

# Lec 8: Doubly Linked Lists

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# Admin

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- In-person lab today
  - Starter code is: `ArrayList.java`, `NodeInt.java`, `LinkedListInt.java` plus files from today
  
- Any questions?

## **Linked List of Ints**

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## Int Linked List Nodes (Review)

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- In our example, students held the information of the list
- We'll store the information in Nodes
  - We'll call them `NodeInt` to clarify that these nodes only help us to store lists of `ints`
- What should the `NodeInt` class look like? Let's code it up together.

# Int Linked List Class

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# Int Linked List Class

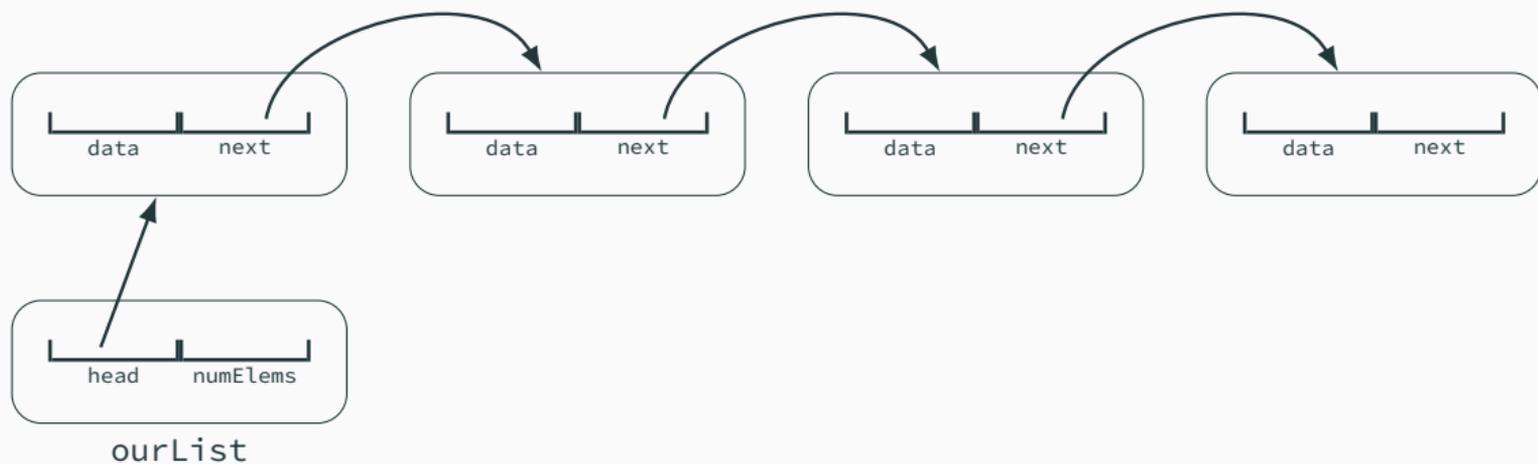
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- The list object only needs to store a reference to the first node!
- We'll also store the number of elements in the list again, like we did for `IntArrayList` (will be handy)
- Maintain the following **invariant**: the `next` pointer of a node is `null` only when the node is the last in the list

