

CSCI 136
Data Structures &
Advanced Programming

Lecture 15

Fall 2019

Instructor: B&S

Announcements

- Mid-Term Review Session
 - Monday, Oct. 14 from 9:00-11:00 am
 - No prepared remarks, so bring questions!
- Mid-term exam is Wednesday, October 16
 - During your normal lab session
 - You'll have 1 hour & 45 minutes (if you come on time!)
 - Closed-book
 - Covers Chapters 1-7 & 9 and all topics up through Linked Lists
 - A “sample” mid-term and study sheet are available online
 - [See Handouts & Problem Sets](#)

Last Time : Linear Structures

- Stack applications
 - Arithmetic Expressions
 - Postscript
 - Mazerunning (Depth-First-Search)

Today: Linear Structures

- Stacks
 - (Implicit) program call stack
- Queues
 - Implementations Details
 - Applications
- Iterators

Recursive “Pseudo-Code” Sketch

Boolean RecSolve(Maze m, Position current)

If (current equals finish) return true

Mark current as visited

next ← some unvisited neighbor of current (or null if none left)

While (next does not equal null && recSolve(m, next) is false)

next ← some unvisited neighbor of current (or null if none left)

Return next ≠ null

- To solve maze, call: *Boolean recSolve(m, start)*
- To prove correct: Induction on distance from *current* to *finish*
- How could we generate the actual solution?

Method Call Stacks

- In JVM, need to keep track of method calls
- JVM maintains stack of method invocations (called frames)
- Stack of frames
 - Receiver object, parameters, local variables
- On method call
 - Push new frame, fill in parameters, run code
- Exceptions print out stack
- Example: StackEx.java
- Recursive calls recurse too far: StackOverflowException
 - Overflow.java

Stacks vs. Queues

- Stacks are LIFO (Last In First Out)
 - Methods: push, pop, peek, empty
 - Sample Uses:
 - Evaluating expressions (postfix)
 - Solving mazes
 - Evaluating postscript
 - JVM method calls
- Queues are FIFO (First In First Out)
 - Another linear data structure (implements Linear interface)
 - Queue interface methods: enqueue (add), dequeue (remove), getFirst (get), peek (get)

Queues



- Examples:
 - Lines at movie theater, grocery store, etc
 - OS event queue (keeps keystrokes, mouse clicks, etc, in order)
 - Printers
 - Routing network traffic (more on this later)

Queue Interface

```
public interface Queue<E> extends Linear<E> {  
    public void enqueue(E item);  
    public E dequeue();  
    public E getFirst(); //value not removed  
    public E peek(); //same as get()  
}
```

Implementing Queues

As with Stacks, we have three options:

QueueArray

```
class QueueArray<E> implements Queue<E> {  
    protected Object[] data; //can't declare E[]  
    int head;  
    int count; // better than storing tail...  
}
```

QueueVector

```
class QueueVector<E> implements Queue<E> {  
    protected Vector<E> data;  
}
```

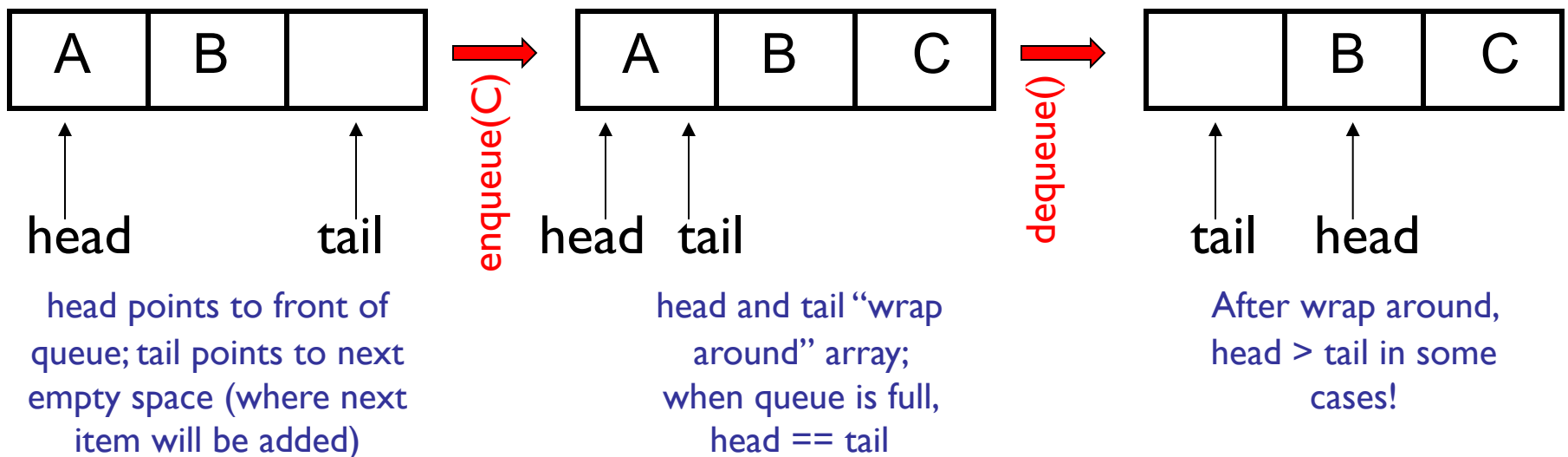
QueueList

```
class QueueList<E> implements Queue<E> {  
    protected List<E> data; //uses a CircularList  
}
```

