	Topics
CSCI 136: Data Structures and Advanced Programming	•Linear ADTs
Lecture 16	Stack ADT
Linear structures	
Instructor: Dan Barowy	
Williams	
Your to-dos	Announcements
	<ul> <li>Colloquium: What I Did Last Summer (Industry), 2:35pm in Wege Auditorium with cookies.</li> </ul>
<ol> <li>Lab 6 (partner lab), due Tuesday 11/1 by 10pm. (two weeks!)</li> </ol>	<ul> <li>Midterm: in lab two weeks from now:</li> <li>Wed, October 26 and</li> <li>Thu, October 27 and</li> </ul>
2. Read <b>before Fri</b> : Bailey, Ch 8-8.3.	<ul> <li>Midterm review: Mon, October 24 in class.</li> <li>No class: Fri, October 28.</li> </ul>

## Radix sort

(a sort where stability matters)

## Recall: Abstract Data Type

An **abstract data type** is a mathematical formulation of a data type. ADTs abstract away **accidental** properties of data structures (e.g., implementation details, programming language). Instead, ADTs contain only **essential** properties and are **concisely defined by their logical behavior** over a **set of values** and a **set of operations**.

In an ADT, precisely how data is **represented** on a computer **does not matter**.

#### By contrast: data structure

A **data structure** is the physical form of a data type, i.e., it is an implementation of an ADT. Generally, data structures are designed to efficiently support the logical operations described by the ADT.

For data structures, precisely how data is **represented** on a computer **matters a lot**. Simple data structures are often composed of simple representations, like primitives, while more complex data structures are composed of other data structures.

## ADT example: List

A **list** is a sequential collection of data elements, whose order is not necessarily given by their placement in memory. Elements may store **any type of value**. A list supports **inserting**, **searching** for, and **deleting** any value in a list, **at any location**, although not necessarily efficiently.

## Linear ADT

A linear ADT is one that presents elements in a sequence, even if the elements are not actually stored that way.

In a linear ADT, **adding** and **removing** elements is **constrained**, meaning that the structure can only be inspected and modified according to certain rules.

We will talk about two this week: **stack** and **queue**.

## Stack ADT

A stack is an abstract data type that stores a collection of any type of element. A stack restricts which elements are accessible: elements may only be added and removed from the "top" of the collection. The "push" operation places an element onto the top of the stack while a "pop" operation removes an element from the top.

#### Stack ADT



## Stack ADT

Also sometimes referred to as a LIFO: "last in, first out."

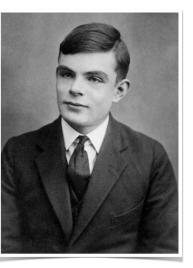
We also sometimes include a "**peek**" operation that lets us look at an element on the top of a stack without removing it, and "**size**" and "**isEmpty**" operations that let us check how many elements are stored and whether a stack stores zero elements, respectively.

## Stack ADT

Interesting history: first appeared in print in a paper by Alan Turing (1946).

Unclear if he actually invented it.

push = bury,
pop = unbury.



## **Application: Arithmetic**

A computer can perform arithmetic using a stack.

```
E.g., 1 + 2 * 3 = 7
```

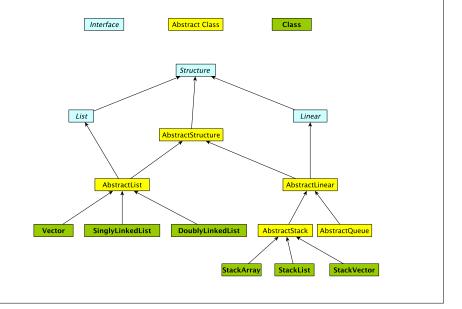
Small problem: order of operations in infix arithmetic depends on the operations themselves.

