CSCI 136
Data Structures &
Advanced Programming

Lecture 7

Spring 2020

Bill and Dan

Administrative Details

- Lab 2 Due Today
  - Anyone Stuck?
- Lab 3 Wednesday
  - Partner Lab
    - One repository where both people have access
    - Beware of merge conflicts!
  - There will be “warm-up” problems, not a design doc
    - We’ll go over warm-up questions at the start of lab, but thinking about them will really get you to practice “thinking recursively”

Last Time

- Measuring Growth
  - Rough discussion of Big-O w.r.t. Vectors
    - We care about trends
    - Goal: determine how performance scales with input size.
    - We often care most about the worst case behavior, but we can analyze best, worst, and average cases

Today

- Revisit `Vector` growing examples
- Recursion
- Mathematical Induction

Vector Operations : Worst-Case

Let `n` = Vector size (*not* capacity!):

- O(1) operations (cost is same regardless of size):
  - `size()`, `capacity()`, `isEmpty()`, `get(i)`, `set(i)`, `firstElement()`, `lastElement()`
- O(n) operations (cost grows proportionally to size):
  - `indexOf()`, `contains()`, `remove(elt)`, `remove(i)`
- What about add methods?
  - If Vector doesn't need to grow
    - `add(elt)` is `O(1)` but `add(elt, i)` is `O(n)`
  - Otherwise, depends on `ensureCapacity()` time
    - Time to copy array: `O(n)`

Vectors: Add Method Complexity

Suppose we grow the Vector's array by a fixed amount `d`.

How long does it take to add `n` items to an empty Vector?

- The array will be copied each time its capacity needs to exceed a multiple of d
  - At sizes `0, d, 2d, … , n/d.`
- Copying an array of size `kd` takes `ckd` steps for some constant `c`, giving a total of

$$\sum_{k=1}^{n/d} ckd \quad = cd \sum_{k=1}^{n/d} k \;= cd \left(\frac{n}{d}\right)\left(\frac{n}{d}+1\right)/2 \;= O(n^2)$$

Vectors: Add Method Complexity

Suppose we grow the Vector's array by doubling.

How long does it take to add `n` items to an empty Vector?

- The array will be copied each time its capacity needs to exceed a power of 2
  - At sizes `0, 1, 2, 4, 8 …, n/2`
- The total number of elements are copied when `n` elements are added is:
  - `1 + 2 + 4 + ... + n/2 = n-1 = O(n)`

- Very cool! (So cool that we'll prove it later using induction!)

Common Complexities

For `n` = measure of problem size:

- $O(1)$: constant time and space (same cost regardless of n)
- $O(\log n)$: divide and conquer algorithms, binary search
- $O(n)$: linear scan (e.g., list lookup)
- $O(n \log n)$: divide and conquer sorting algorithms
- $O(n^2)$: matrix addition, selection sort
- $O(n^3)$: matrix multiplication
- $O(n^k)$: cell phone switching algorithms
- $O(2^n)$: subset sum, graph 3-coloring, satisfiability, ...
- $O(n!)$: traveling salesman problem (in fact $O(n^2 2^n)$)

Recursion

- General problem solving strategy
  - Break problem into sub-problems of same type
  - Solve sub-problems
  - Combine sub-problem solutions into solution for original problem
    - Recursive leap of faith!

Recursion

- Many algorithms are recursive
  - Can be easier to understand (and prove correctness/state efficiency of) than iterative versions
  - They feel *elegant*
- Today we will review recursion and then talk about techniques for reasoning about recursive algorithms

Think Recursively

- In recursion, we always use the same basic approach
- What's our base case? [Sometimes "cases"]
  - n=0? list.isEmpty()?
- What' s the recursive relationship?
  - How can we use the solution to a smaller version of the problem to answer the question?

In-person Demo

- How much money do I have in my Jar?
  - Noah says: “I have too much money to count”
  - Alec says: “That is just not true”
  - Noah says: “It’s just impossible for anyone to count that high”
  - Alec says: “I’m on vacation, and I’ve got nothing better to do. You’re sitting here and we’re counting this money”

  - Noah had $37.22, a few rocks, and an enamel pin

In-person Demo

- How much money do I have in my Jar?
  - Suppose I know the value of any coin, and I can add two numbers.
  - What is the base case? (What is the simplest jar that I can look at and know how much money it contains?)
  - What is the recursive step? (How might I take a full jar, and decompose it into smaller cases that I know how to handle?)

