

Heapsort

This handout goes through an example of how heapsort functions on the following array A .*

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
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Heapsort on an array of size n works in two steps. First, `heapify()` is called on the array to turn it into a heap. Then, for $i = n - 1$ to 0, `remove()` is called, and the resulting element is stored in slot i ; the heap is then considered to be one element smaller.

1 Heapify

As we discussed in class, we want to use Bottom-Up Heapify, as it runs in $O(n)$ time. To accomplish this, for $j = n - 1$ to 0, we call `pushDownRoot(j)`. On the right we show what the array looks like *after* the corresponding step, with arrows to indicate any swaps that were performed *during* that step. We highlight in red the nodes from j to $n - 1$; these nodes satisfy the heap property.

`pushDownRoot(14)`: The children of 14 are 29 and 30; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

`pushDownRoot(13)`: The children of 13 are 27 and 28; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

`pushDownRoot(12)`: The children of 12 are 25 and 26; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

`pushDownRoot(11)`: The children of 11 are 23 and 24; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

`pushDownRoot(10)`: The children of 10 are 21 and 22; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

`pushDownRoot(9)`: The children of 9 are 19 and 20; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

`pushDownRoot(8)`: The children of 8 are 17 and 18; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

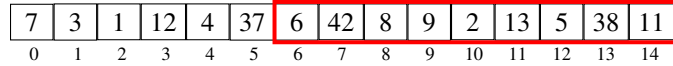
`pushDownRoot(7)`: The children of 7 are 15 and 16; both are not stored in the heap, so the heap property is satisfied.

7	3	1	12	4	37	6	42	8	9	2	13	5	38	11
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

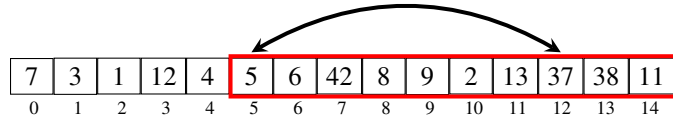
*We use an array in this example, but heapsort would work just as well on a vector.

A side note: at this point you may have noticed that we've wasted a lot of time looking at elements that do not have children. You're right! A good heapify implementation would skip right to elements with children: it would call $\text{pushDownRoot}(j)$ for $j = (n - 2)/2$ to 0.

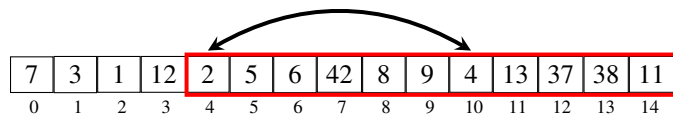
$\text{pushDownRoot}(6)$: The children of 6 are 13 and 14. $A[6] = 6$ is smaller than $A[13] = 38$ and $A[14] = 11$, so no swaps are necessary.



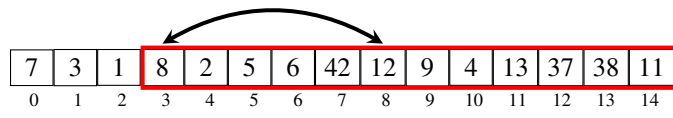
$\text{pushDownRoot}(5)$: The children of 5 are 11 and 12. $A[5] = 37$ is not smaller than $A[11] = 13$ (or $A[12] = 5$), so we need to swap. $\text{pushDown}()$ always swaps with the smaller child. Since $A[12] < A[11]$, we swap $A[5]$ with $A[12]$. The children of 12 are not in the heap so we are done.



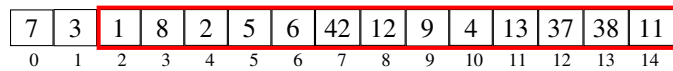
$\text{pushDownRoot}(4)$: The children of 4 are 9 and 10. $A[4] = 4$ is smaller than $A[9] = 9$, but is larger than $A[10] = 2$. We swap with the smaller child. The children of 10 are not in the heap so we are done.