In postfix arithmetic, order is always the same: left to right

E.g., 1 2 3 \* + (note: fixed the confusing class example)

Once in this form, processing is easy. (Example)

## structure5 Stack implementations



## Activity: Arithmetic

Convert infix to postfix: x\*y+z\*w

1. Add parens to preserve order of operations:

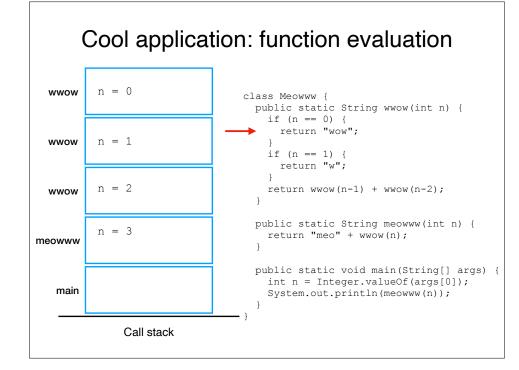
((x \* y) + (z \* w))2. Move all operators to the end of each parenthesized expression:

((x y \*)(z w \*) +)

3. Remove parens:

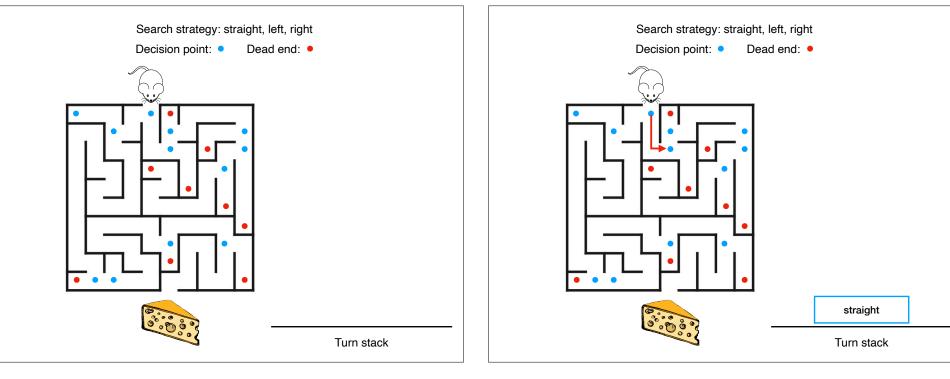
Evaluate these using a stack:

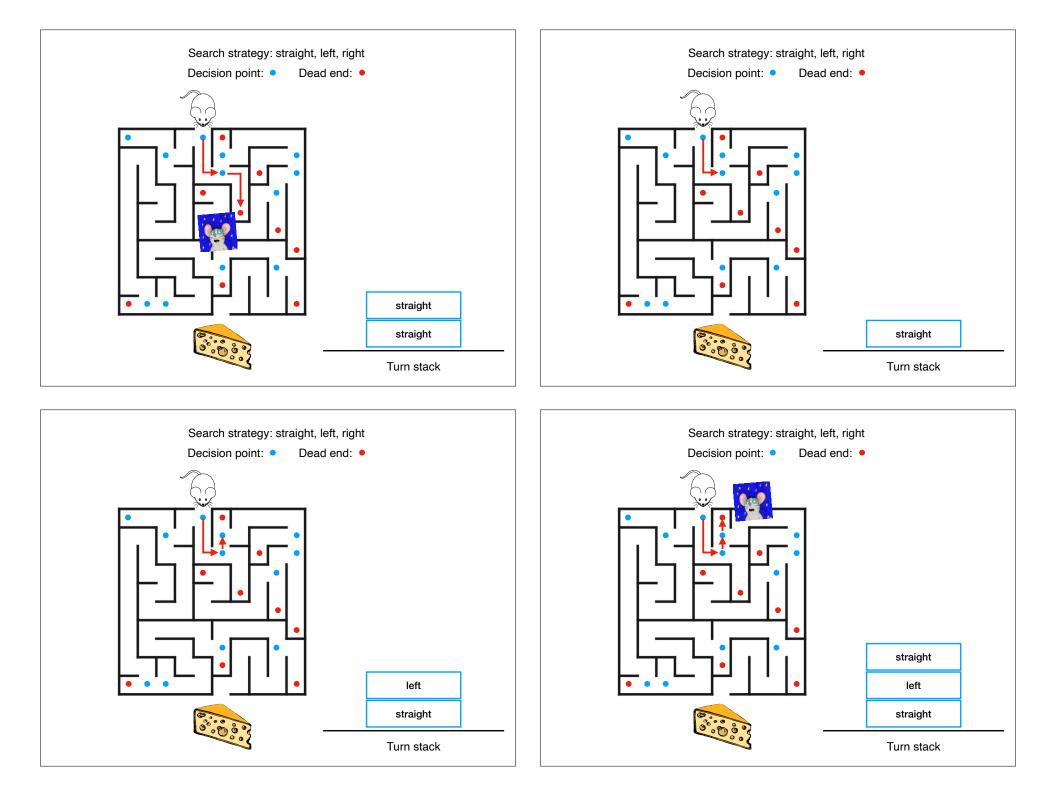
1.4 + 1 \* 82.5 \* (6 + 2) - 12 / 4

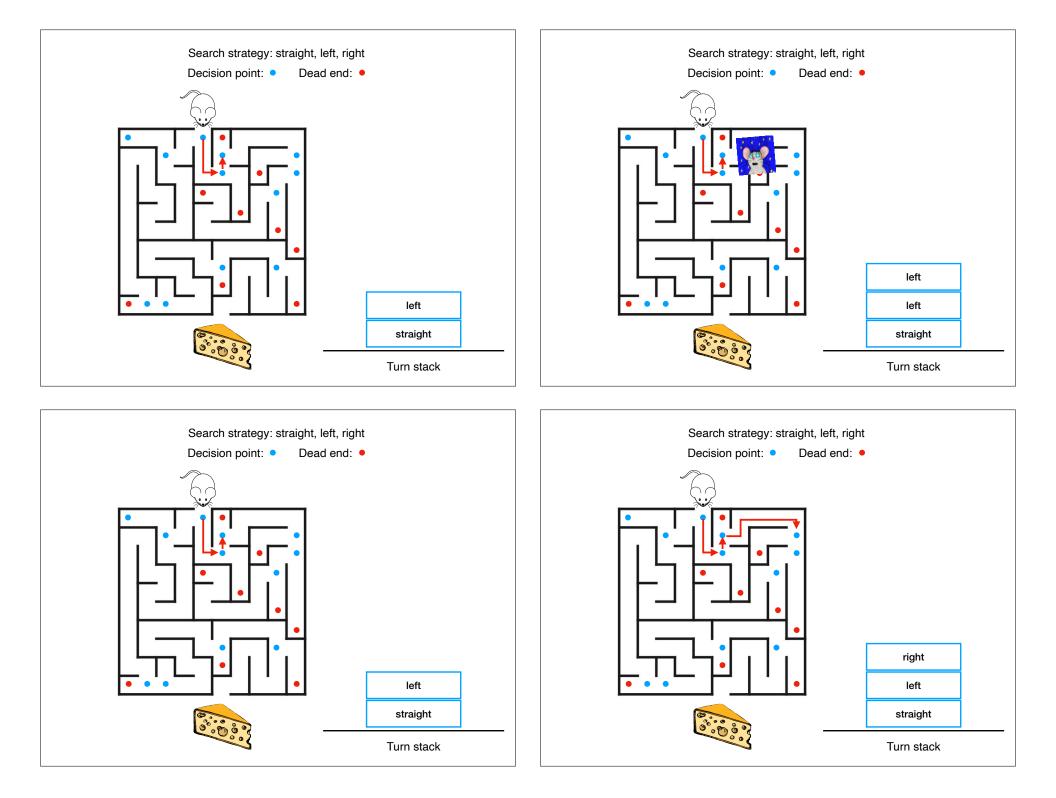


