

CSCI 136:
Data Structures
and
Advanced Programming

Lecture 32

Hash tables, part 4

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Announcements

Lab 10 Part 1 extended: due with Part 2

Part 2: MTF hashtable removed

All graded labs back this week (exc. 10)

Outline

Hashtable recap

Hashtables: big picture

We use **hash tables** when our problem meets **two criteria**:

1. A **table-like structure** is convenient (i.e., a `Map<K, V>`).
2. We are willing to trade an **ordering** of elements for **high performance** (**$O(1)$** put/get), on average.

Hashtables: big picture

What we **want**:

key ↗ ↘ value

lunchMenu =

Monday	nuggets
Tuesday	mac n chez
Wednesday	pizza
Thursday	burgers
Friday	fish stix
Saturday	taco salad
Sunday	soup

Intuitively:
You can find what
you're looking for
quickly.

More generally, a **table ADT** is a `Map<K, V>`.

Hashtables: big picture

In a table, the **key** must be **unique**.

key ↗ ↘ value

lunchMenu =

Monday	nuggets
Tuesday	mac n chez
Wednesday	pizza
Thursday	burgers
Friday	fish stix
Saturday	taco salad
Sunday	soup

Hashtables: big picture

Also, the **order of entries** are not particularly important.

key ↗ ↘ value

lunchMenu =

Friday	fish stix
Tuesday	mac n chez
Sunday	soup
Monday	nuggets
Wednesday	pizza
Thursday	burgers
Saturday	taco salad

Hashtables: big picture

We always build complex data structures **from more primitive parts**.

What we **have**:

0	1	2	3	4	5	6	7

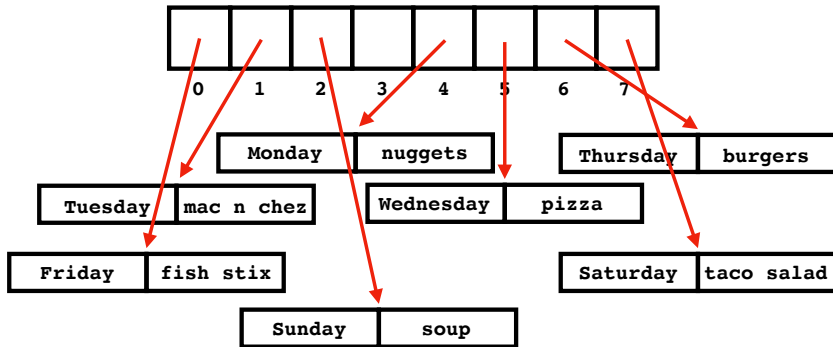
An **array**. **Not a table!**

Monday	nuggets
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Also, **associations**. Also **not a table!**
But with one more ingredient: **a table!**

Hashtables: big picture

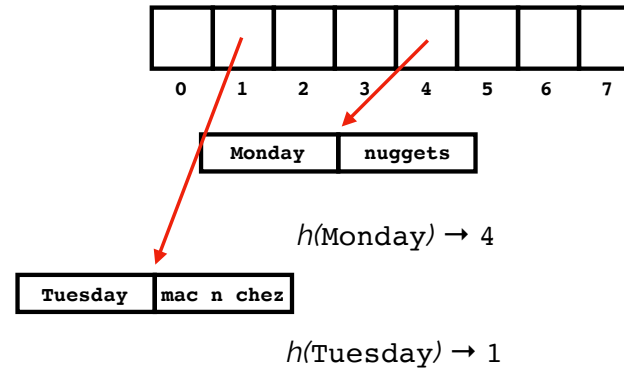
Let's **build a table** from our **parts**.



Close. But why did we **insert** in those **buckets**?