# Overview of Plan

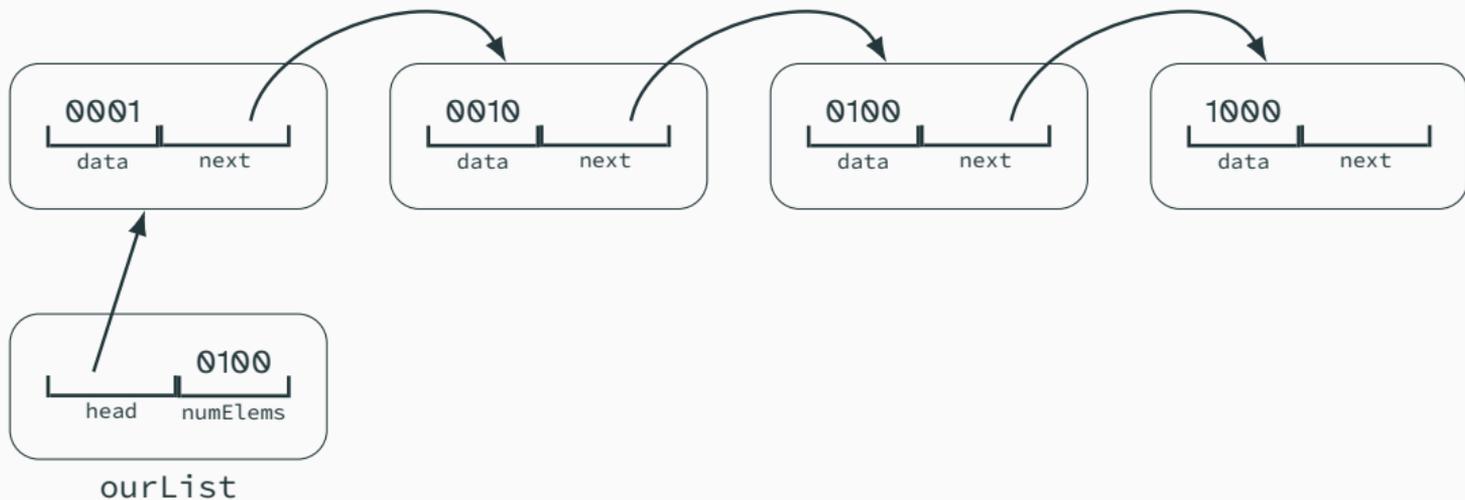
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## Overview of Plan: With Data

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The following list stores: [1, 2, 3, 4]. ints are truncated to 4 bits for readability.



## get and set

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- `int get(int index)` returns element `index` stored in the linked list. We'll start by accessing the Node stored in `head`; its value is correct for index `0`. Then, we'll go to the next element of that Node, then its next, etc., `index` times. Don't forget to do bounds checking!

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- `int get(int index)` returns element `index` stored in the linked list. We'll start by accessing the `Node` stored in `head`; its value is correct for index `0`. Then, we'll go to the next element of that `Node`, then its next, etc., `index` times. Don't forget to do bounds checking!
- `int set(int index, int newElement)` sets the element stored at `index` to be `newElement`, and returns the element *previously* stored in slot `index`. We'll accomplish this similarly: scan through the list until the `i`th node, but now we set its value rather than getting it.

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- Both methods need to scan through `index` nodes. Let's set up a method to do that!

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(Fun video)

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- Finally, we need to “join” our new node to the end of the list: the next value of the node at the end of the list should be the new node we made.

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- Finally, we need to “join” our new node to the end of the list: the next value of the node at the end of the list should be the new node we made.
- Let's do this in the code

## Adding at an index

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- `void add(int index, int newElement)`: adds `newElement` to index `index` in the list, pushing all later elements in the list down by one. In a linked list, we just place the new node in the correct location. We do not need to manually “push down” elements: all later elements will automatically be one further down in the list, since there is now a new node in front of them.

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- This method requires some careful “surgery” on the linked list: we need to disconnect the reference to insert the new node. Let’s draw a picture of what that looks like. Let’s say we have found the  $(index - 1)$ -st node in our list; its next node is the current `index`-th node. We want our new node to go in between them.

## Adding at an index

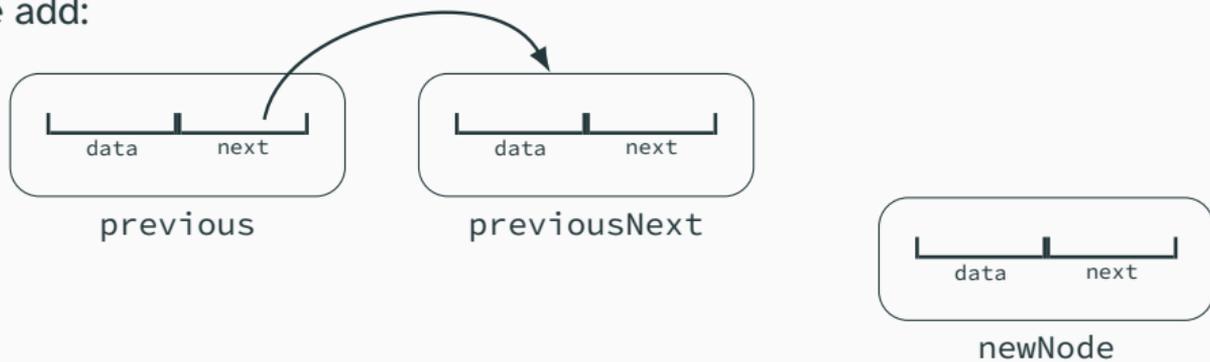
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- (Diagram on next slide)