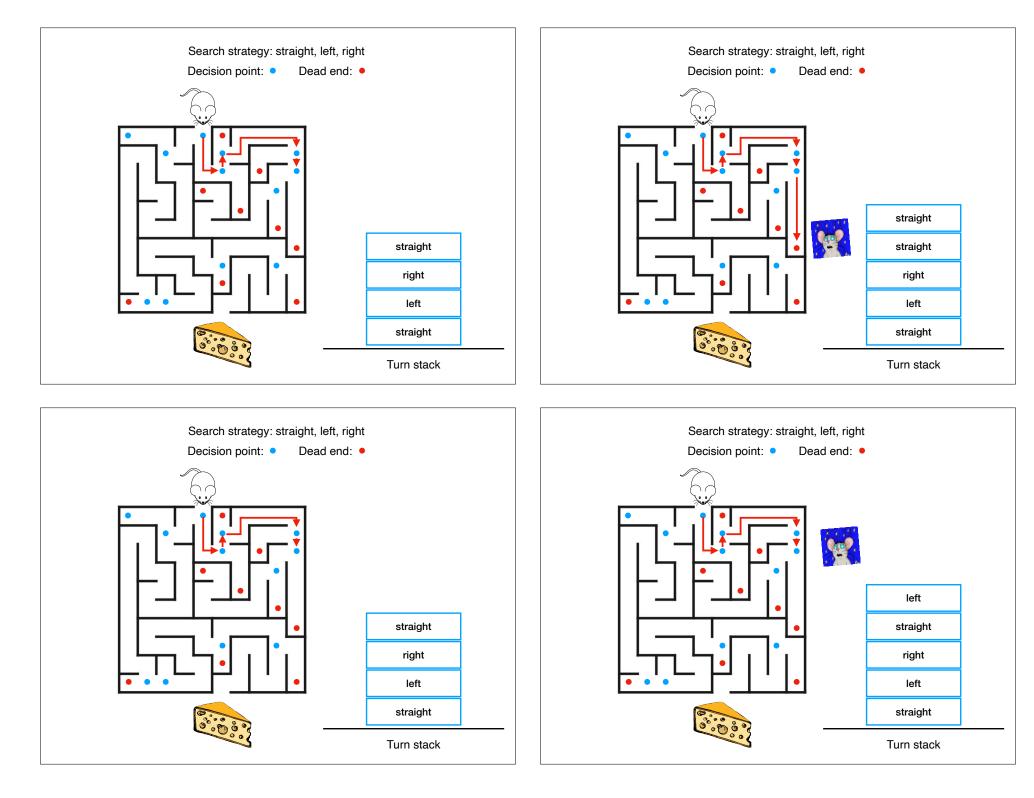
## Cool application: backtracking search