All three of these also extend **AbstractQueue**

QueueArray

- Let's look at an example...
- How to implement?
 - enqueue(item), dequeue(), size()



```
public class queueArray<E> {

    protected Object[] data;          // Must use object because...
    protected int head;
    protected int count;

    public queueArray(int size) {
        data = new Object[size]; // ... can't say "new E[size]"
    }

    public void enqueue(E item) {
        Assert.pre(count<data.length,"Queue is full.");
        int tail = (head + count) % data.length;
        data[tail] = item;
        count++;
    }

    public E dequeue() {
        Assert.pre(count>0,"The queue is empty.");
        E value = (E)data[head];
        data[head] = null;
        head = (head + 1) % data.length;
        count--;
        return value;
    }

    public boolean empty() {
        return count>0;
    }
}
```

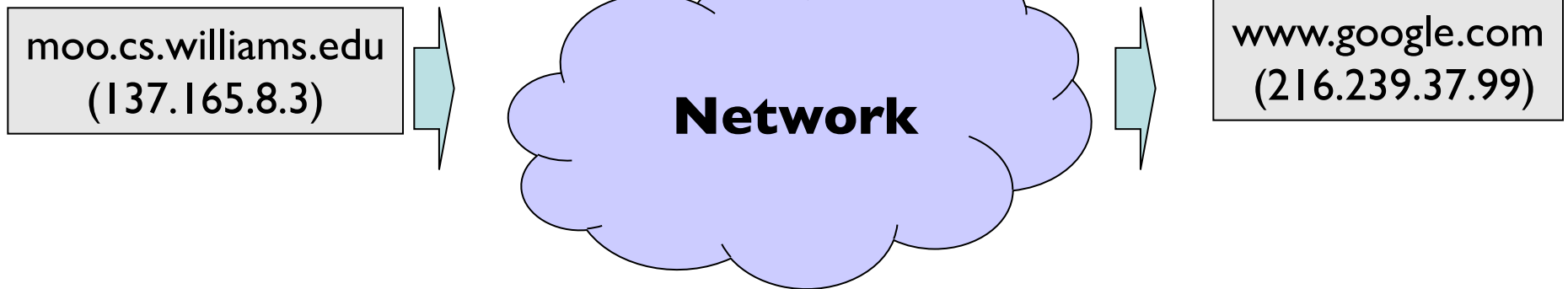
Tradeoffs:

- **QueueArray:**
 - enqueue is $O(1)$
 - dequeue is $O(1)$
 - Faster operations, but limited size
- **QueueVector:**
 - enqueue is $O(1)$ (but $O(n)$ in worst case - ensureCapacity)
 - dequeue is $O(n)$
- **QueueList:**
 - enqueue is $O(1)$ (addLast)
 - dequeue is $O(1)$ (CLL removeFirst)