Hashtables: big picture

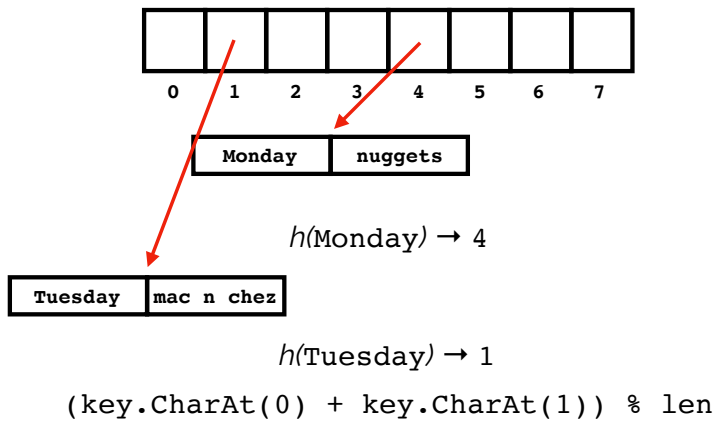
Last ingredient: a **function** that maps **keys** to array **indices**.



Hashtables: big picture

How might we obtain such a function?

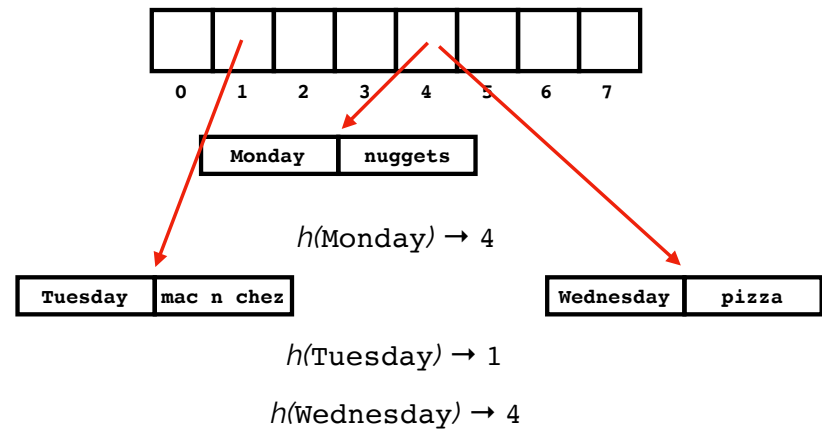
Any function from $T \rightarrow \text{int}$ will do.



$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

Hashtables: big picture

Some functions are **better** than others.

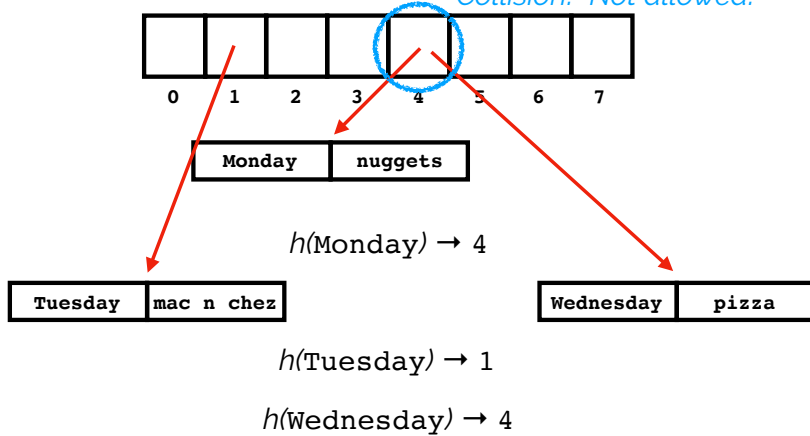


$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

Hashtables: big picture

Some functions are **better** than others.

"Collision!" Not allowed!

