CSCI 136:
Data Structures
and
Advanced Programming
Lecture 17

Linear structures, part 2

Instructor: Dan Barowy

Williams

Topics

- Stack data structure
- Queue ADT
- Queue data structure
- Resubmission procedure

Your to-dos

- 1. Lab 6 (partner lab), due Tuesday 11/1 by 10pm. (two weeks!)
- 2. Review readings for midterm.

Announcements

- Colloquium: What I Did Last Summer (Research),
 2:35pm in Wege Auditorium with cookies.
- Practice midterm posted on the course website.
- TA feedback.



Announcements

Please consider being a TA next semester (especially for this class!)

Applications due Friday, October 28.

https://csci.williams.edu/tatutor-application/

Queue ADT

A queue is an abstract data type that stores a collection of any type of element. A queue restricts which elements are accessible: elements may only be added to the "end" of the collection and elements may only be removed from the "front" of a collection. The "enqueue" operation places an element at the end of a queue while a "dequeue" operation removes an element from the front.

Queue ADT



Queue ADT

Also sometimes referred to as a FIFO: "first in, first out."

(a stack would be an annoying way to process a line at Starbucks!)

Frequently used as a **buffer** to hold work to do later.

We also frequently include a "peek" operation that lets us look at an element on the top of a queue without removing it, and "size" and "isEmpty" operations that let us check how many elements are stored and whether a queue stores zero elements, respectively.

Queue implementations

QueueArray

A QueueArray is a queue implemented using an array for element storage.

Pros: enqueue and dequeue are O(1) operations.

Cons: data structure has a maximum capacity.

Queue implementations

QueueVector

A Queue Vector is a queue implemented using a Vector for element storage.

Pros: enqueue and dequeue are amortized O(1) operations. There is no maximum capacity.

Cons: Most of the time, they take O(1) time, but occasionally-when the underlying array needs to grow--an O(n) cost is incurred. This may be fine for most applications, but if the application cannot tolerate wide variation in time, this is a bad choice. Also, unless the underlying array is completely full, Vectors waste some space.

Queue implementations

QueueList

A QueueList is a queue implemented using a List (usu. DLL or CL) for element storage.

Pros: enqueue and dequeue are O(1) operations. There is no maximum capacity. enqueue and dequeue costs are predictable (always the same), unlike QueueVector.

Cons: because of the way computer hardware is implemented, a **QueueList**'s constant-time cost is likely to be much higher than a **QueueVector**'s. So a **QueueList**'s performance may be more predictable than a **QueueVector**, but it will likely be slower on average.

Other queue-like ADTs

One very useful and interesting variant of the Queue ADT is the **Priority Queue ADT**. We'll talk about priority queues after the midterm!

Resubmission procedure

Resubmission procedure



Remember: the goal of this course is mastery.

Resubmission procedure

Allows you to earn up to 50% of the lost points.

E.g., if you got a 50% on the midterm, you can get a 75% on resubmission.

Midterm is 25% of your final grade.

This is worth doing!

Resubmission procedure

- You have until the end of reading period.
- Resubmission must include both the original work and the new submission.
- 3. Must be accompanied by an **explanation document**, written in plain English.

Resubmission procedure

Explanation document must identify:

- 1. What the mistake is.
- 2. How you fixed the mistake.
- 3. Why the new version is correct.

Resubmission procedure

Resubmit code **electronically** (i.e., using git).

Resubmit exam on paper (i.e., hand it to me or put in mailbox).

Resubmission procedure Sample from CS334:

2. Troubleshooting

My fix was slightly wrong. Righr before calling random_string(), I added

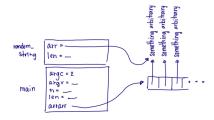
char * arrarr[i] = malloc(sizeof(char)*MAXLEN);

when what I should have added is

arrarr[i] = malloc(sizeof(char)*MAXLEN);
mcheck(arrarr[i]);

There is no need for "char *" because I am not declaring arrarr.

I got my explanation and drawing wrong. In my drawing, I had arrarr[i] pointing back to a call stack because I thought the program would automatically allocate memory on a call stack if we did not malloc(). What I should have said is that without allocating sub-array arrarr[i], the address currently living in the sub-array is arbitrary so the referred to by the sub array is also arbitrary. When we call memset() or manipulating arrarr[i] in $random_string()$, we are likely to get memory errors. Below is what I should have drawn.



