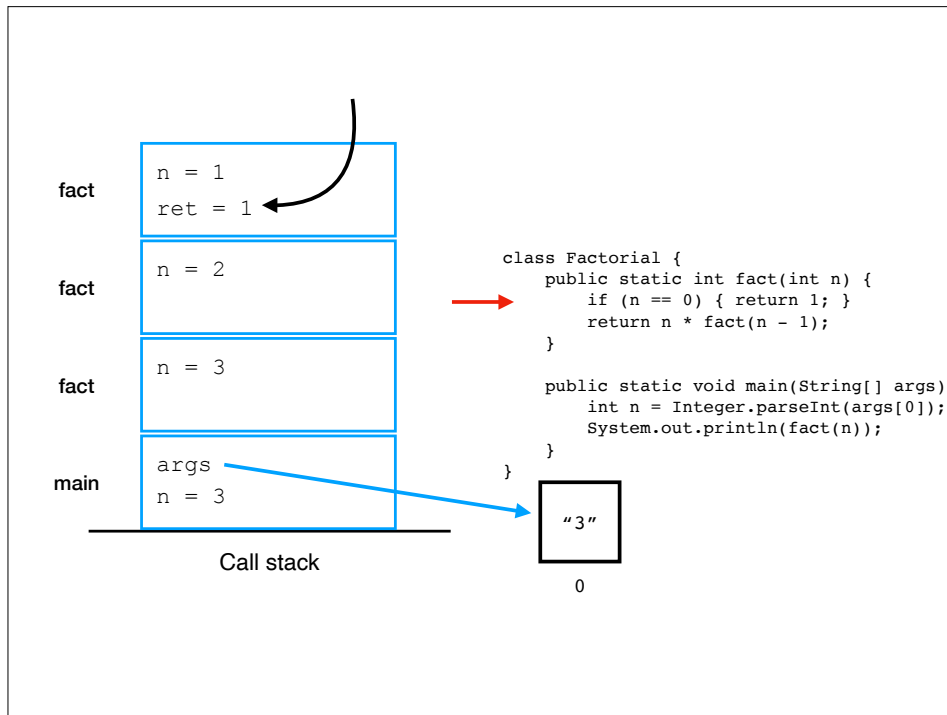


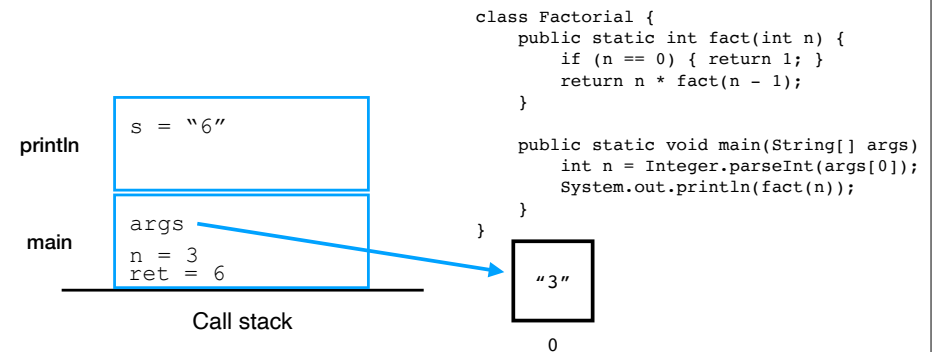
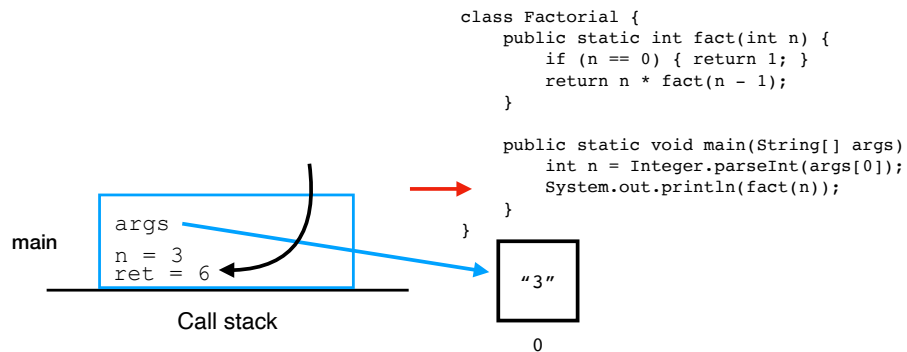
I skipped a subtlety here; did you spot it?