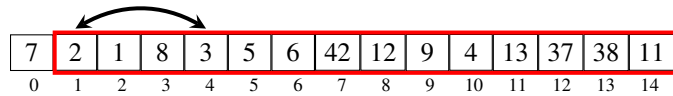
$\text{pushDownRoot}(3)$: The children of 3 are 7 and 8. $A[3] = 12$ is larger than $A[8] = 8$; we swap with the smaller child. The children of 8 are not in the heap so we are done.



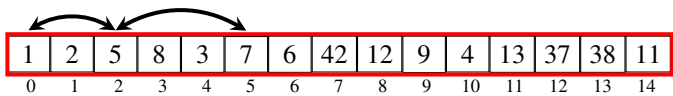
$\text{pushDownRoot}(2)$: The children of 2 are 5 and 6. $A[2]$ is smaller than $A[5]$ and $A[6]$, so we don't swap.



$\text{pushDownRoot}(1)$: The children of 1 are 3 and 4. $A[1]$ is larger than $A[4]$ and $A[4] < A[3]$ so we swap with $A[4]$. The children of 4 are 9 and 10. $A[4]$ is smaller than $A[9]$ and $A[10]$ so we don't swap.



$\text{pushDownRoot}(0)$: The children of 0 are 1 and 2. $A[0]$ is larger than $A[2]$ and $A[2] < A[1]$ so we swap with $A[2]$. The children of 2 are 5 and 6. $A[2] = 7$ is larger than $A[5] = 5$ and $A[5] < A[6] = 6$ so we swap with $A[5]$. After this, $A[5] < A[11]$ and $A[5] < A[12]$ so we are done.

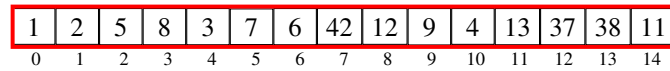


After $O(n)$ calls to pushDownRoot , we have a heap! As we showed in class, this actually takes $O(n)$ total work in the worst case. Now we need to sort.

2 Sort By Repeatedly Removing the Minimum

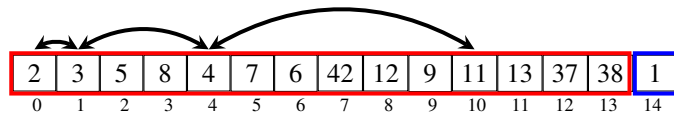
Now we sort. As discussed in class, for $i = n - 1$ to 0 , we remove the minimum element of the heap and place it in $A[i]$. This is particularly easy for a heap, as the first step in `remove()` is to swap the first and last item in the heap. After that, `pushDownRoot(0)` is called.

So after round i , the i smallest elements are in reverse-sorted order from $A[i]$ to $A[n-1]$; meanwhile $A[0]$ to $A[i-1]$ is a heap. Again we go round by round. We show the state of the array *after* each round, and show the swaps made during the round using arrows. Again, cells highlighted in red form a heap; cells highlighted in blue are in reverse-sorted order. We begin with the heap obtained from before:

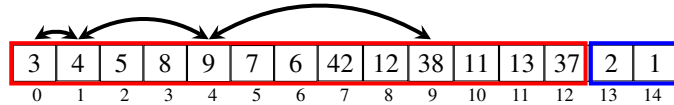


Now we begin our calls to `remove`.

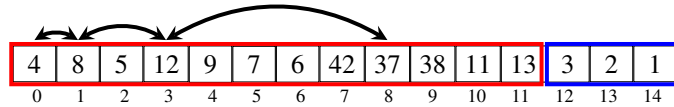
($i = 14$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[1]$, $A[1]$ is swapped with $A[4]$, and $A[4]$ is swapped with $A[10]$.



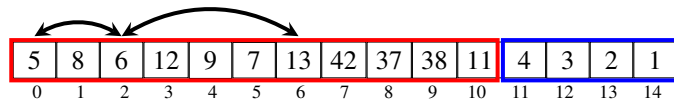
($i = 13$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[1]$, $A[1]$ is swapped with $A[4]$, and $A[4]$ is swapped with $A[9]$.