Iterators

What do the following have in common?

```
double[] a
// ... initialize a ...
double sum = 0.0;
for (int i = 0; i < a.length; i++) {
    sum += a[i];
}

List<Double> ls = new SinglyLinke
// ... initialize ls ...
double sum = 0.0;
for (int i = 0; i < ls.size(); i
    sum += ls.get(i);
}

Stack<Double> s = new StackVect
// ... initialize s ...
double sum = 0.0;
while (!s.isEmpty()) {
    sum += s.pop();
}
```

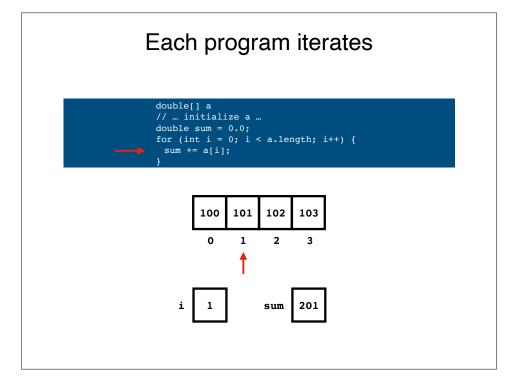
Iteration

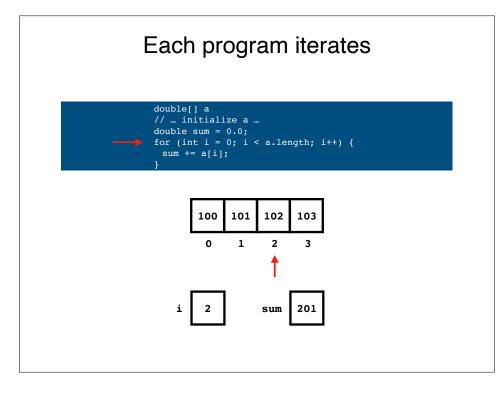
Iteration is the **repetition of a process** in order to generate a (possibly unbounded) **sequence of outcomes**. Each step in an iteration performs the given process **once**; the result of each step is the **starting point** of the **next step**.

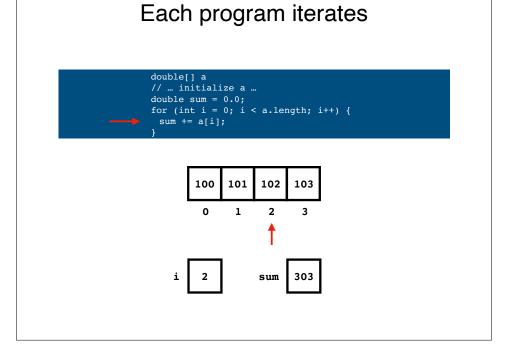
Each program iterates

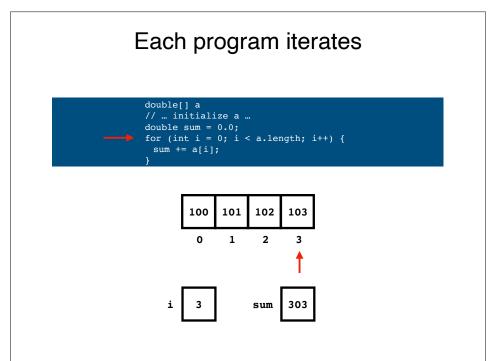
Each program iterates

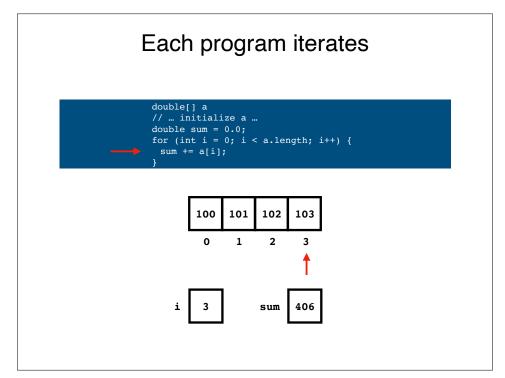
double[] a // ... initialize a ... double sum = 0.0; for (int i = 0; i < a.length; i++) { sum += a[i]; }</pre> i 1 sum 100