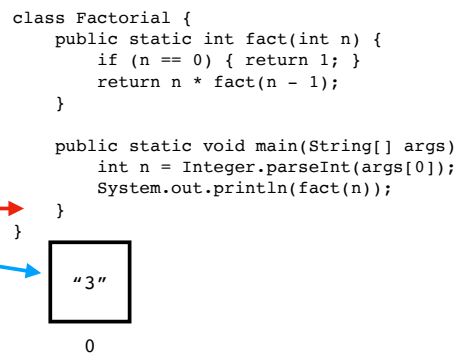








I skipped another subtlety here; did you spot it?



Recursion tradeoffs

- Advantages
 - Often **easier** to construct recursive solution
 - Code is usually **cleaner**
 - Some problems do not have **obvious** non-recursive solutions
- Disadvantages
 - Time cost** of recursive calls
 - Memory cost** (need to store state for each recursive call until base case is reached)

Mathematical Induction



A note about “formal methods”



If the problem “fits” the mold, there is a procedure for determining truth.

Mathematical Induction

- The **mathematical cousin** of **recursion** is **induction**
- Induction is a **proof technique**
- Purpose: to **simultaneously prove** an **infinite number** of theorems!

Principle of Mathematical Induction

Let $P(n)$ be a **predicate** that is defined for **integers** n , and let a be a **fixed integer**.

If the following two statements are **true**:

1. $P(a)$ is **true**.
2. For all integers $k \geq a$, **if** $P(k)$ **is true then** $P(k + 1)$ **is true**.

then the statement

for all integers $n \geq a$, $P(n)$ **is true**

is **also true**.

Principle of Mathematical Induction (variant)

Let $P(n)$ be a **predicate** that is defined for **integers** n , and let a be a **fixed integer**.

If the following two statements are **true**:

1. $P(a)$ is **true**.
2. For all integers $k > a$, if $P(k-1)$ is **true** then $P(k)$ is **true**.

then the statement

for all integers $n \geq a$, $P(n)$ is **true**

is **also true**.

To be clear:

If you want to prove that $P(n)$ is **true** for all integers $n \geq a$,

1. You must first prove that $P(a)$ is **true**.

2. Then you must prove that:

For all integers $k \geq a$, if $P(k)$ is **true** then $P(k+1)$ is **true**.

Critically, when proving #2, **assume** that $P(k)$ is **true** and **show** that $P(k+1)$ **must also be true**.

Names for things and “form”

Hypothesis: $P(n)$ is **true** for all integers $n \geq a$,

1. Base case: $P(a)$ is **true**.

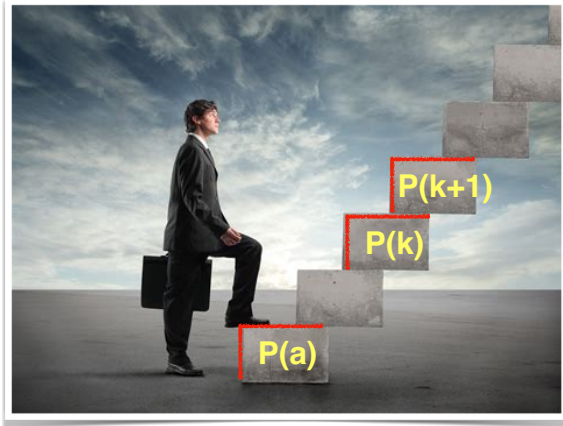
2. Inductive step:

For all integers $k \geq a$, if $P(k)$ is **true** then $P(k+1)$ is **true**.