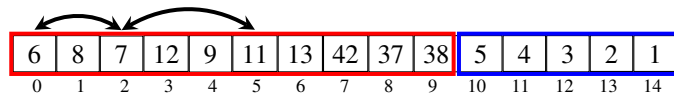
($i = 12$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[1]$, $A[1]$ is swapped with $A[3]$, and $A[3]$ is swapped with $A[8]$.



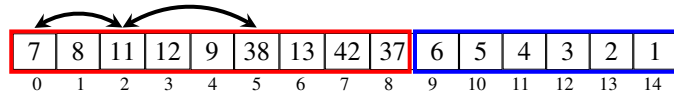
($i = 11$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[2]$ and $A[2]$ is swapped with $A[6]$.



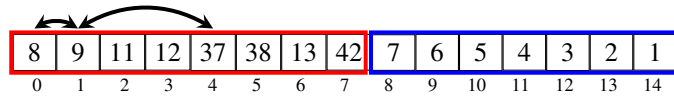
($i = 10$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[2]$ and $A[2]$ is swapped with $A[5]$.



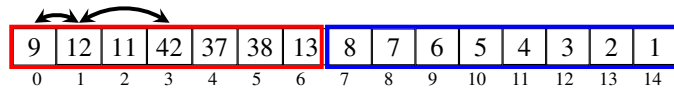
($i = 9$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[2]$ and $A[2]$ is swapped with $A[5]$.



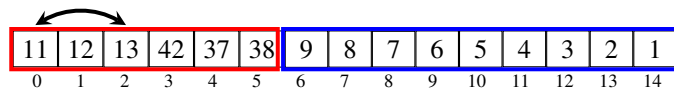
($i = 8$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[1]$ and $A[1]$ is swapped with $A[4]$.



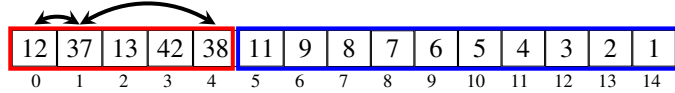
($i = 7$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[1]$ and $A[1]$ is swapped with $A[3]$.



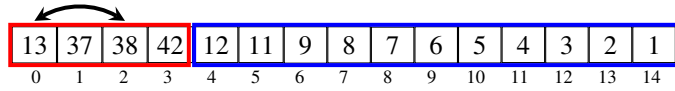
($i = 6$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[2]$.



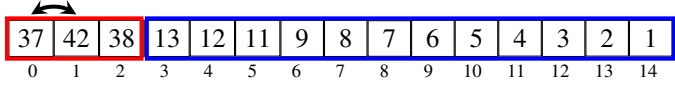
($i = 5$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[4]$, and $A[1]$ is swapped with $A[4]$.



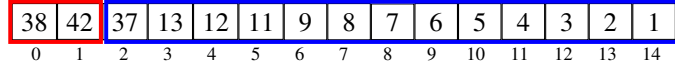
($i = 4$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[2]$.



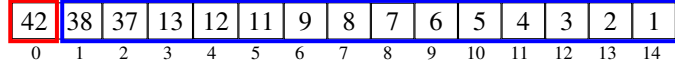
($i = 3$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. $A[0]$ is swapped with $A[1]$.



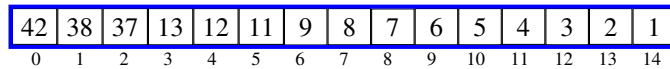
($i = 2$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. No swaps are needed.



($i = 1$): We swap $A[0]$ with $A[i]$ and call `Pushdown(0)`. No swaps are needed.



And we're done! Our array is in sorted order:



Well, it's in reverse sorted order. If we want, we can reverse it in $O(n)$ time.

Side note: There are techniques one can use to avoid this final step—for example, by using a “max heap,” which is exactly like a heap but with the property that each node is *larger* than its children. Calling `remove()` on a max heap gives the largest node, ultimately resulting in a sorted array with no reversals needed.