CSCI 136: Data Structures and Advanced Programming Lecture 23 Trees, part 3

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Williams

Tree terminology

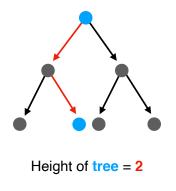
Your to-dos

- 1. Read before Fri: review Bailey, Ch 14.
- 2. Lab 8 (solo lab), due Tuesday 11/15 by 10pm. Note: you will implement a tree data structure called a **trie** for lab 8; the structure is described in the lab handout. Please bring a short design document to your lab meeting.

Activity: Binary Tree Height

Topics

The **height** of a tree is the length of the longest path between the root and any leaf.



Binary Tree Height

Let's think about some corner cases.

What is the height of a tree with just one node?

The **height** of a tree is the length of the longest path between the root and any leaf.

Height of tree = 0

Binary Tree Height

Let's think about some corner cases.

What about the empty tree?



The **height** of a tree is the length of the longest path between the root and any leaf.

Height of tree = -1

Binary Tree Height

Here's a more formal definition.

The **height** of a tree is defined as:

- -1 if the tree is empty, or
- height(left) or height(right), whichever is bigger, + 1



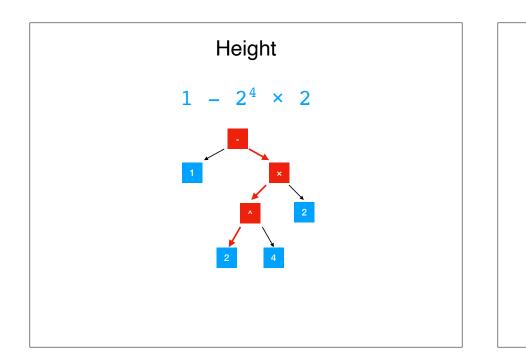
empty tree: -1

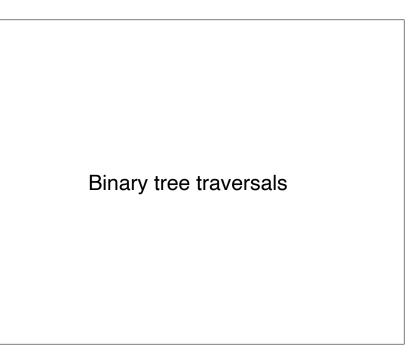


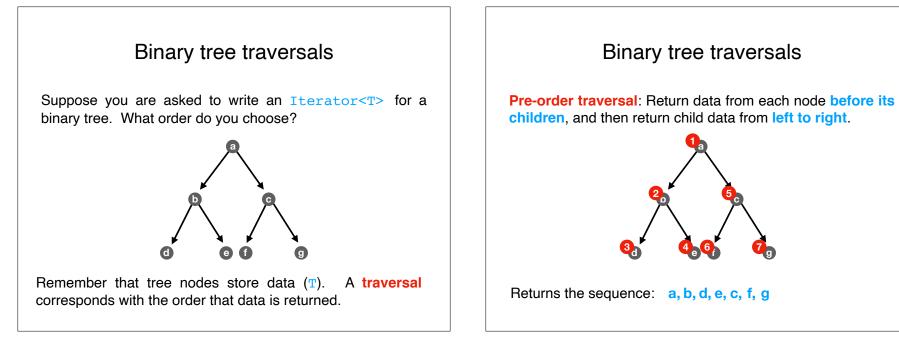
any other tree: longest path

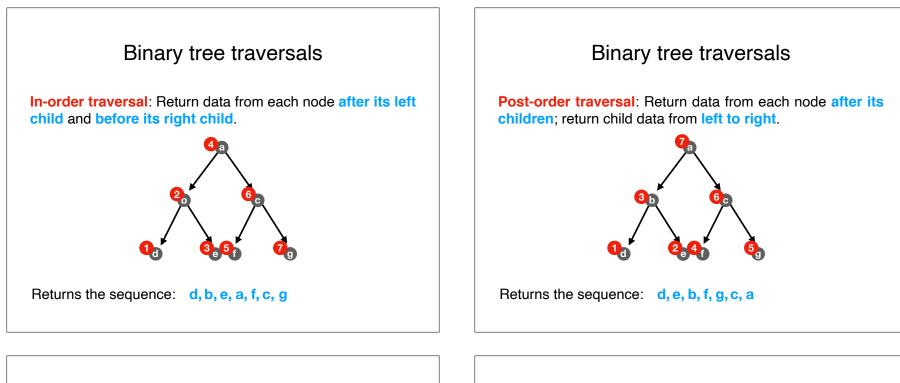
Activity: Binary Tree Height

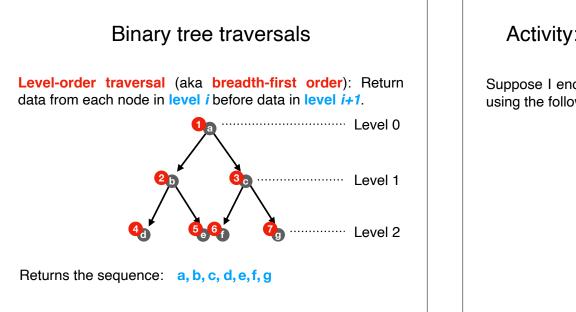
How might we implement getHeight()?