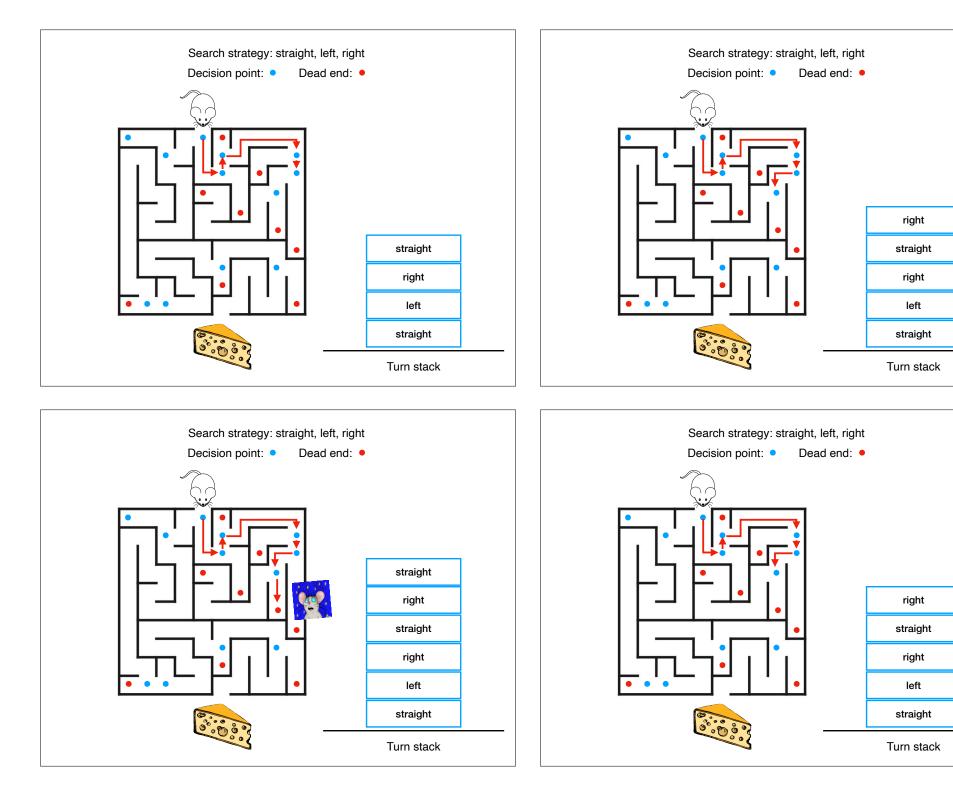


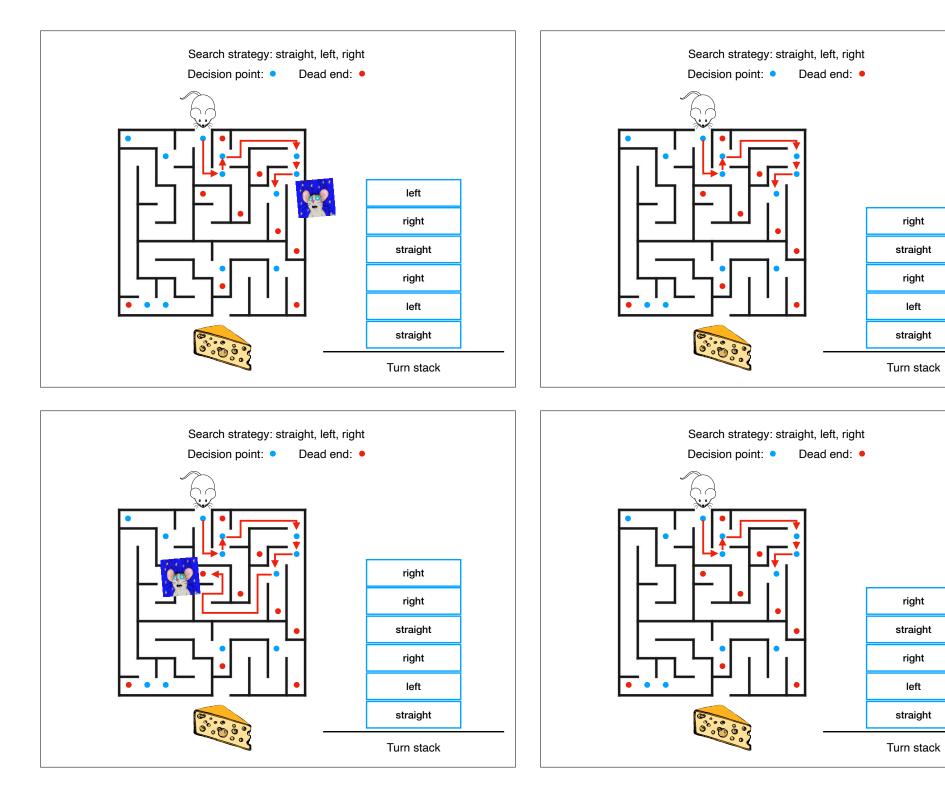


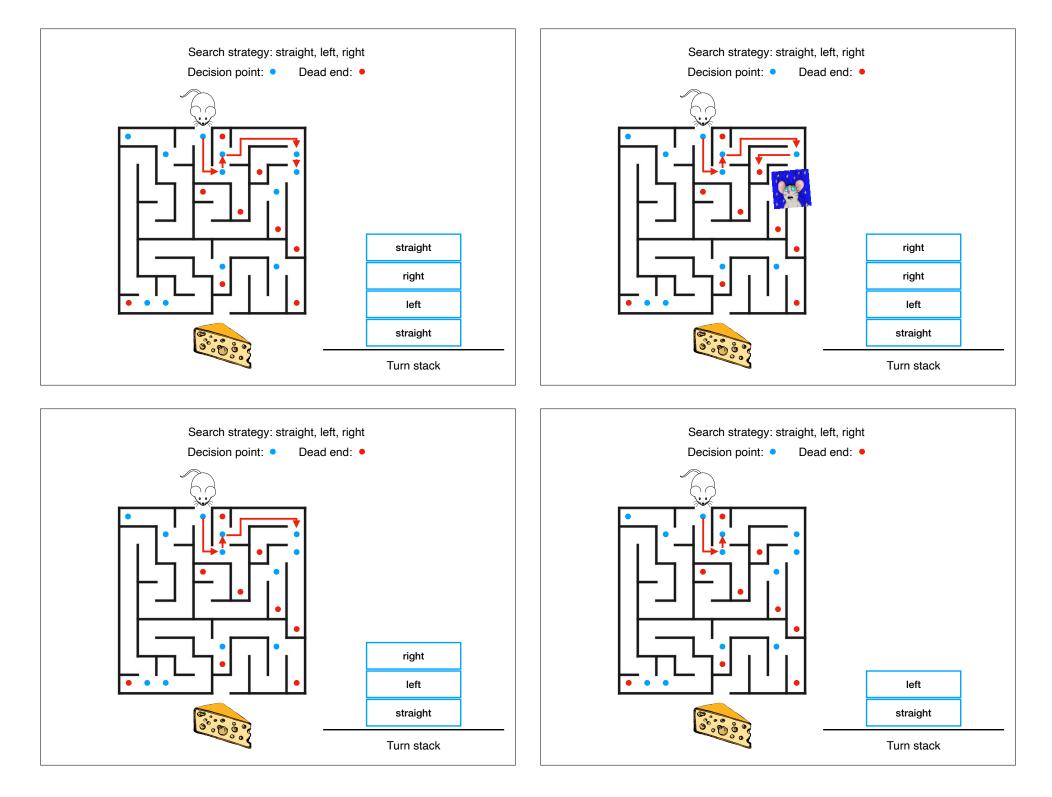


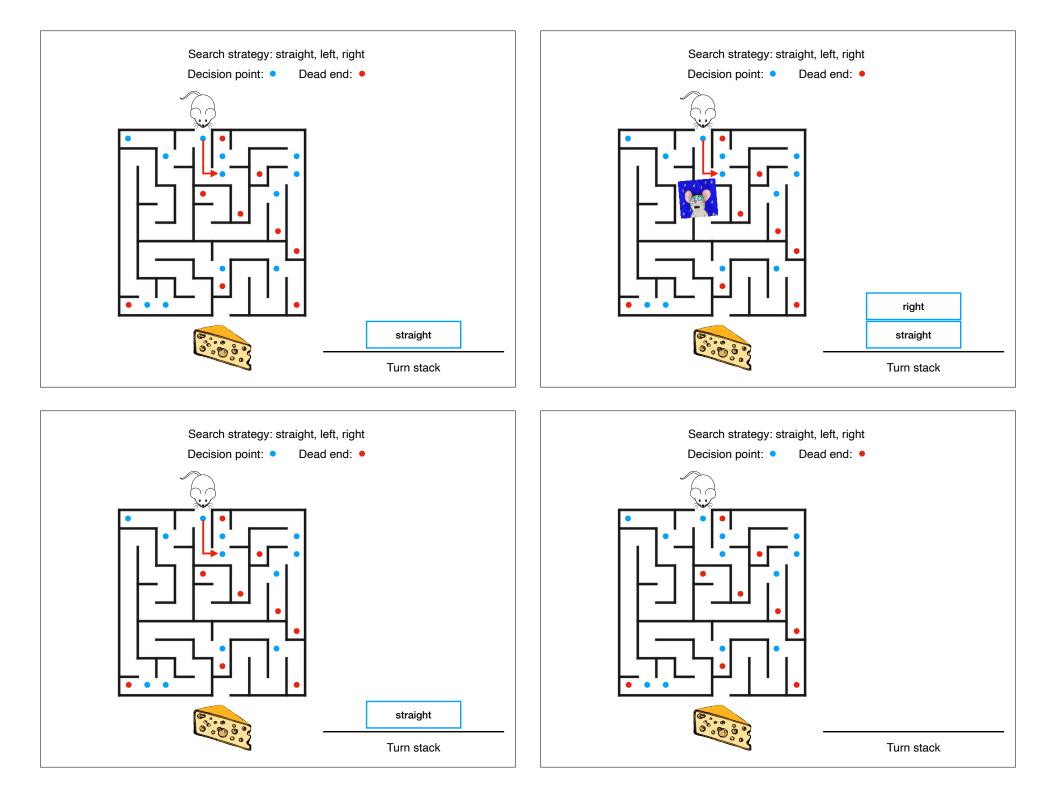


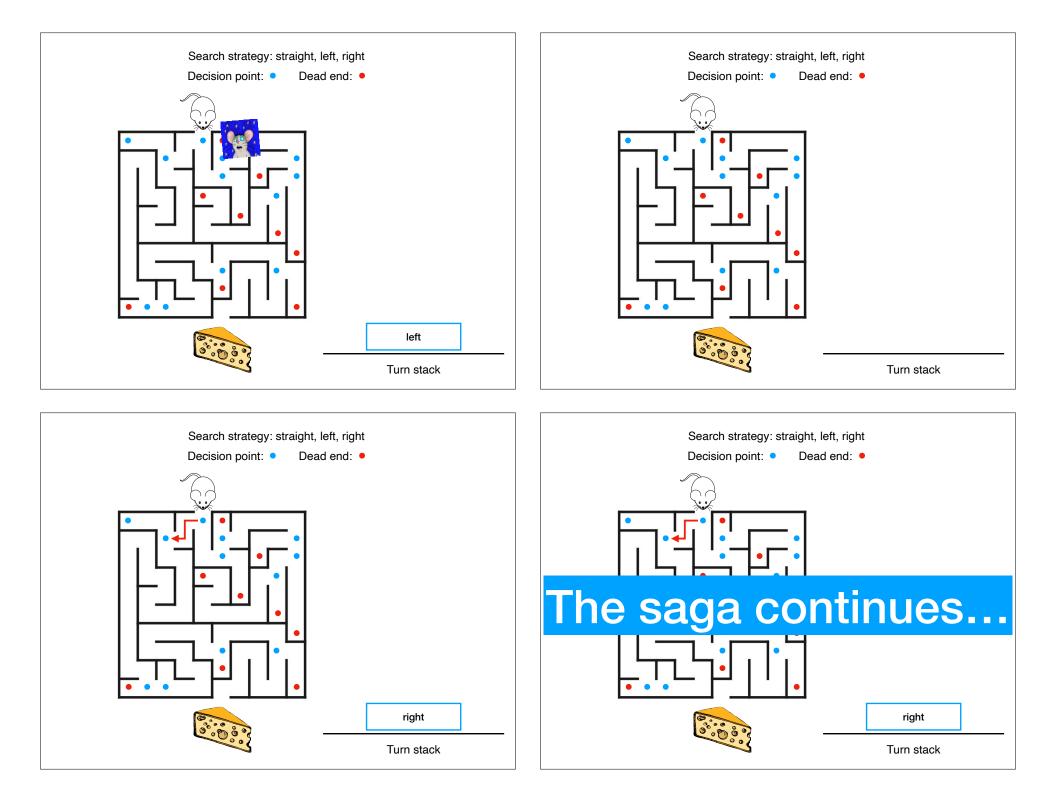












## Stack implementation

## Stack data structures

#### **StackArray**

A **StackArray** is a stack implemented using an **array** for element storage.

**Pros**: push and pop are O(1) operations.

**Cons**: data structure has a maximum **capacity**.

## Stack data structures

**StackVector** 

A **StackVector** is a stack implemented using a **Vector** for element storage.