Practice Writing Code: Factorial

Definition

- $n! = n \cdot (n-1) \cdot (n-2) \cdot \ldots \cdot 1$
- How can we implement this?
  - We could use a for loop…
- But we could also write it recursively
  - $n! = n \cdot (n-1)!$
  - $0! = 1$
- With a partner, try to write fact()

Practice Writing Code: Factorial

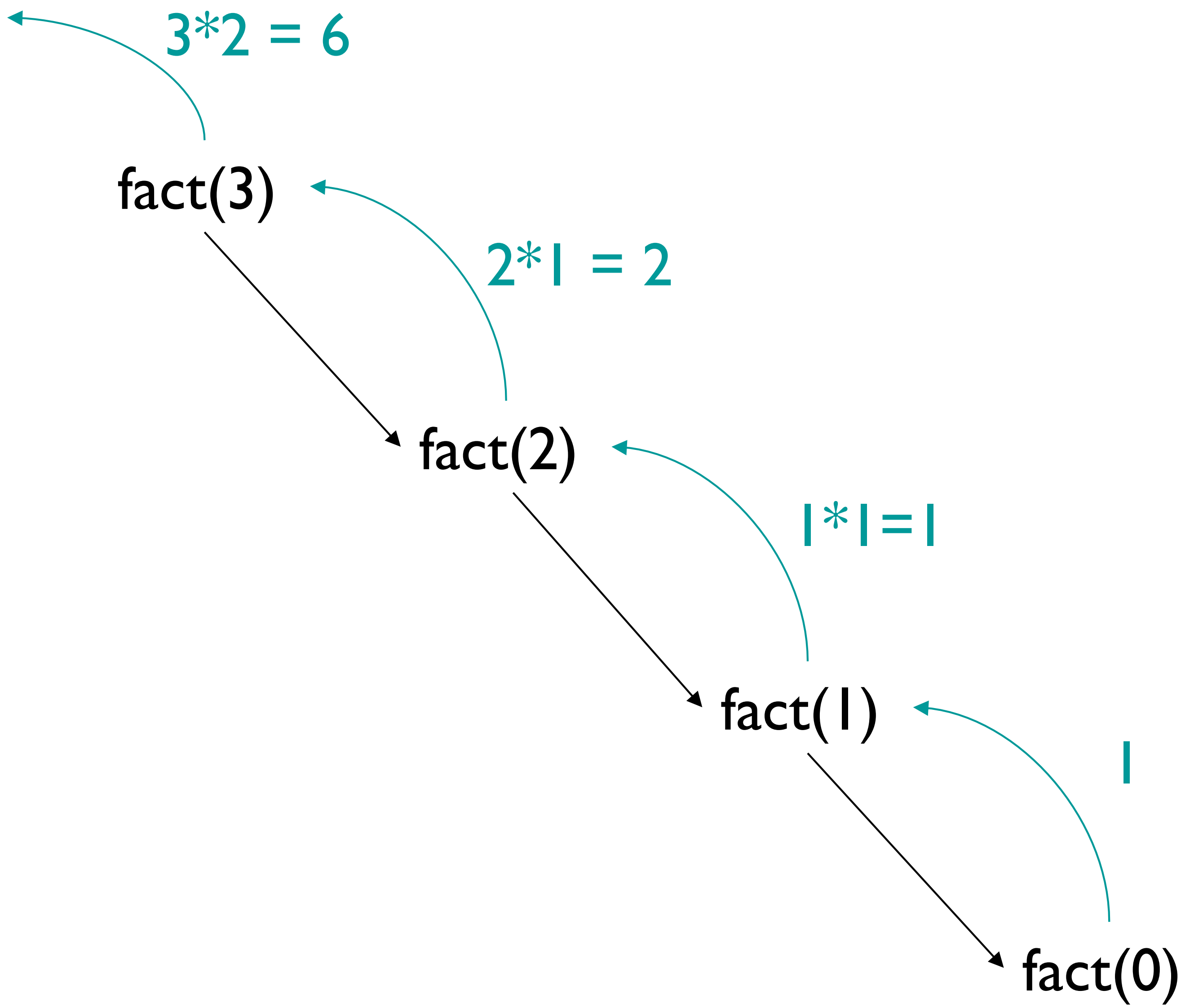
Definition

- $n! = n \cdot (n-1)!$
- $0! = 1$

```
public static ___ fact(_______) {
      // base case?

      // recursive case?
}
```

Factorial



Fact.java

```
public class Fact{

    // Pre: n >= 0
    public static int fact(int n) {
        // base case
        if (n==0) {
            return 1;
        }
        // recursive leap of faith
        else {
            return n*fact(n-1);
        }
    }

    public static void main(String args[]) {
        System.out.println(fact(Integer.valueOf(args[0]).intValue()));
    }

}
```

Recursion Tradeoffs

- Advantages
  - Often easier to construct recursive solution
  - Code is usually cleaner (so *elegant*!)
  - Some problems do not have obvious non-recursive solutions
- Disadvantages
  - Overhead of recursive calls
  - Can use lots of memory (need to store state for each recursive call until base case is reached)
    - E.g. recursive fibonacci method