Routing With Queues

Slides by Stephen Freund

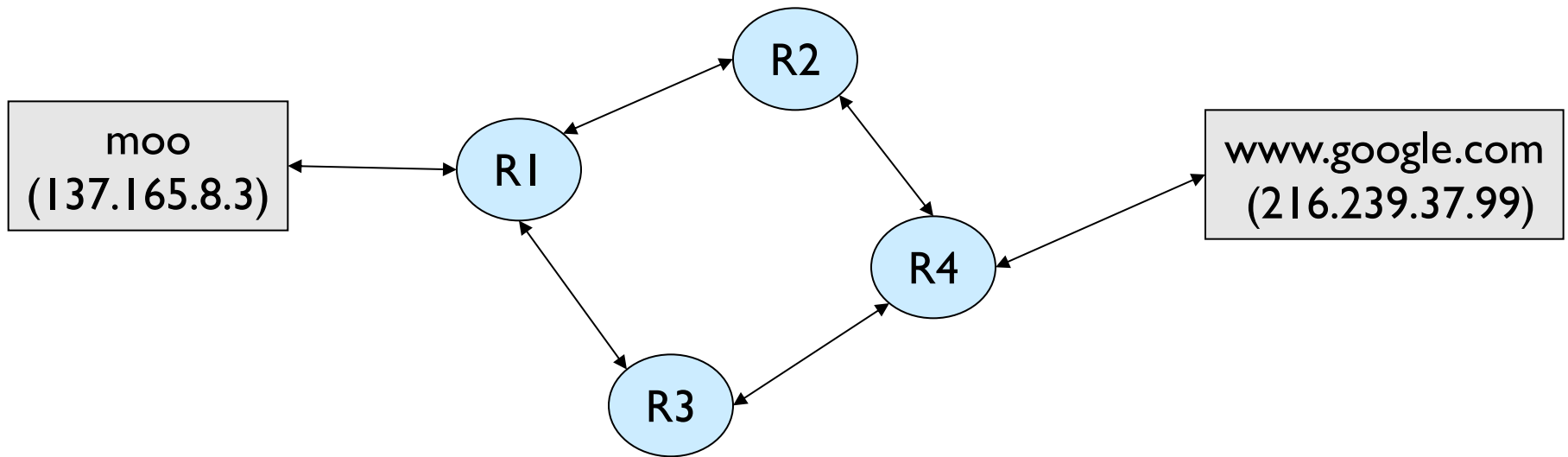
The Network



Message:

137.165.8.3	216.239.37.99	"Search for ..."
-------------	---------------	------------------

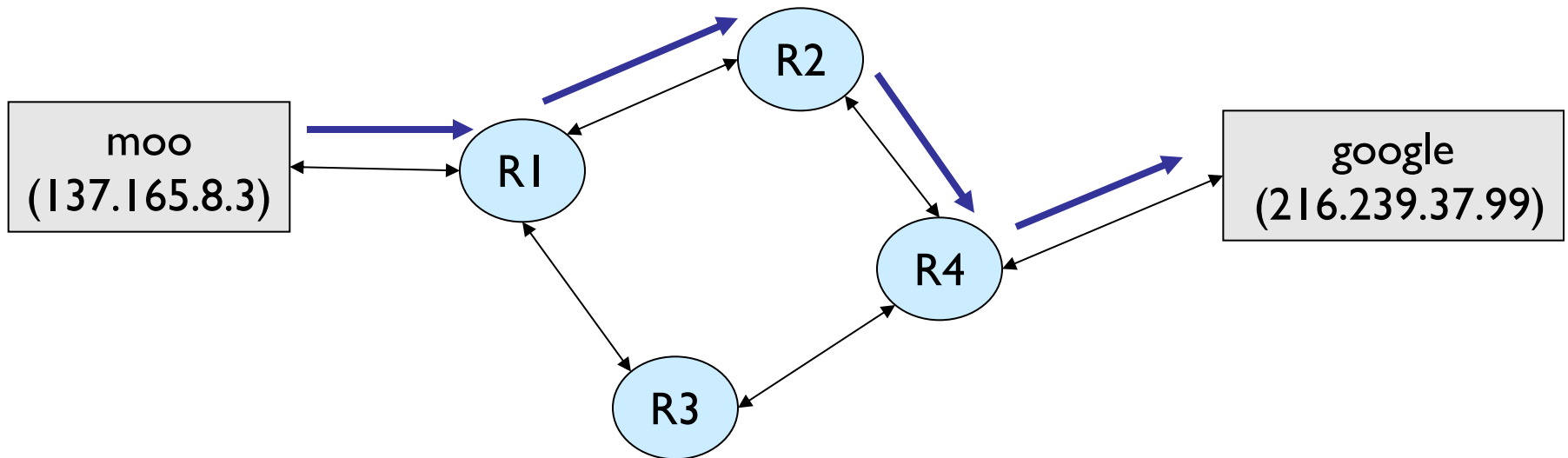
Routers



Message:

137.165.8.3	216.239.37.99	"Search for ..."
-------------	---------------	------------------

