CSCI 136: Data Structures and Advanced Programming Lecture 23 Trees, part 3 Instructor: Dan Barowy

Williams

Tree terminology

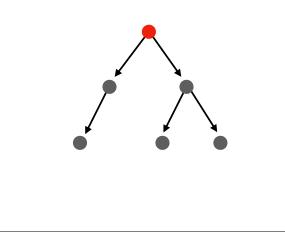
Topics

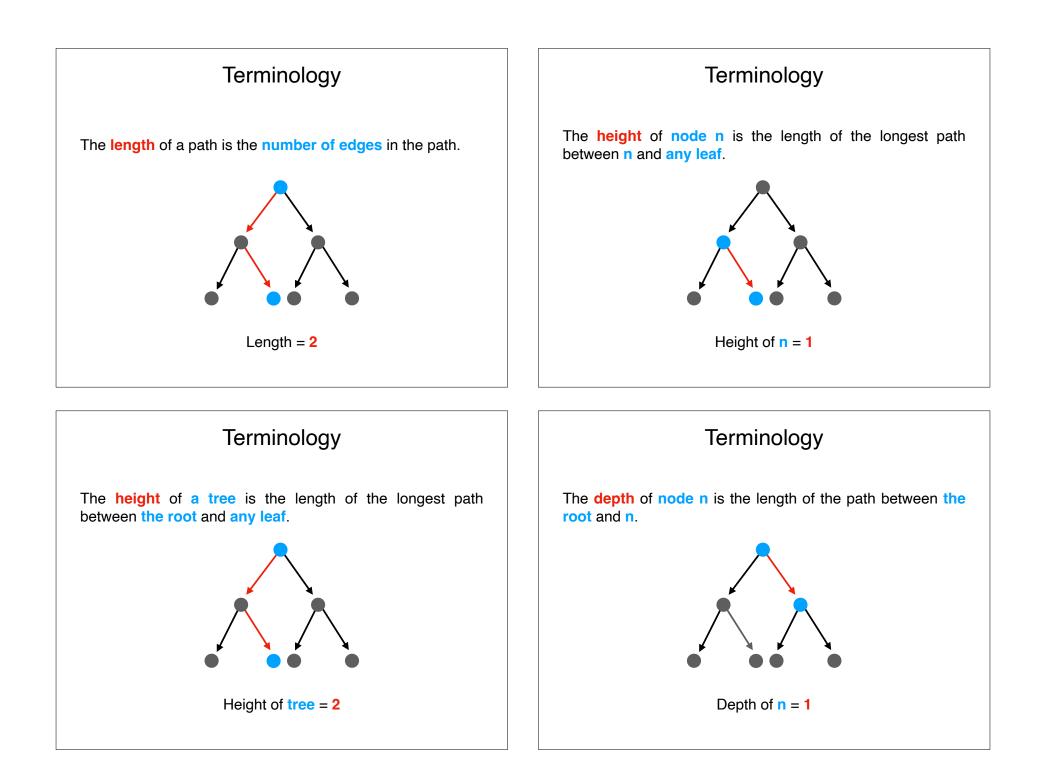
Your to-dos

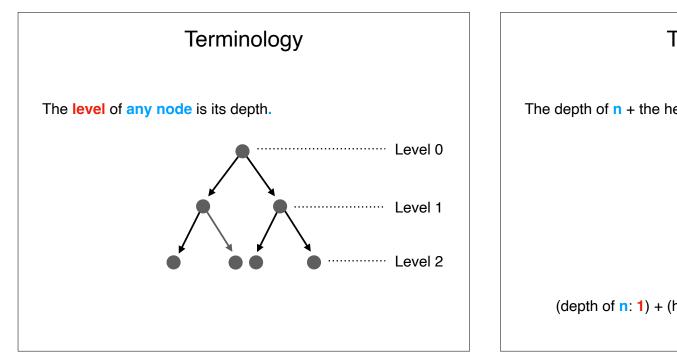
- 1. Read before Fri: review Bailey, Ch 14.
- 2. Lab 8 (solo lab), due Tuesday 11/15 by 10pm. Note: you will implement a tree data structure called a **trie** for lab 8; the structure is described in the lab handout. Please bring a short design document to your lab meeting.