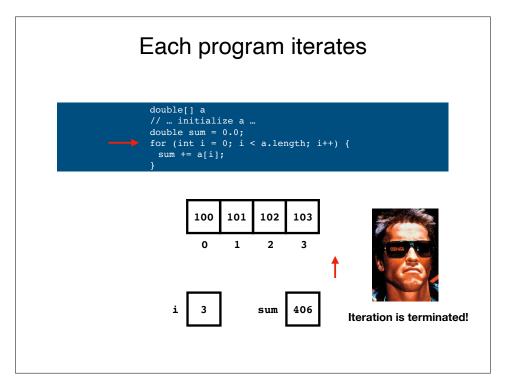


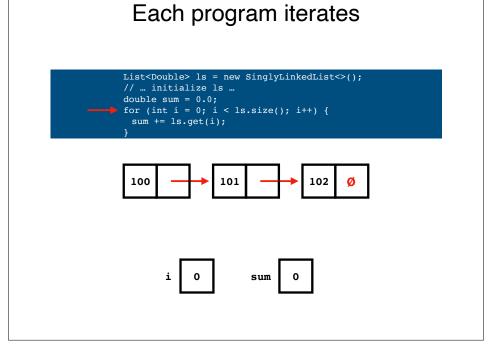


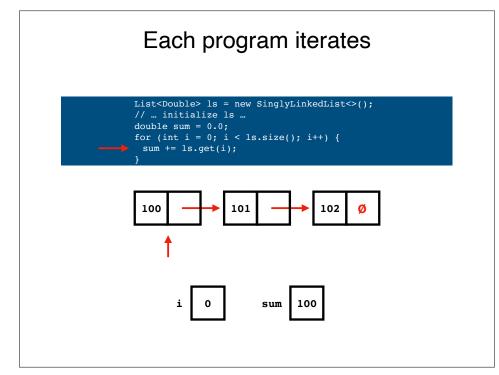


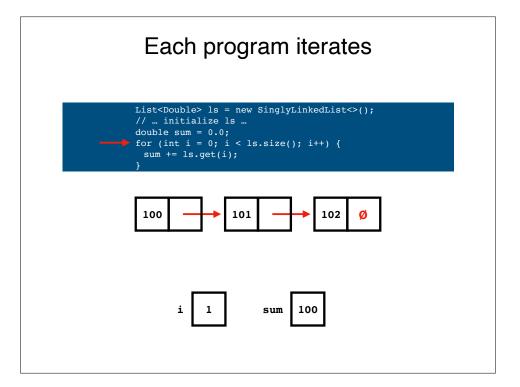


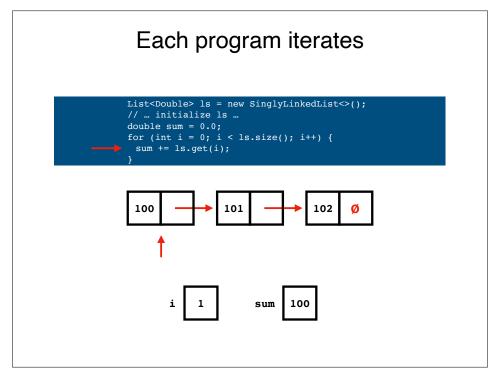


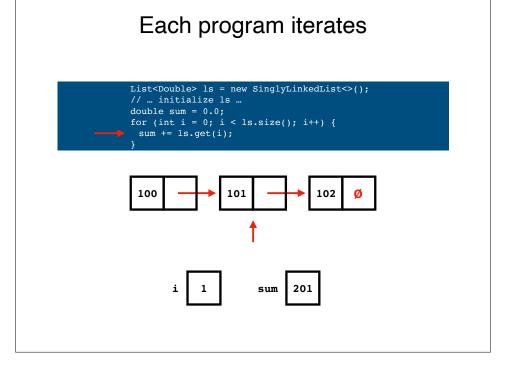


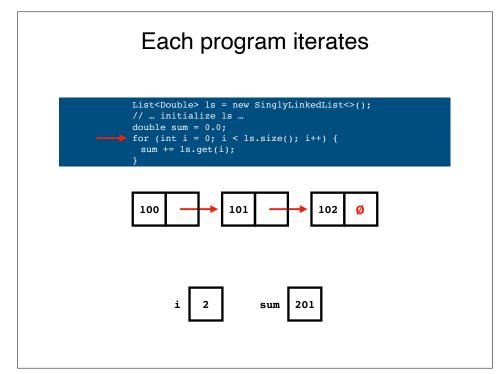


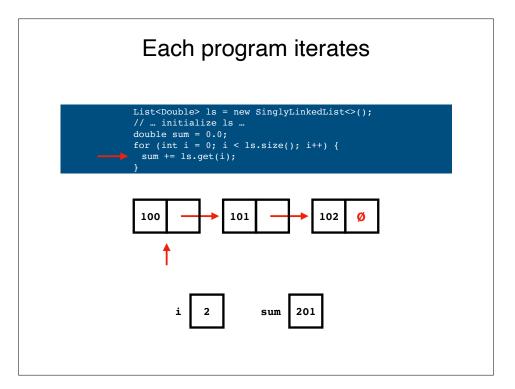


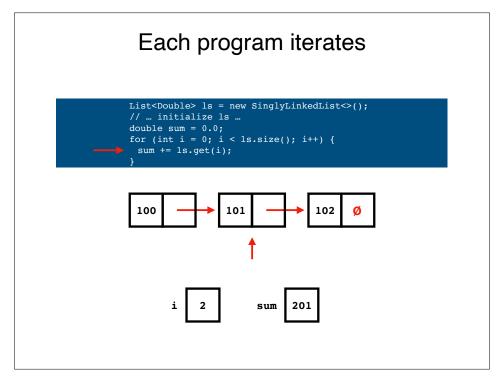


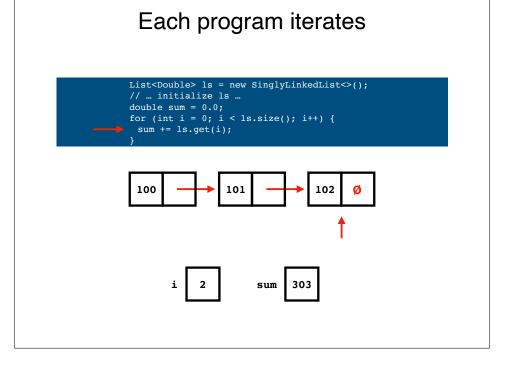


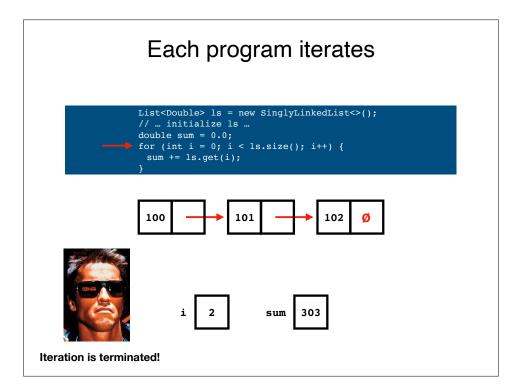


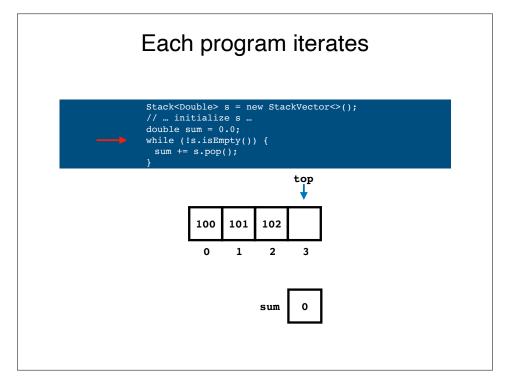


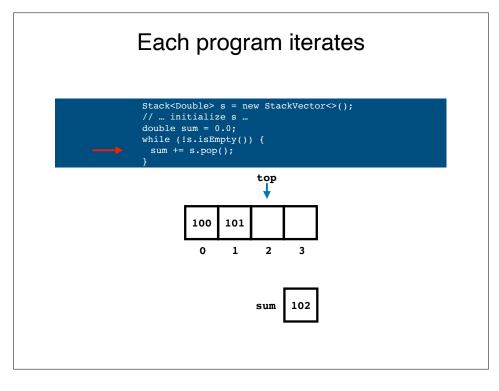


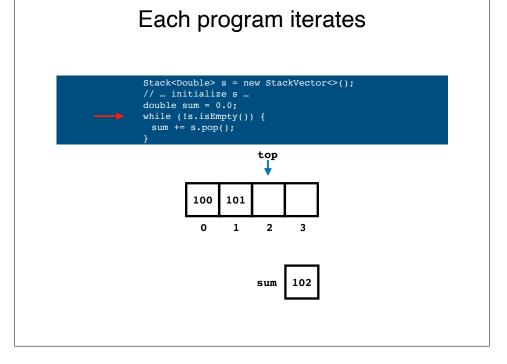


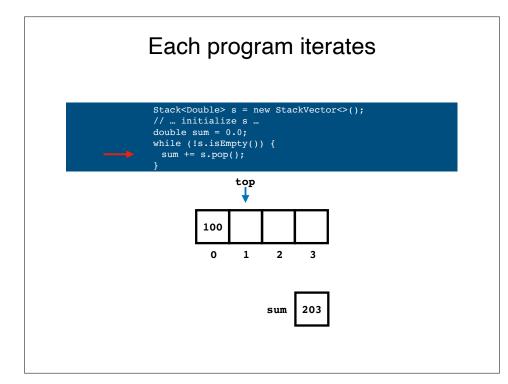


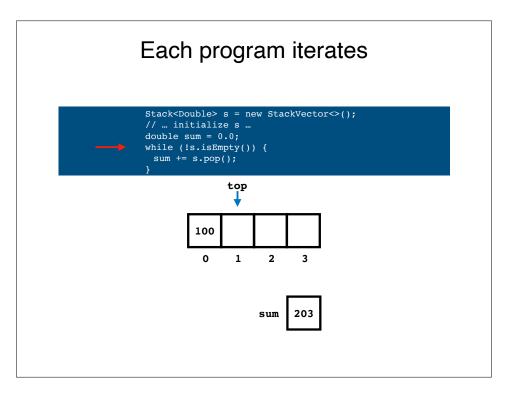


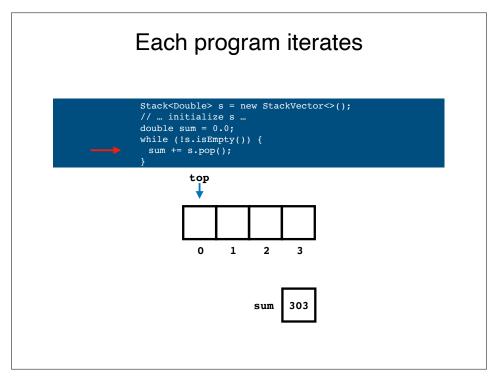


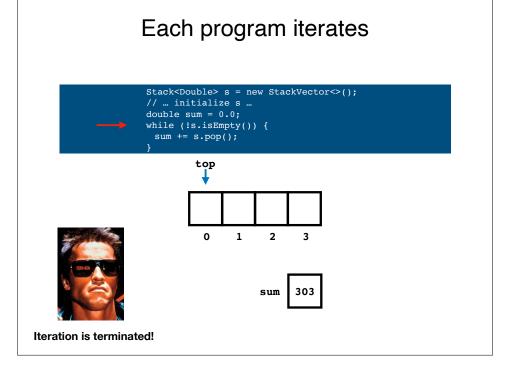












Essentially the same algorithm!

```
double[] a
// ... initialize a ...
double sum = 0.0;