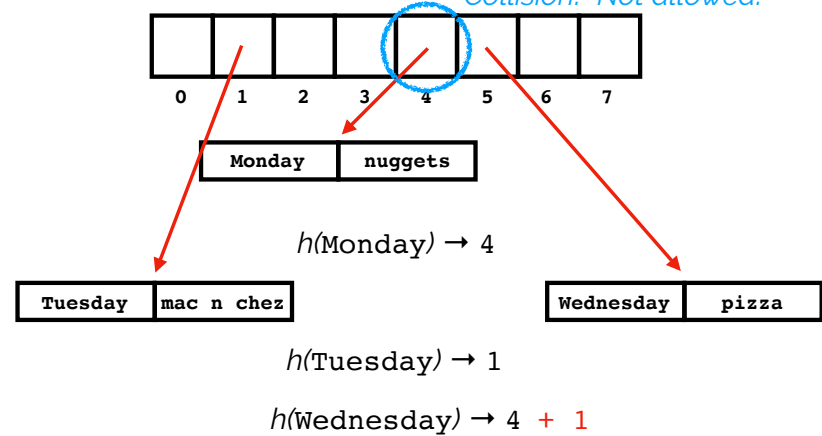


$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

Hashtables: big picture

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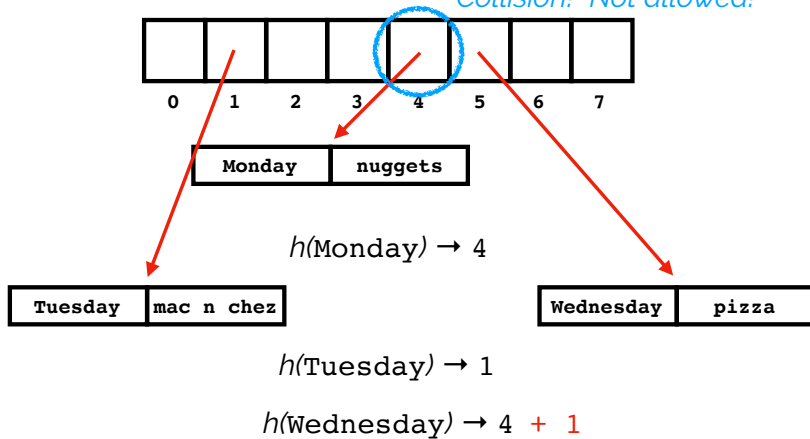


$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

Hashtables: big picture

When hashed keys **collide**, we need a **recovery method**.

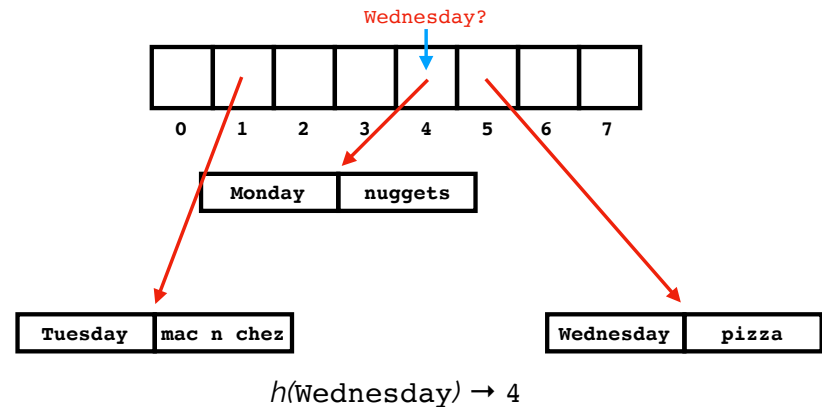
"Collision!" Not allowed!



$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

Hashtables: big picture

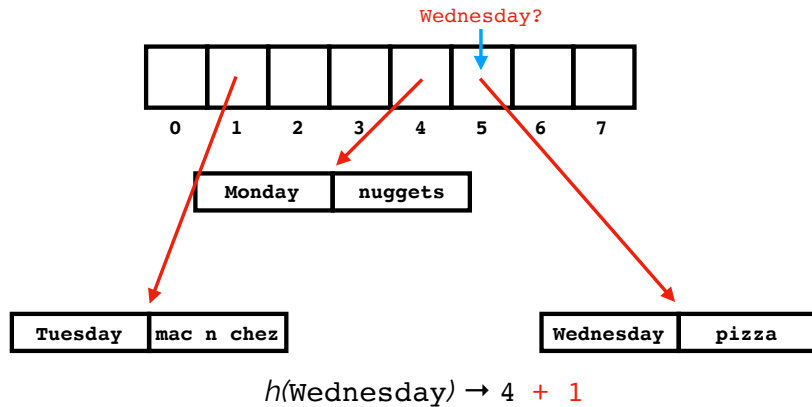
This turns out to be important on **lookup** too.