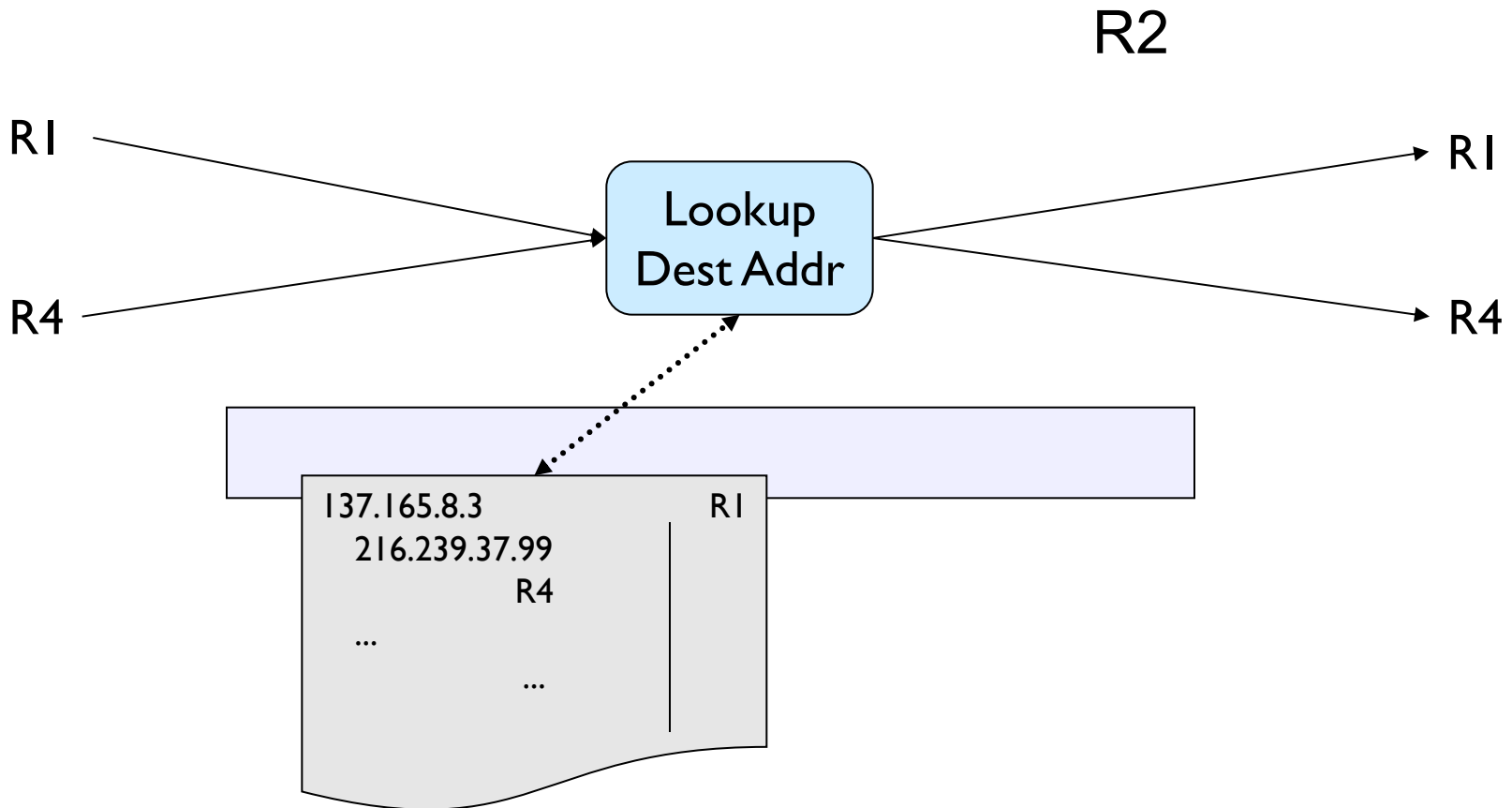
Routers



Routing Algorithm

1. Receive message
2. Look up Destination Address
 - a) Deliver message to Dest
 - b) Forward to next Router

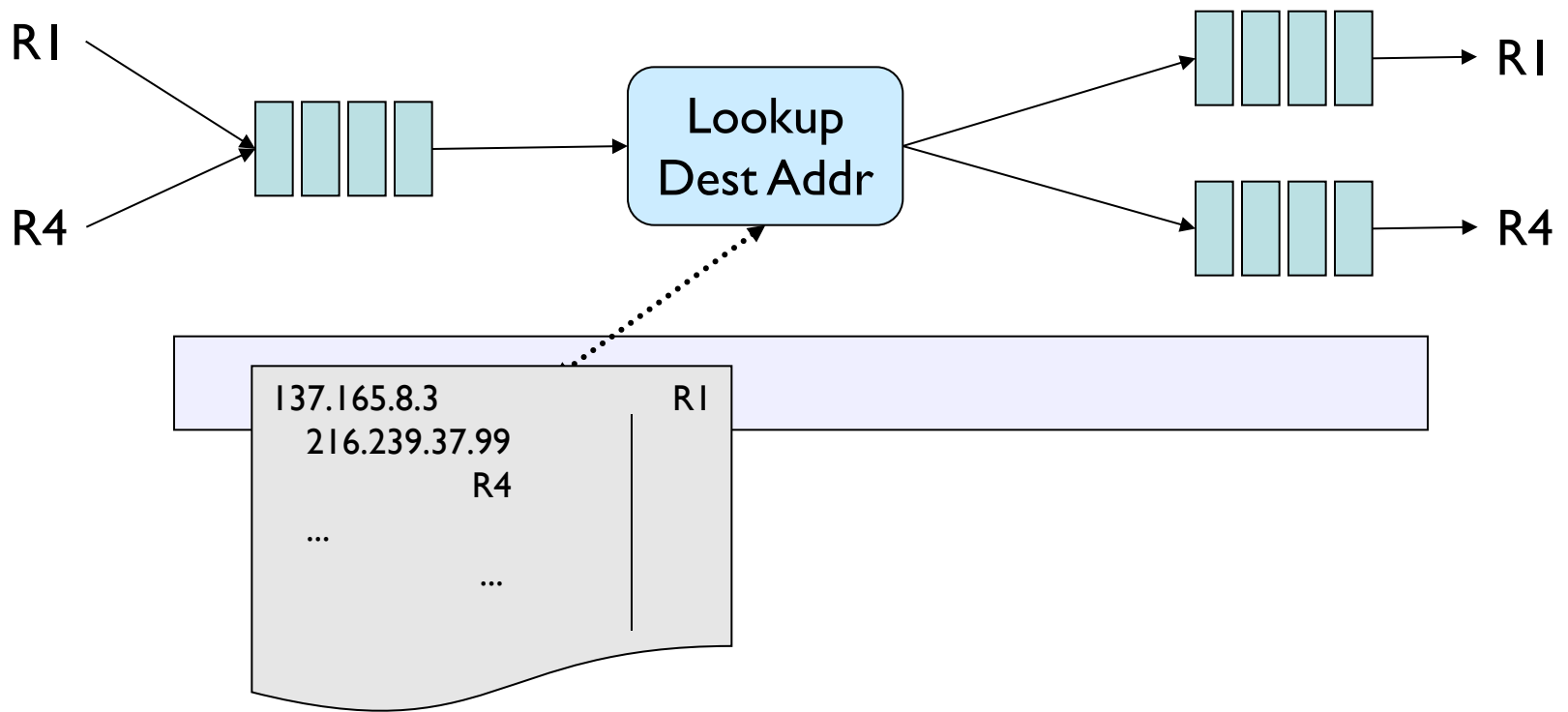
Router Internals



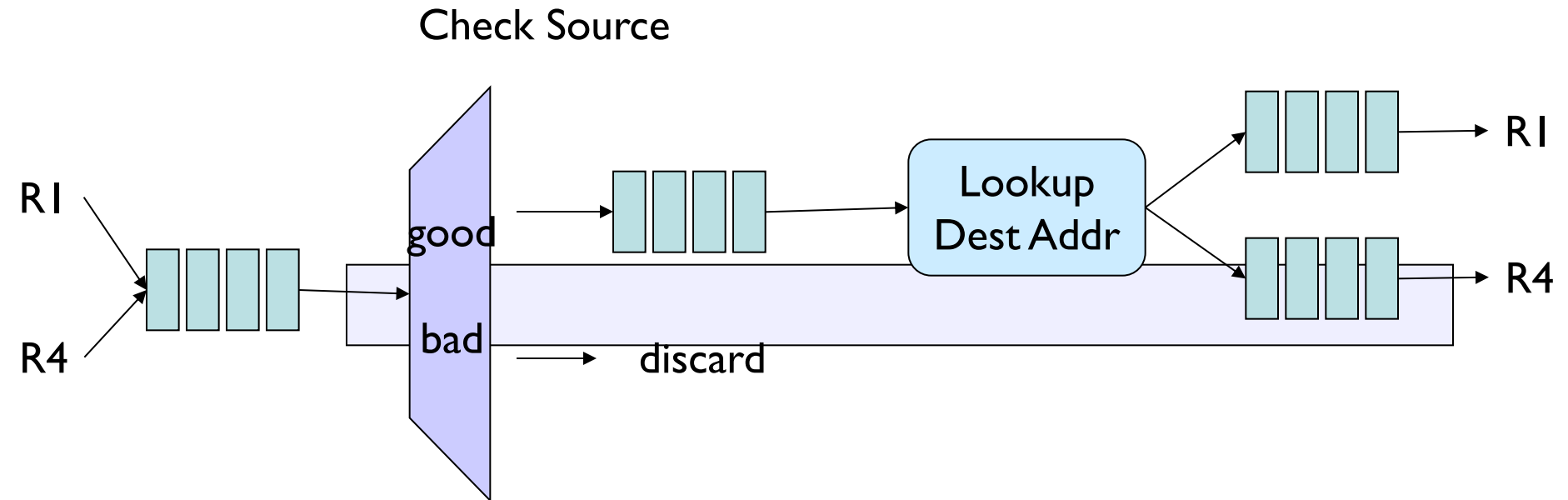
Buffering Messages

- There may be delays
 - Router receives messages faster than it can process and send
 - Some links are slower than others
 - Common speeds: 10 Mbs, 100Mbs, 1Gbs.
 - Wireless, satellite, infra-red, telephone line, ...
 - Hardware problems
- Want to be able to handle short-term congestion problems