**Pros**: **push** and **pop** are amortized **O(1)** operations. There is no maximum capacity.

**Cons**: Most of the time, ops take O(1) time, but occasionally (when the underlying array needs to grow) an O(n) cost is incurred. This may be fine for most applications, but if the application cannot tolerate wide variation in time, this is a bad choice.

Also, unless the underlying array is completely full, Vectors **waste some space**.

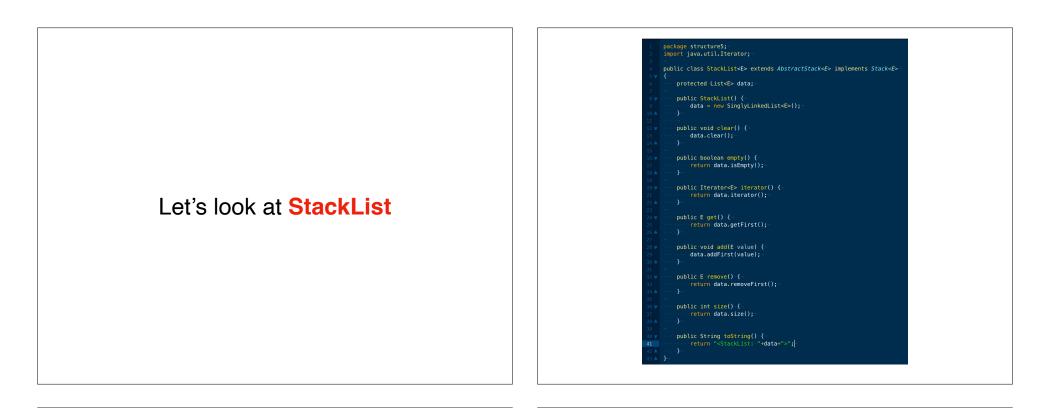
## Stack data structures

#### StackList

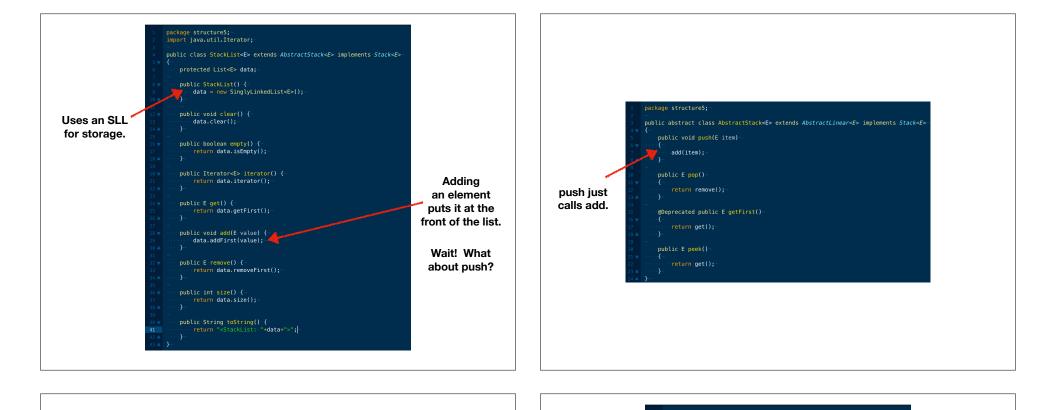
A **StackList** is a stack implemented using a **List** (usu. **SLL**) for element storage.

**Pros:** push and pop are O(1) operations. There is no maximum capacity, and no wasted space. push and pop costs are predictable (always the same), unlike **StackVector**.

**Cons**: because of the way computer hardware is implemented, a **StackList**'s constant-time cost is likely to be much higher than a **StackVector**'s. So a **StackList**'s performance may be more predictable than a **StackVector**, but it will likely be **slower on average**.









Adding an element puts it at the front of the list.

# Recap & Next Class

# Today:

Linear ADTs

Stack

## **Next class:**

Queue, etc.