$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

Hashtables: big picture

This turns out to be important on **lookup** too.



Hashtable recipe

Ingredients:

1. Array (or array-like structure).
2. Association of key and value.
3. Hash function.
4. Collision recovery method.

Insertion (or lookup) procedure:

1. Hash the key to obtain index.
2. If index is occupied (not what we're looking for), find new index using recovery method.
3. Insert (or return) key.

Hashtable expansion

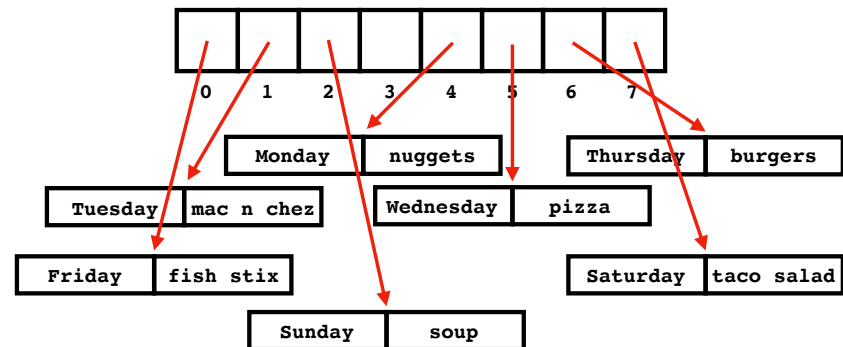
A **Vector** is a convenient way to store key-value Associations in a hash table because, when the hash table fills up, a **Vector** can expand.

But: on expansion, all keys need to be **hashed again and reinserted**. Why? Because hash value depends on **Vector length**.

$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$

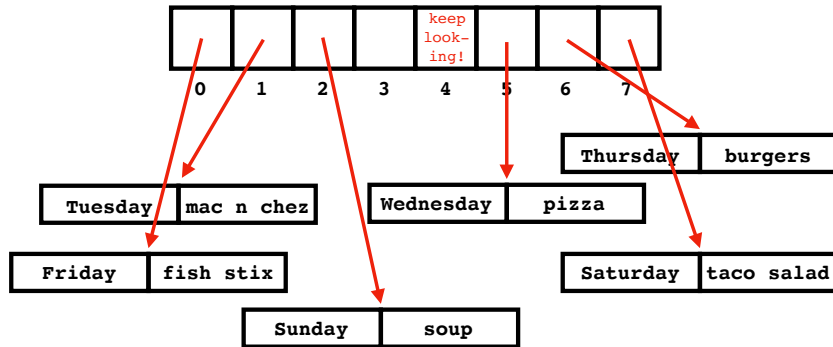
Hashtable deletion

Also: **deletion** of a key may require leaving a **placeholder** behind.



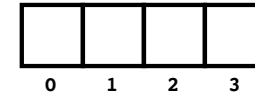
Hashtable deletion

Also: **deletion** of a key may require leaving a **placeholder** behind.



`HashAssociation` has a `reserved` flag for this purpose.

Hashtable activity



$$(\text{key.CharAt}(0) + \text{key.CharAt}(1)) \% \text{len}$$

Harry Potter, Gryffindor
Draco Malfoy, Slytherin
Luna Lovegood, Ravenclaw
Gilderoy Lockhart, Ravenclaw
Cedric Diggory, Hufflepuff
Hermione Grainger, Gryffindor
Dolores Umbridge, Slytherin

Recap & Next Class

Today we learned:

Hashtable recap

Next class:

Shortest paths