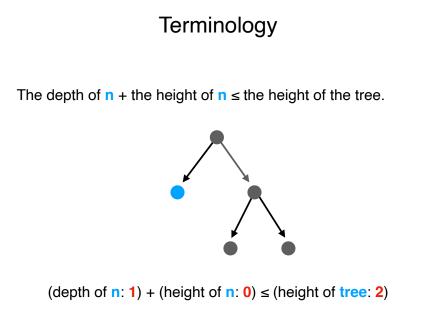
Properties of trees

**Cycle-free**: no path will ever revisit the same node.





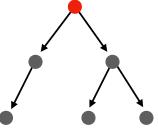




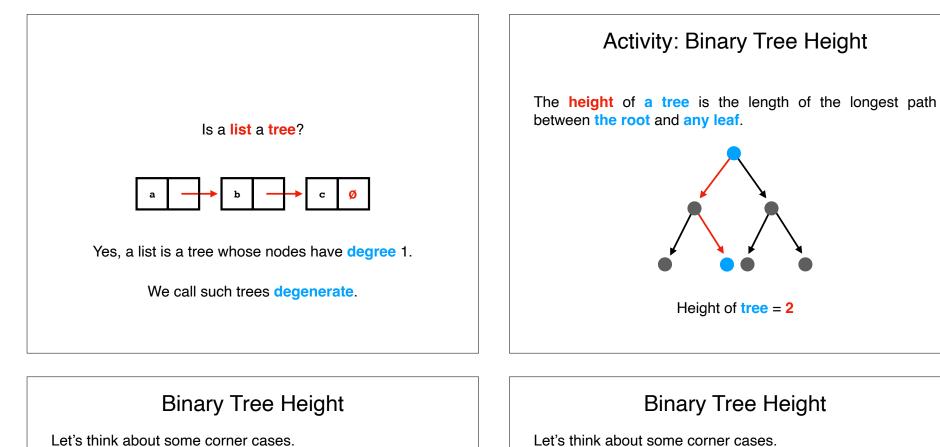
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# Properties of trees

**Directed** or **undirected**: trees can be either directed, meaning that traversals can only happen in one direction, or undirected, meaning that traversals can happen in any direction.



The tree shown here is directed. We can represent an undirected tree using back edges.



What is the height of a tree with just one node?

The height of a tree is the length of the longest path between the root and any leaf.

Height of tree = 0

# **Binary Tree Height**

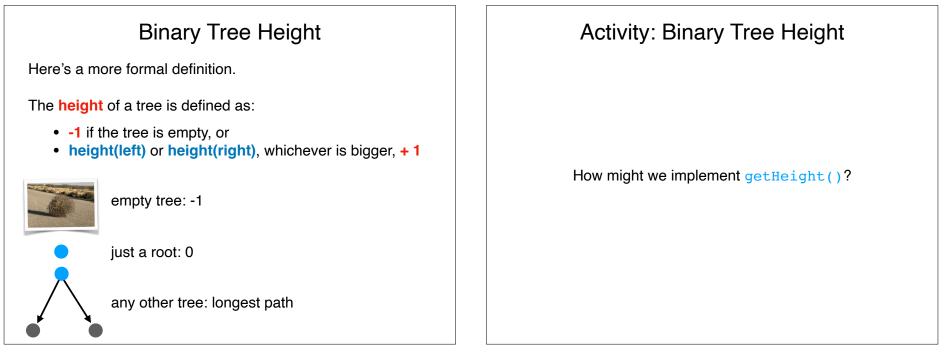
Let's think about some corner cases.

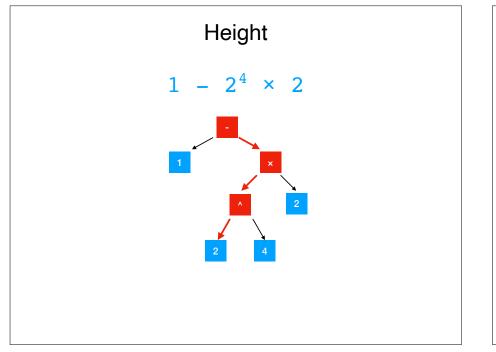
What about the empty tree?



