

**CSCI 136**  
**Data Structures &**  
**Advanced Programming**

**Simple Sorting Algorithms**

# Introduction to Sorting

- Along with search, sorting is among the most ubiquitous computations
- Simple Examples
  - Reordering a list of integers so that they are in increasing order
  - Reordering a list of strings so that they are in alphabetical (*lexicographic*) order
  - Ordering the results of a Google search by decreasing relevance

# Introduction to Sorting

- More Examples
  - Reordering a list of strings so that they are in *lexicographical order by length*
    - Shorter strings come before longer strings and
    - Equal-length strings are in lexicographic order
  - Ordering tracks in a music collection by artist
    - Breaking ties by title
      - Breaking remaining ties by release date
- If you can order it, you can sort it!
- So, let's look at some simple sorting methods

# Introduction to Sorting

In this video we will

- Introduce three simple sorting algorithms
  - Bubble Sort, Insertion Sort, and Selection Sort
- Discuss their implementation
- Discuss their time and space complexity
  - Identify some sorting-specific measures of complexity
- Introduce some notation for describing *lower bounds*

# Bubble Sort

- First Pass:
  - ( **5** 1 3 2 9 ) → ( 1 **5** 3 2 9 )
  - ( 1 **5** 3 2 9 ) → ( 1 3 **5** 2 9 )
  - ( 1 3 **5** 2 9 ) → ( 1 3 2 **5** 9 )
  - ( 1 3 2 **5** 9 ) → ( 1 3 2 5 9 )
- Second Pass:
  - ( **1** 3 2 5 9 ) → ( **1** 3 2 5 9 )
  - ( 1 **3** 2 5 9 ) → ( 1 2 **3** 5 9 )
  - ( 1 2 **3** 5 9 ) → ( 1 2 3 5 9 )
- Third Pass:
  - ( **1** 2 3 5 9 ) → ( **1** 2 3 5 9 )
  - ( 1 **2** 3 5 9 ) → ( 1 **2** 3 5 9 )
- Fourth Pass:
  - ( **1** 2 3 5 9 ) → ( **1** 2 3 5 9 )
- Finished!

<http://www.visualgo.net/sorting>

<http://www.youtube.com/watch?v=lyZQPjUT5B4>

# Bubble Sort

Bubble sort uses a utility method swap

```
private static void swap(int[]A, int i, int j) {
    int temp = a[i];
    A[i] = A[j];
    A[j] = temp;
}

public static void bubbleSort(int[] A) {
    for(int i = 1; i < A.length; i++)
        // Process all but last i-1 elements
        for(int j = 0; j < A.length -i; j++)
            if( A[j] > A[j+1] ) swap(A,j,j+1);
}
```

The only subtlety: Do loops start and end at reasonable points!

# Bubble Sort

- Repeatedly scans through the list to be sorted, comparing two items at a time and swapping them if they are in the wrong order
  - Works on smaller initial slice each time
  - Can be improved to stop after a "swap-free" scan
- Gets its name from the way larger elements "bubble" to the end of the list
- Time complexity?
  - $O(n^2)$  : Might perform  $O(n^2)$  compares *and*  $O(n^2)$  swaps
- Space complexity?
  - $O(n)$  total (very little additional space is required)
- It's a *Stable* sorting method
  - Equal elements remain in same relative positions

# Bubble Sort

```
public static void bubbleSort(int[] A) {  
    for(int i = 1; i < A.length; i++)  
        // Process all but last i-1 elements  
        for(int j = 0; j < A.length - i; j++)  
            if( A[j] < A[j+1] ) swap(A, j, j+1);  
}
```

---

## Counting Operations (where $n = A.length$ )

- Outer loop executes  $n-1$  times
  - for  $i^{\text{th}}$  iteration: inner loop executes  $n-i$  times
    - Performing, say at most, say, 5 operations each time

$$\sum_{i=1}^{n-1} 5(n-i) = \sum_{i=1}^{n-1} 5n - \sum_{i=1}^{n-1} 5i = 5n(n-1) - 5 \sum_{i=1}^{n-1} i = 5n(n-1) - 5n(n-1)/2$$

- This equals  $\frac{5n(n-1)}{2}$  which is  $O(n^2)$



# Aside: Lower Bound Notation

There are situations in which bubble sort must necessarily perform a quadratic number of operations.