Like recursion, there is an analogy



Like recursion, there is an analogy



Example

Prove that the sum of the first n integers is:

$$\frac{n(n+1)}{2}$$

Example

Put another way, prove

$$P(n) : 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

for all $n \geq 1$.

We have an unbounded number of hypotheses ("for all $n \geq 1$ ").

Use **mathematical induction**.

Remember the template!

Step 1: Prove **$P(a)$**

Step 2: Prove **$P(k) \Rightarrow P(k+1)$**

Therefore,

$$P(n) : 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

For all $n \geq 1$.

Is **true**.

Example

Step 1: Prove **P(a)**

What would a good **a** be?

$$P(n) : 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

The “simplest” instance is **a = 1**. Let’s start there.

Example

Step 1: Prove **P(a)**

$$P(a) : 1 = \frac{1(1+1)}{2}$$

Is this statement true? **Yes.**

$$\text{Proof: } \frac{1(1+1)}{2} = \frac{2}{2} = 1$$

Example

Step 2: Prove **P(k) \Rightarrow P(k+1)**

Assume the following is true:

$$\mathbf{P(k)} : 1 + 2 + 3 + \dots + k = \frac{k(k+1)}{2}$$

Prove:

$$\mathbf{P(k+1)} : 1 + 2 + 3 + \dots + (k + 1) = \frac{(k+1)((k+1)+1)}{2}$$

Example

Step 2: Prove **P(k) \Rightarrow P(k+1)**

$$\mathbf{P(k+1)} : 1 + 2 + 3 + \dots + (k + 1) = \frac{(k+1)((k+1)+1)}{2}$$

Let’s handle the left side first.

$$1 + 2 + 3 + \dots + (k + 1)$$

Looks familiar. Isn’t it the same as:

$$(1 + 2 + 3 + \dots + k) + (k + 1)$$

Example

Step 2: Prove $P(k) \Rightarrow P(k+1)$

$$(1 + 2 + 3 + \dots + k) + (k + 1)$$

According to $P(k)$, which is true,
it must be equal to:

$$(1 + 2 + 3 + \dots + k) + (k + 1) = \frac{k(k+1)}{2} + (k + 1)$$

Example

Step 2: Prove $P(k) \Rightarrow P(k+1)$

$$\text{Simplify} \quad = \frac{k(k+1)}{2} + (k + 1)$$

$$= \frac{k(k+1)}{2} + \frac{2(k+1)}{2}$$

$$= \frac{k(k+1) + 2(k+1)}{2}$$

Let's stop here.
The left side is

$$= \frac{(k+1)(k+2)}{2}$$

Example

Step 2: Prove $P(k) \Rightarrow P(k+1)$

$$P(k+1) : 1 + 2 + 3 + \dots + (k + 1) = \frac{(k+1)((k+1)+1)}{2}$$

Let's handle the right side now.

$$\frac{(k+1)((k+1)+1)}{2}$$

Simplify

$$\frac{(k+1)(k+2)}{2}$$

Let's stop here.

Example

Step 2: Prove $P(k) \Rightarrow P(k+1)$

$$P(k+1) : 1 + 2 + 3 + \dots + (k + 1) = \frac{(k+1)((k+1)+1)}{2}$$

We just showed that the left side

$$\frac{(k+1)(k+2)}{2}$$

equals the right side

$$\frac{(k+1)(k+2)}{2}$$

Example

Step 1: Prove $P(a)$ ✓

Step 2: Prove $P(k) \Rightarrow P(k+1)$ ✓

Therefore,

$$P(n) : 1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

For all $n \geq 1$.

Is true. ✓

Recap & Next Class

Today:

- Recursion costs
- Mathematical induction

Next class:

- ADTs
- Lists