for (int i = 0; i < a.length; i++) {
    swn += a[i];
}

List<Double> ls = new SinglyLinkedList<>();
// ... initialize ls ...
double sum = 0.0;
for (int i = 0; i < ls.size()) i++) {
    syn += ls.get(i);
}

Stack<Double> s = new StackVector<>();
// ... initialize s ...
double sum = 0.0;
while [!s.ieEmpty()) {
    swn += s.pop();
}
```

But the code looks different.

Problems

- Different data structures yield different code for same algorithm.
- Data hiding potentially causes efficiency problems.
- Inspecting data structure "from the outside" can change the state of a data structure (e.g., pop() ing a Stack).

What if I told you that you could solve



all of these problems with **abstraction**?

Iteration abstraction to the rescue.

```
double[] a
// ... initialize a ...
double sum = 0.0;
for (double d : a) {
   sum += d;
}
```

```
List<Double> ls = new SinglyLinkedList<>();
// ... initialize ls ...
double sum = 0.0;
for (double d : ls) {
   sum += d;
}
```

```
Stack<Double> s = new StackVector<>();
// ... initialize s ...
double sum = 0.0;
for (double d : s) {
   sum += d;
}
```

Brought to you by **Iterators**.

Iterators are a really good idea.

- Invented by Barbara Liskov in 1974.
- Incidentally, abstract data types were also invented by Barbara Liskov in 1974.
- (1974 was a good year!)
- Both debuted in an influential PL named CLU.
- Barbara won the Turing Award in 2008 for this work and more.



Recap & Next Class

Today:

Queue ADT

Queue implementation

Resubmissions

Next class:

Iterators

Search

Remember this from Python?

```
for num in nums:
    # do something
```

Java has it too!

```
for (int num : nums) {
   // do something
}
```

(admittedly, it is less pretty in Java)

This is called a "for-each loop"