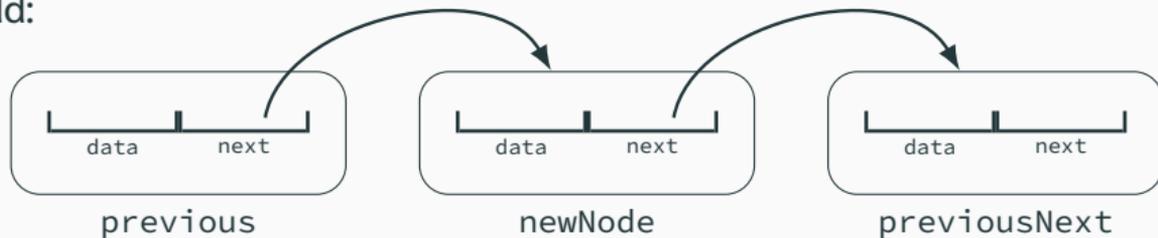
# Linked List Surgery

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Before add:



After add:



## **More on Memory**

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# Arrays

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- Let's look back at the `expandCapacity` method of `ArrayListInt.java`

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- When we use it to access instance variables, it's like "giving directions"
- Also used in `CoinStrip.java`

```
1 public class Student{
2     String name;
3     int graduationYear;
4
5     public Student(String name, int graduationYear) {
6         //this.name is instance variable; name is parameter
7         this.name = name;
8         this.graduationYear = graduationYear;
9     }
10 }
```

## Multiple References to the Same Object

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In pairs: what happens in this example?

```
1 Student s1 = new Student("Victoria", 2005);
2 Student s2 = new Student("Victoria", 2005); //s2 references
    a new student with the same data
3 System.out.println(s2.getGraduationYear()); //prints?
4 System.out.println(s1.getGraduationYear()); //prints?
5 s2.setGraduationYear(1997);
6 System.out.println(s1.getGraduationYear()); //prints?
7 System.out.println(s2.getGraduationYear()); //prints?
```

## Multiple References to the Same Object

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In pairs: what happens in *this* example? (Second line changed!)

```
1 Student s1 = new Student("Freida",2005);
2 Student s2 = s1;
3 System.out.println(s2.getGraduationYear()); //prints?
4 System.out.println(s1.getGraduationYear()); //prints?
5 s2.setGraduationYear(1997);
6 System.out.println(s1.getGraduationYear()); //prints?
7 System.out.println(s2.getGraduationYear()); //prints?
```

## Multiple References to the Same Object

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```
1 Student s1 = new Student("Freida",2005);
2 Student s2 = s1; //s2 and s1 reference the same place
3 System.out.println(s2.getGraduationYear()); //prints 2005
4 System.out.println(s1.getGraduationYear()); //prints 2005
5 s2.setGraduationYear(1997);
6 System.out.println(s1.getGraduationYear()); //prints 1997
7 System.out.println(s2.getGraduationYear()); //prints 1997
```

## `.equals()`

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## .equals()

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- Special Java method for objects
- Always test equality of objects using `.equals()`
  - If you use `==` it tests if the *references* are the same
  - Need to fill in yourself when you create a class (we'll come back to this)
- Always use `.equals()` for Strings!