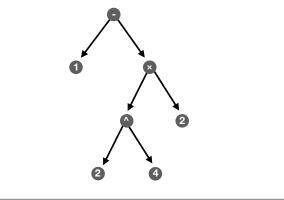






Activity: What traversal should I use?

Suppose I encode the arithmetic expression $1 - 2^4 \times 2$ using the following tree.



Ordered Trees

Binary search tree (alternative)

A **binary search tree** is a binary tree that maintains the **binary search property** as elements are added or removed. In other words, the **key** in each node:

• must be \geq any **key** stored in the left subtree, and

• must be < any key stored in the right subtree.

As with other ordered structures, order is maintained on insertion.

Binary search tree

A **binary search tree** is a binary tree that maintains the **binary search property** as elements are added or removed. In other words, the **key** in each node:

must be > any key stored in the left subtree, and
must be ≤ any key stored in the right subtree.

As with other ordered structures, order is maintained on insertion.

Key, Value nodes

Note that I said key instead of element.

Storing a **key** and a **value** in each node allows the greatest flexibility when arranging a tree. I.e., the key type K need not be the value type V.

<u>Restriction</u>: keys must be **comparable** in some way (e.g., Comparable<K> or Comparator<K>).

Example

Insert the following elements: 71, 20, 27, 17, 91, 14, 87

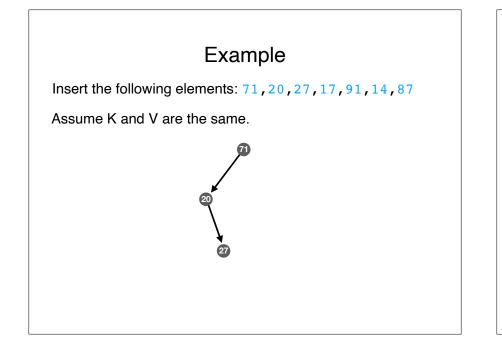
Assume K and V are the same.

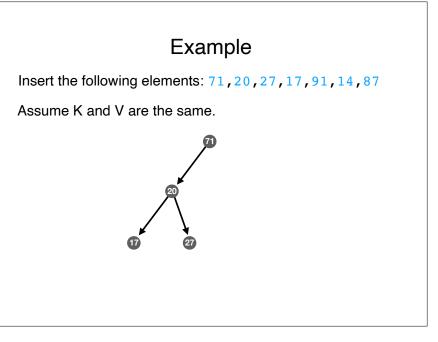
Example

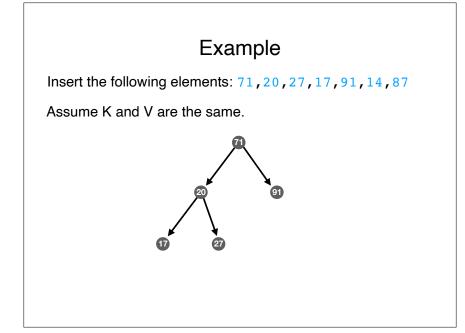
Insert the following elements: 71, 20, 27, 17, 91, 14, 87

Assume K and V are the same.





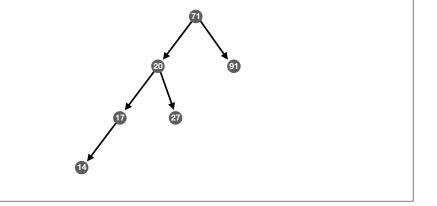


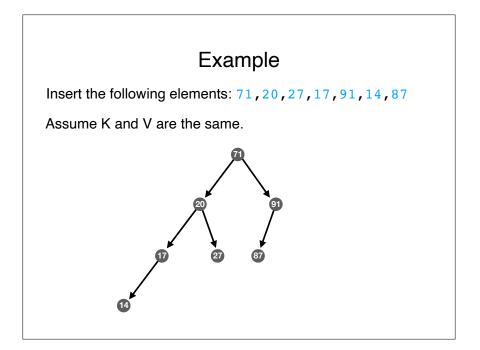


Example

Insert the following elements: 71, 20, 27, 17, 91, 14, 87

Assume K and V are the same.







Insert the following elements:

Assume K and V are the same.

Binary Search Tree

Let's start implementing this together.

Recap & Next Class

Today:

Tree terminology

Tree traversals

Next class:

Binary search trees