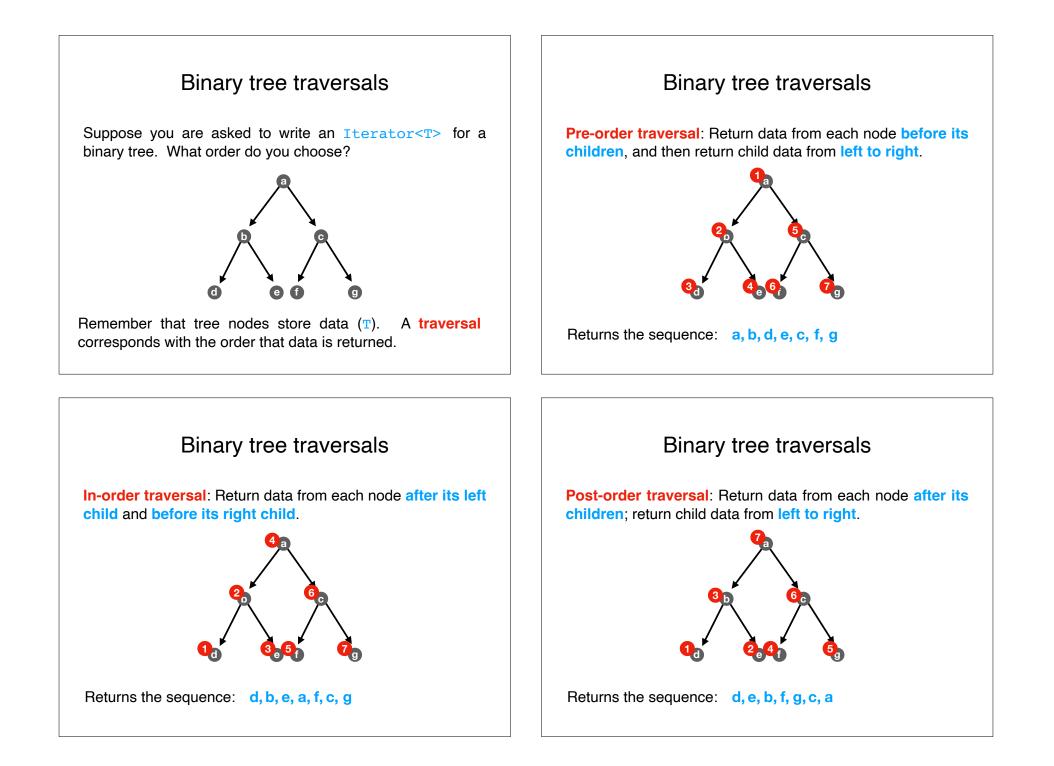
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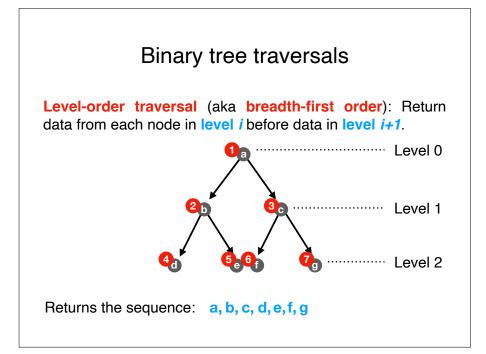
Height of tree = -1