```
1 if(string1.equals(string2)) {
2     //string 1 and string 2 are equal
3 }
4 if(string1 == string2) {
5     //if string 1 and string 2 are same, usually get here (
6     //but not always)
7     //don't use this!
8 }
```

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- Let's look at the main method of `LinkedListInt.java`

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- To `add()` to the end need to traverse the entire list!

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- Modification to the linked list from last time
  - What we implemented Monday in `LinkedListInt.java` is a “Singly Linked List”
- In a singly linked list, you can only go forward to the next node; can't go backwards
- To `add()` to the end need to traverse the entire list!
- In a **Doubly Linked List**, we add in the capability to go forwards or backwards

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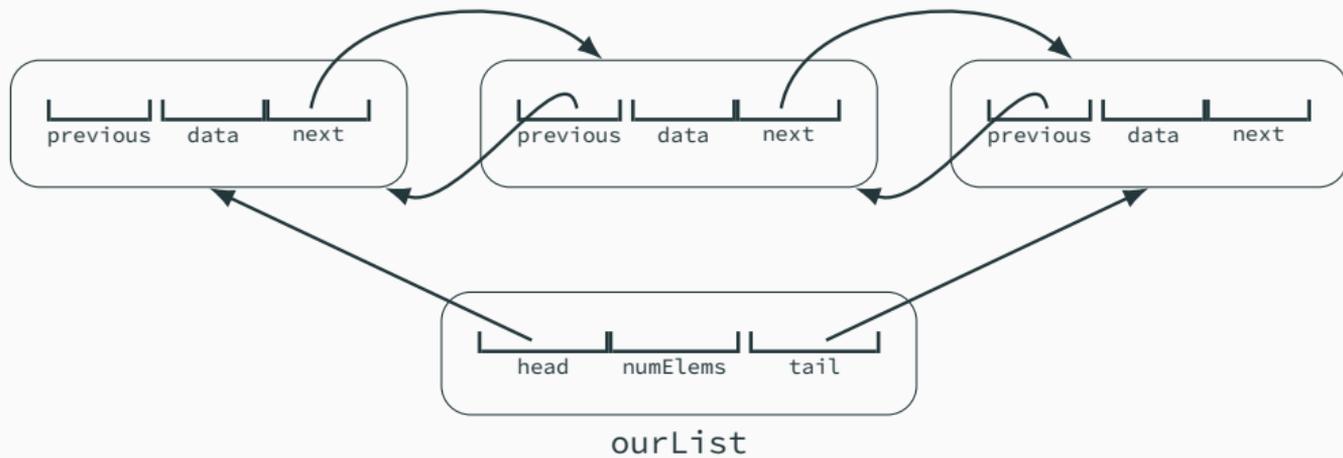
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- If the list has one element, `head` and `tail` are the same

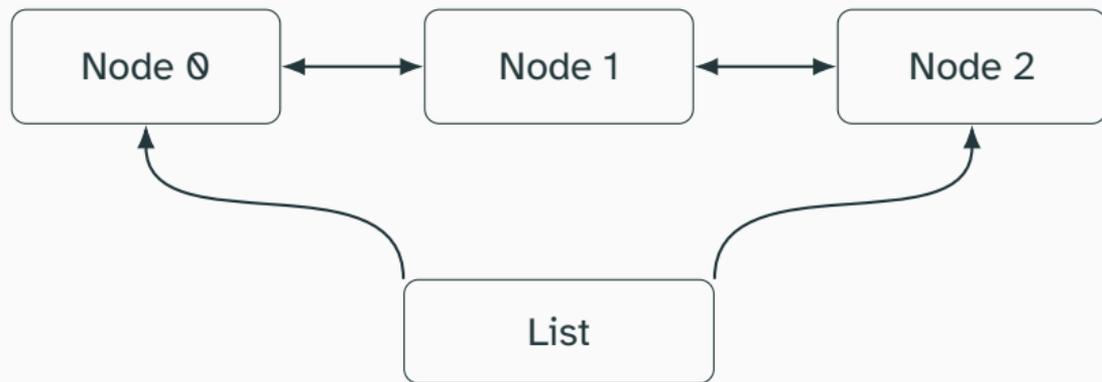
# Doubly Linked List

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## Doubly Linked List: Simplified Diagram

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## Doubly Linked List: Some Code

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- How can we fill in `add()` (at the end of the list)?
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- How can we fill in our private method `getNode()`?

## Doubly Linked List: Some Code

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- Let's look at `DLLNodeInt.java`
- How can we fill in `add()` (at the end of the list)?
  - Let's do this **in pairs** if it's before 9:35
- How can we fill in our private method `getNode()`?
  - Do we want to traverse the list forward or backward?

## **Comparing Efficiency**

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# What We Mean By Efficiency

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- We'll talk soon about how to analyze efficiency rigorously
- For today: does the method need to traverse the entire data structure?
- If your list has millions or billions of elements, traversing the data structure takes lots of time
- Example: `indexOf()` always traverses the whole data structure. It's a relatively slow operation

# Efficiency of Methods

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- What about `set()`?
- What about `add()`?