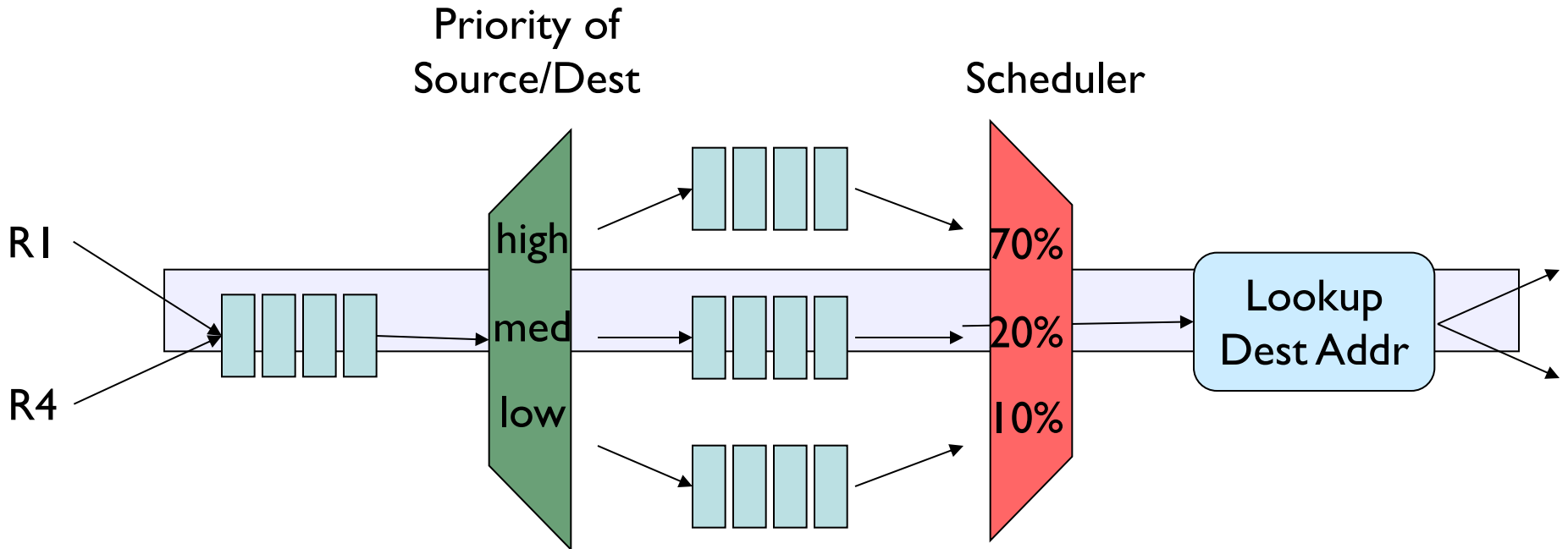
Router Internals



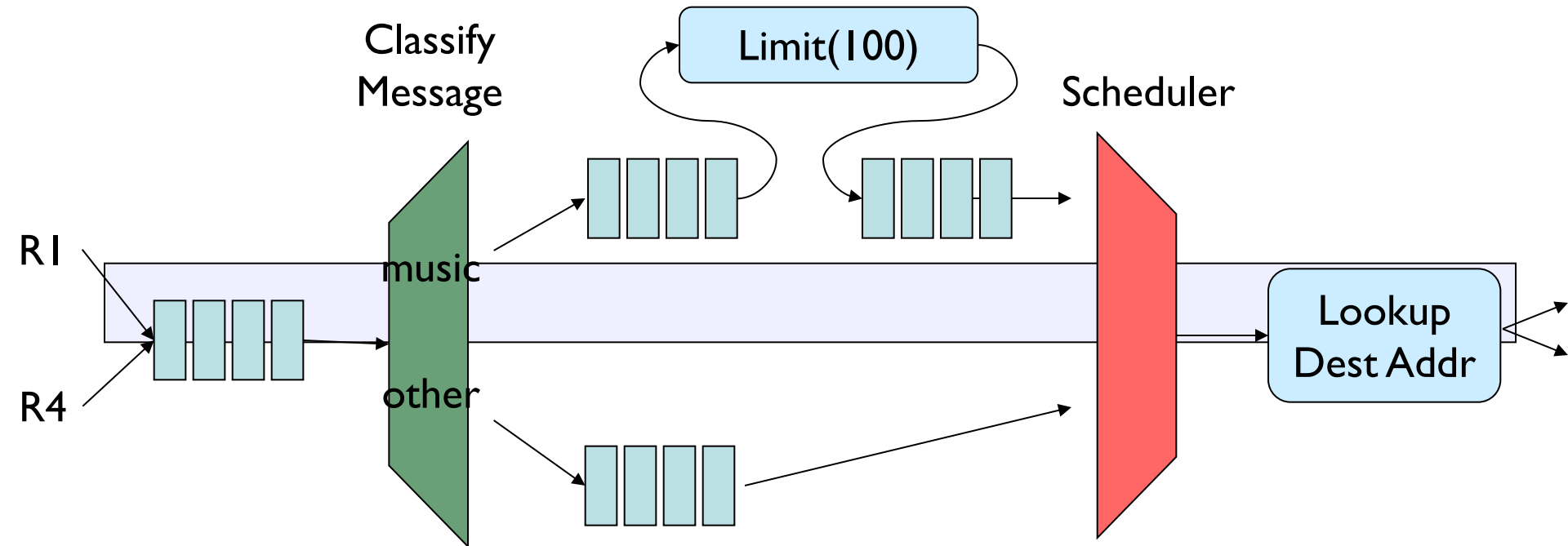
Firewalls



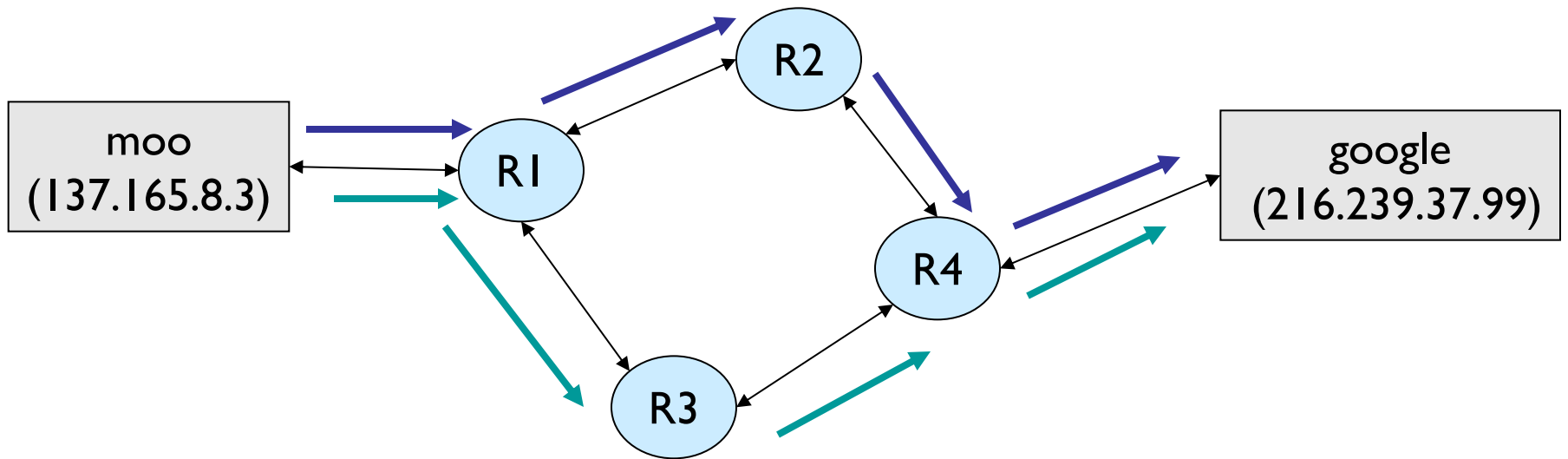
Priority Scheduling



Bandwidth Shaper



Choosing The Best Route



Choosing Routes

- Routers exchange information periodically
 - Attempt to route on "best" path to destination
 - Not easy to determine:
 - Network congestion varies (evening vs. morning)
 - Hardware added/removed or failures
- Dijkstra's algorithm (later)

Visiting Data from a Structure

- Write a method (`numOccurs`) that counts the number of times a particular `Object` appears in a structure

```
public int numOccurs (List data, E o) {
    int count = 0;
    for (int i=0; i<data.size(); i++) {
        E obj = data.get(i);
        if (obj.equals(o)) count++;
    }
    return count;
}
```

- Does this work on all structures (that we have studied so far)?

Problems

- `get()` not defined on Linear structures (i.e., stacks and queues)
- `get()` is “slow” on some structures
 - $O(n)$ on SLL (and DLL)
 - So `numOccurs` = $O(n^2)$ for linked lists
- How do we traverse data in structures in a general, efficient way?
 - Goal: data structure-specific for efficiency
 - Goal: use same interface to make general

Recall : Structure Operations

- `size()`
- `isEmpty()`
- `add()`
- `remove()`
- `clear()`
- `contains()`

- But also
 - Method for efficient data traversal
 - `iterator()`

Iterators

- **Iterators** provide support for *efficiently* visiting all elements of a data structure
- An Iterator:
 - Provides generic methods to dispense values
 - Traversal of elements : *Iteration*
 - Production of values : *Generation*
 - Abstracts away details of how elements are retrieved
 - Uses different implementations for each structure

```
public interface Iterator<E> {  
    boolean hasNext() – are there more elements in iteration?  
    E next() – return next element  
    default void remove() – removes most recently returned value
```

- Default : Java provides an implementation for remove
 - It throws an UnsupportedOperationException exception