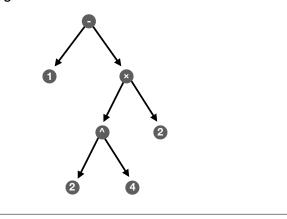
Binary tree traversals

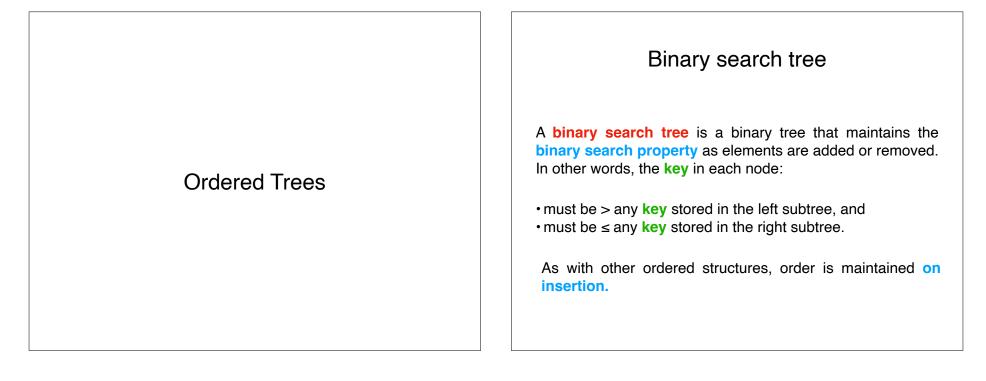




#### Activity: What traversal should I use?

Suppose I encode the arithmetic expression  $1 - 2^4 \times 2$  using the following tree.





#### Binary search tree (alternative)

A **binary search tree** is a binary tree that maintains the **binary search property** as elements are added or removed. In other words, the **key** in each node:

must be ≥ any key stored in the left subtree, and
must be < any key stored in the right subtree.</li>

As with other ordered structures, order is maintained on insertion.

#### Key, Value nodes

Note that I said key instead of element.

Storing a **key** and a **value** in each node allows the greatest flexibility when arranging a tree. I.e., the key type K need not be the value type V.

<u>Restriction</u>: keys must be **comparable** in some way (e.g., Comparable<K> or Comparator<K>).

#### Example

Insert the following elements: 71, 20, 27, 17, 91, 14, 87

Assume K and V are the same.

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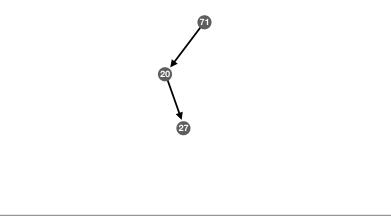
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Insert the following elements: 71, 20, 27, 17, 91, 14, 87

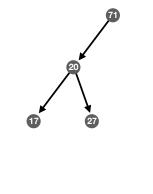
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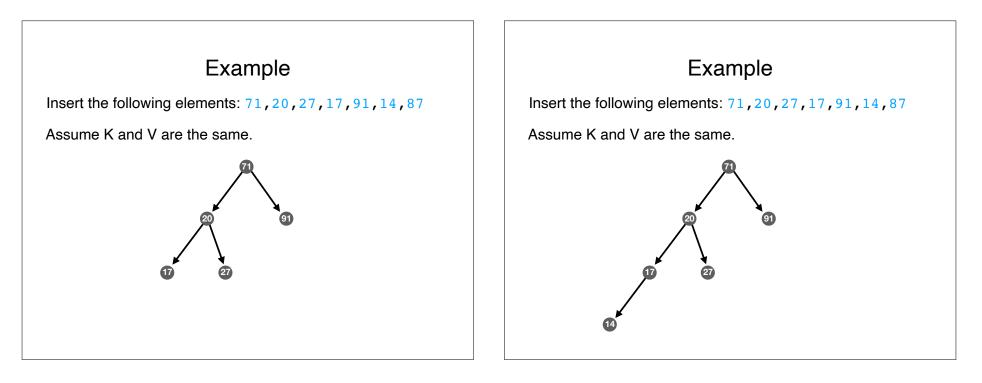


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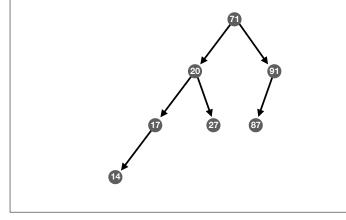




#### Example

Insert the following elements: 71, 20, 27, 17, 91, 14, 87

Assume K and V are the same.



#### **Binary Search Tree**

Let's start implementing this together.

# Activity

Insert the following elements:

Assume K and V are the same.

#### Recap & Next Class

### Today:

Tree terminology

Tree traversals

# Next class:

Binary search trees