

CSCI 136:
Data Structures
and
Advanced Programming
Lecture 12
Asymptotic analysis, part 3

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Williams

Outline

Study tip

Proof: doubling is good strategy

Interfaces

Inheritance

Announcements

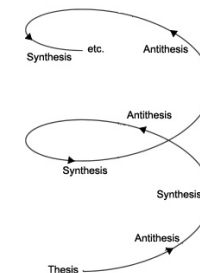
Feedback: How should I study for midterm?

Grades will be determined as follows:

Final exam:	20%
Midterm exam:	20%
Programs/Labs:	35%
Code reviews:	10%
Engagement:	15%

Life skill #9

Experimentation



1. A **thesis** is an intellectual proposition (i.e., a **T/F** statement).
2. An **antithesis** is an alternative proposition using the same facts.
3. **Synthesis** reconciles the two hypotheses by gathering new facts and proposing a new thesis.

Life skill #9

Thesis: **You tell me**. How should you study?

Life skill #9

Antithesis:

1. Revisit your **glossary** of terms.
2. Revisit **homework**. What did you get wrong?
3. What **activities** have we done in class?
 1. Expect to have to **code on paper**.
 2. Expect to have to **prove inductively**.
4. **Practice problems** in book (you have solutions!)

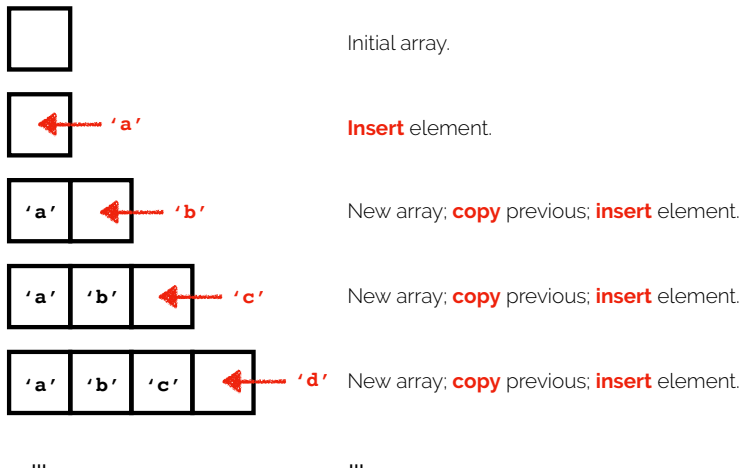
Life skill #9

Synthesis: How should you study?

From last class

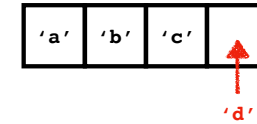
Why is the **array doubling** strategy for Vector **better** than expanding the array **one element at a time**?

One-at-a-time expansion



Insertion into an array

How much does **array insertion** cost?



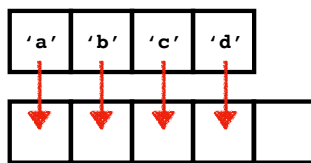
It costs **$O(1)$** .

In fact, lookup and insertion both cost **$O(1)$** .

Tradeoff: arrays are fixed size.

Copying an array

How much does an **array copy** cost?

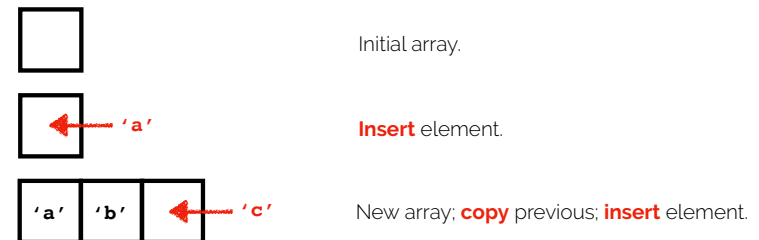


It costs **$O(1) \times m$** , where **m** is the size of the original array.

$\approx O(m)$

One-at-a-time expansion costs?

(in the worst case, each time)



$O(m) + O(1) \approx O(m)$, where **m** is the size of the original array.

Cost is **dominated by the size of the array** being copied.

How many copies?

of copies for one-at-a-time expansion:

$$\begin{array}{ccccccc} & \mathbf{1} & + & \mathbf{2} & + & \mathbf{3} & + \dots + \mathbf{(n-1)} \\ \text{add()} & \text{2nd} & & \text{3rd} & & \text{4th} & & \text{nth} \\ & \text{elem.} & & \text{elem.} & & \text{elem.} & & \text{elem.} \end{array}$$

Recall theorem: $1 + 2 + 3 + \dots + k = k(k+1)/2$

$$\begin{aligned} \text{Sub } n-1 \text{ for } k: & (n-1)((n-1)+1)/2 = n(n-1)/2 \\ & = n^2/2 - n/2 \end{aligned}$$

One-at-a-time expansion costs $\approx O(n^2)$

How many copies?

of copies for doubling expansion:

$$\begin{array}{ccccccc} & \mathbf{1} & + & \mathbf{2} & + & \mathbf{4} & + \dots + \mathbf{(n/2)} \\ \text{add()} & \text{up to} & & \text{up to} & & \text{up to} & & \text{up to} \\ & \text{2nd} & & \text{4th} & & \text{8th} & & \text{nth} \\ & \text{elem.} & & \text{elem.} & & \text{elem.} & & \text{elem.} \end{array}$$

Neat theorem: $1 + 2 + 4 + \dots + 2^{k-1} = 2^k - 1$

Suppose $n = 2^k$.

$$\begin{aligned} \text{Then } 1 + \dots + n/2 &= 1 + \dots + 2^{k/2} \\ &= 1 + \dots + 2^{k-1} = 2^k - 1 = n - 1 \end{aligned}$$

Doubling expansion costs $\approx O(n)$

Which is faster?



🤩 One-at-a-time expansion costs $\approx O(n^2)$ 🤩

😎 Doubling expansion costs $\approx O(n)$ 😎

Doubling is Vin Diesel-approved.

Interfaces

Interface

An **interface** defines boundary between two systems across which they share information. An interface is a **contract**: calling a method defined in an interface returns the data as promised.

A key principle of object-oriented design is to **deny access** to all data (i.e., to make `private`) by default, allowing access only through methods specified by the interface.

(code)

Inheritance

Inheritance

Inheritance is a **mechanism** for defining a class in terms of another class. It is a labor-saving device employed to reduce **code duplication**. Inheritance allows programmers to specify a new implementation while :

1. **maintaining the same behavior**,
2. **reusing code**, and
3. **extending the functionality** of existing software.

(code)

Recap & Next Class

Today we learned:

Sample Big-O analysis

Interfaces

Inheritance

Next class:

Sorting