A Simple Iterator

- Example: FibonacciNumbers

```
public class FibonacciNumbers implements Iterator<Integer> {
    private int next= 1, current = 1;
    private int length= 10; // Default

    public FibonacciNumbers() {}
    public FibonacciNumbers(int n) {length= n;}
    public boolean hasNext() { return length>=0;}
    public Integer next() {
        length--;
        int temp = current;
        current = next;
        next = temp + current;
        return temp;
    }
}
```

Why Is This Cool? (it is)

- We could calculate the i^{th} Fibonacci number each time, but that would be slow
 - Observation: to find the n^{th} Fib number, we calculate the previous $n-1$ Fib numbers...
 - But by storing some state, we can easily generate the next Fib number in $O(1)$ time
- Knowledge about the structure of the problem helps us traverse the Fib space *efficiently* one element at a time
 - Let's do the same for data structures

Iterators Of Structures

Goal: Have data structures produce iterators that return the values of the structure in some order.

How?

- Define an iterator class for the structure, e.g.

```
public class VectorIterator<E>  
    implements Iterator<E>;
```

```
public class SinglyLinkedListIterator<E>  
    implements Iterator<E>;
```

- Provide a method in the structure that returns an iterator

```
public Iterator<E> iterator() { ... }
```


Iterators Of Structures

The details of `hasNext()` and `next()` depend on the specific data structure, e.g.

- `VectorIterator` holds an array reference and index of next element
 - A reference to the data array of the `Vector`
 - The index of the next element whose value to return
- `SinglyLinkedListIterator` holds
 - a reference to the head of the list
 - A reference to the next node whose value to return

Iterator Use : numOccurs

```
public int numOccurs (List<E> data, E o) {
    int count = 0;
    Iterator<E> iter = data.iterator();
    while (iter.hasNext())
        if(o.equals(iter.next())) count++;
    return count;
}
// Or...
```

```
public int numOccurs (List<E> data, E o) {
    int count = 0;
    for(Iterator<E> i = data.iterator();
        i.hasNext();)
        if(o.equals(i.next())) count++;
    return count;
}
```

Implementation Details

- We use both an Iterator interface and an AbstractIterator class
- All concrete classes in structure5 extend AbstractIterator
 - AbstractIterator partially implements Iterator
- Importantly, AbstractIterator *adds* two methods
 - get() – peek at (but don't take) next element, and
 - reset() – reinitialize iterator for reuse
- Methods are specialized for each data structure

Iterator Use : numOccurs

Using an AbstractIterator allows more flexible coding
(but requiring a cast to AbstractIterator)

Note: It has the form of a standard 3-part for statement

```
public int numOccurs (List<E> data, E o) {
    int count = 0;
    for(AbstractIterator<E> i =
        (AbstractIterator<E>) data.iterator();
        i.hasNext(); i.next())
        if(o.equals(i.get())) count++;
    return count;
}
```

Implementation : SLLiterator

```
public class SinglyLinkedListIterator<E> extends AbstractIterator<E> {
    protected Node<E> head, current;

    public SinglyLinkedListIterator(Node<E> head) {
        this.head = head;
        reset();
    }

    public void reset() { current = head;}

    public E next() {
        E value = current.value();
        current = current.next();
        return value;
    }

    public boolean hasNext() { return current != null; }

    public E get() { return current.value(); }
}
```

In SinglyLinkedList.java:

```
public Iterator<E> iterator() {
    return new SinglyLinkedListIterator<E>(head);
}
```

More Iterator Examples

- How would we implement `VectorIterator`?
- How about `StackArrayIterator`?
 - Do we go from bottom to top, or top to bottom?
 - Doesn't matter! We just have to be consistent...
- We can also make “specialized” iterators
 - Another SLL Example: `Skipliterator.java`
 - `Reverseliterator.java`

Iterators and For-Each

Recall: with arrays, we can use a simplified form of the for loop

```
for( E elt : arr) {System.out.println( elt );}
```

Or, for example

```
// return number of times o appears in data
public int numOccurs (E[] data, E o) {
    int count = 0;
    for(E current : data)
        if(o.equals(current)) count++;
    return count;
}
```

We can do this with classes that provide an iterator() method...

The Iterable Interface

We can use the “for-each” construct...

```
for( E elt : boxOfStuff ) { ... }
```

...as long as `boxOfStuff` implements the *Iterable* interface

```
public interface Iterable<T> {  
    public Iterator<T> iterator();  
}
```

Since `Structure<E>` extends `Iterable<E>`, we can write

```
public int numOccurs (List<E> data, E o) {  
    int count = 0;  
    for(E current : data)  
        if(o.equals(current)) count++;  
    return count;  
}
```


General Rules for Iterators

1. Understand order of data structure
 2. **Always call hasNext() before calling next()!!!**
 3. Use remove with caution!
 1. Don't use remove....
 4. Don't add to structure while iterating: TestIterator.java
- Take away messages:
 - Iterator objects capture state of traversal
 - They have access to internal data representations
 - They should be fast and easy to use