- Any "almost reverse-sorted" list will cause this

This observation describes a lower bound on the (worst-case) running time of the algorithm

- It's useful to have notation for lower-bound claims, similar to the Big-O notation for upper bound
- It exists: It's called "Big- $\Omega$ " (Big Omega) notation

# Aside: Lower Bound Notation

**Definition:** A function  $f(n)$  is  $\Omega(g(n))$  if for some constant  $c > 0$  and all  $n \geq n_0$

$$f(n) \geq c g(n)$$

So,  $f(n)$  is  $\Omega(g(n))$  exactly when  $g(n)$  is  $O(f(n))$

All three sorting algorithms have time complexity

- $O(n^2)$  : Never use more than  $cn^2$  operations
- $\Omega(n^2)$  : Sometimes use at least  $cn^2$  operations

When  $f(n)$  is  $O(g(n))$  and  $f(n)$  is  $\Omega(g(n))$  we write:

$$f(n) \text{ is } \Theta(g(n))$$

(pronounced "Big-Theta")

# Bubble Sort Complexity

## Time complexity?

- $\Theta(n^2)$  : That is, both  $O(n^2)$  and  $\Omega(n^2)$ 
  - $O(n^2)$  : Never performs more than  $c n^2$  operations
  - $\Omega(n^2)$  : Sometimes uses at least  $cn^2$  operations
    - Might perform  $O(n^2)$  compares *and*  $O(n^2)$  swaps

## Space complexity?

- $\Theta(n)$  : That is, both  $O(n)$  and  $\Omega(n)$ 
  - $\Omega(n)$  : *Needs to store* an  $n$ -element array of constant-sized values
  - $O(n)$  : *Only stores* the array *plus a few* other constant-sized variables

# Insertion Sort

- 5 7 0 3 4 2 6 1
- 5 7 0 3 4 2 6 1
- 0 5 7 3 4 2 6 1
- 0 3 5 7 4 2 6 1
- 0 3 4 5 7 2 6 1
- 0 2 3 4 5 7 6 1
- 0 2 3 4 5 6 7 1
- 0 1 2 3 4 5 6 7

<http://www.visualgo.net/sorting>

# Insertion Sort

- Simple sorting algorithm that works by building a sorted list one entry at a time
- Less efficient on large lists than more advanced algorithms
- Advantages:
  - Simple to implement and efficient on small lists
  - Efficient on data sets which are already mostly sorted
- Time complexity (Worst Case):  $\Theta(n^2)$ 
  - $O(n^2)$  : Only perform  $O(n^2)$  compares and  $O(n^2)$  moves
  - $\Omega(n^2)$  : Could perform  $\Omega(n^2)$  compares and  $\Omega(n^2)$  moves
- Space complexity :  $\Theta(n)$
- Stable: Yes

# Selection Sort

- 11    3    27    5    16    Swap 27 with 16
- 11    3    16    5    27    Swap 16 with 5
- 11    3    5    16    27    Swap 11 with 5
- 5    3    11    16    27    Swap 5 with 3
- 3    5    11    16    27    Done!

- The algorithm works as follows:
  - Find the maximum value in the list
  - Swap it with the value in the last position
  - Repeat the steps above for remainder of the list (ending at the second to last position)
  - Continue on progressively smaller portions of array

# Selection Sort

- Similar to insertion sort
- Noted for its simplicity and performance advantages when compared to complicated algorithms
- Time complexity (Worst Case):  $\Theta(n^2)$ 
  - $O(n^2)$  : Only perform  $O(n^2)$  compares and  $O(n)$  moves
  - $\Omega(n^2)$  : Could perform  $\Omega(n^2)$  compares and  $\Omega(n)$  moves
- Space Complexity :  $\Theta(n)$
- Stable : Yes

# Implementation Details

Selection sort uses two utility methods

Uses a swap method

```
private static void swap(int[]A, int i, int j) {  
    int temp = a[i];  
    A[i] = A[j];  
    A[j] = temp;  
}
```

And a max-finding method

```
// Find position of largest value in A[0 .. last]  
public static int findPosOfMax(int[] A, int last) {  
    int maxPos = 0;    // A wild guess  
    for(int i = 1; i <= last; i++)  
        if (A[maxPos] < A[i]) maxPos= i;  
    return maxPos;  
}
```



# Iterative & Recursive Selection Sort

## An Iterative Selection Sort

```
public static void selectionSort(int[] A) {  
    for(int i = A.length - 1; i>0; i--)  
        int big= findPosOfMax(A,i);  
        swap(A, i, big);  
    }  
}
```

## A Recursive Selection Sort (just the helper method)

```
public static void recSSHelper(int[] A, int last) {  
    if(last == 0) return; // base case  
  
    int big= findPosOfMax(A, last);  
    swap(A,big,last);  
    recSSHelper(A, last-1);  
}
```

# Notes

Three simple algorithms: Bubble, Insertion, Selection

- All perform two basic operations: *comparisons* and *swaps/moves*

Comparisons vs Swaps/Moves (worst case)

- Bubble Sort performs (naïve version)
  - quadratic number of comparisons in all cases
  - quadratic number of swaps for almost reverse-sorted data
- Insertion Sort performs
  - linear number of comparisons on almost-sorted data and quadratic number on almost reverse-sorted data
  - quadratic number of moves on almost reverse-sorted data
- Selection Sort
  - quadratic number of comparisons in all cases
  - at most (and sometimes) a linear number of swaps

# Coming Up Next

How can we adapt our sorting algorithms to work on non-primitive data types?

Can we break the  $\Omega(n^2)$  performance bottleneck of our sorting algorithms?

Spoiler: Yes!

Stay tuned....